# **Protection**

Chapter 14

March 29, 2008

**OS:** Protection

1



- Discuss the goals and principles of protection in a modern computer system
- Explain how protection domains combined with an access matrix are used to specify the resources a process may access
- Examine capability and language-based protection systems

#### Chapter Outline

- Goals and Principles of Protection
- Policy and Mechanism
- Domain of Protection
- Access Matrix
- Implementation of Access Matrix
- Access Control
- Revocation of Access Rights
- Capability-Based Systems (skip)
- Language-Based Protection (skip)

### - Goals and Principles of Protection

- Goals of Protection
  - Operating system consists of a collection of objects, hardware or software
  - Each object has a unique name and can be accessed through a well-defined set of operations.
  - Protection problem ensure that each object is accessed correctly and only by those processes that are allowed to do so.
- Guiding principle principle of least privilege
  - Programs, users and systems should be given just enough privileges to perform their tasks



- Good to separate protection policy from mechanism
- Policy
  - User dictates policy.
  - Who can access what object and in what mode.

#### Mechanism

- Operating system provides access-matrix + rules.
- It ensures that the matrix is only manipulated by authorized agents and that rules are strictly enforced.



- The ability to execute an operation on an object is an access-right
- Access-right = < object-name, rights-set> where rights-set is a subset of all valid operations that can be performed on the object.
- **Domain** = set of access-rights



## -- Example: Domain Implementation (UNIX)

- System consists of 2 domains:
  - User
  - Supervisor
- UNIX
  - Domain = user-id
  - Domain switch accomplished via file system (dynamic).
    - Each file has associated with it a domain bit (setuid bit).
    - When file is executed and setuid = on, then user-id is set to owner of the file being executed. When execution completes user-id is reset.

### -- Example: Domain Implementation (MULTICS)

- Let  $D_i$  and  $D_j$  be any two domain rings.
- If  $j < I \Rightarrow D_j \subseteq D_j$





- View protection as a matrix (*access matrix*)
- Rows represent domains
- Columns represent objects
- Access(i, j) is the set of operations that a process executing in Domain<sub>i</sub> can invoke on Object<sub>i</sub>
- Access matrix design separates mechanism from policy.

#### ... - Access Matrix (static)





- If a process in Domain D<sub>i</sub> tries to do "op" on object O<sub>j</sub>, then "op" must be in the access matrix.
- Can be expanded to **dynamic** protection.
  - Operations to add, delete access rights.
  - Special access rights:
    - owner of  $O_i$
    - **copy** op from  $O_i$  to  $O_j$
    - control D<sub>i</sub> can modify D<sub>j</sub> access rights
    - *switch transfer from domain* D<sub>i</sub> to D<sub>j</sub>



object domain	F <sub>1</sub>	$F_2$	F <sub>3</sub>	laser printer	<i>D</i> <sub>1</sub>	D <sub>2</sub>	<i>D</i> <sub>3</sub>	$D_4$
<i>D</i> <sub>1</sub>	read		read			switch		
$D_2$				print			switch	switch
<i>D</i> <sub>3</sub>		read	execute					
$D_4$	read write		read write		switch			

### -- Access Matrix with Copy Rights

object domain	F <sub>1</sub>	$F_2$	F <sub>3</sub>			
<i>D</i> <sub>1</sub>	execute		write*			
D <sub>2</sub>	execute	read*	execute			
<i>D</i> <sub>3</sub>	execute					
(a)						
object domain	F <sub>1</sub>	$F_2$	F <sub>3</sub>			
<i>D</i> <sub>1</sub>	execute		write*			
D <sub>2</sub>	execute	read*	execute			
<i>D</i> <sub>3</sub>	execute	read				
(b)						

#### **OS:** Protection

#### -- Access Matrix With Owner Rights

object domain	F <sub>1</sub>	F <sub>2</sub>	F <sub>3</sub>			
D <sub>1</sub>	owner execute		write			
D <sub>2</sub>		read* owner	read* owner write			
<i>D</i> <sub>3</sub>	execute					
(a)						
object domain	F <sub>1</sub>	F <sub>2</sub>	$F_3$			
D <sub>1</sub>	owner execute		write			
D <sub>2</sub>		owner read* write*	read* owner write			
<i>D</i> <sub>3</sub>		write	write			
(b)						

#### **OS:** Protection



object domain	F <sub>1</sub>	F <sub>2</sub>	F <sub>3</sub>	laser printer	<i>D</i> <sub>1</sub>	D <sub>2</sub>	<i>D</i> <sub>3</sub>	$D_4$
<i>D</i> <sub>1</sub>	read		read			switch		
D <sub>2</sub>				print			switch	switch control
$D_3$		read	execute					
$D_4$	write		write		switch			



- Protection can be applied to non-file resources
- Solaris 10 provides role-based access control to implement least privilege
  - Privilege is right to execute system call or use an option within a system call
  - Can be assigned to processes
  - Users assigned roles granting access to privileges and programs

#### -- Role-based Access Control in Solaris 10



- Global Table: Simple but usually too big to be kept in memory and difficult to take advantage of special grouping of objects or domains.
- Each column = Access-control list for one object Defines who can perform what operation.

Domain 1 = Read, Write Domain 2 = Read Domain 3 = Read

 Each Row = Capability List (like a key) Fore each domain, what operations allowed on what objects. Object 1 – Read Object 4 – Read, Write, Execute Object 5 – Read, Write, Delete, Copy

- Access List Delete access rights from access list.
  - Simple
  - Immediate
- Capability List Scheme required to locate capability in the system before capability can be revoked.
  - Reacquisition
  - Back-pointers
  - Indirection
  - Keys

# End of Chapter 14