

COE 301 – Computer Organization

Assignment 2 Solution: MIPS Instructions and Assembly Language

1. (2 pts) Bits have no inherent meaning. Given the 32-bit pattern:

1010 1101 0001 0000 0000 0000 0000 0010

What does it represent, assuming it is ...

- a) A 2's complement signed integer?
- b) A MIPS instruction?

Solution:

a) -1,391,460,350

b) `Op = 1010112 = 0x2b = sw - store word (I-Type format)`

`rs = 010002 = r8 = $t0`

`rt = 100002 = r16 = $s0`

`immediate16 = 0000 0000 0000 00102 = 2`

`MIPS instruction = sw $s0, 2($t0)`

2. (2 pts) Find the shortest sequence of MIPS instructions to:

- a) Determine if there is a carry out from the addition of two registers `$t3` and `$t4`. Place the carry out (`0` or `1`) in register `$t2`. It can be done in two instructions.
- b) Determine the absolute value of a signed integer. Show the implementation of the following pseudo-instruction using three real instructions:

`abs $t1, $t2`

Solution:

a) `addu $t5, $t3, $t4`

`sltu $t2, $t5, $t3 # there is carry if sum < any operand`

b) `addu $t1, $t2, $zero`

`bgez $t2, next`

`subu $t1, $zero, $t2`

`next:`

3. (4 pts) For each pseudo-instruction in the following table, produce a minimal sequence of actual MIPS instructions to accomplish the same thing. You may use the `$at` for some of the sequences. In the following table, `imm32` refers to a 32-bit constant.

Pseudo-instruction	Solution
<code>move \$t1, \$t2</code>	<code>addu \$t1, \$t2, \$zero</code>
<code>clear \$t5</code>	<code>addu \$t5, \$zero, \$zero</code>
<code>li \$t5, imm32</code>	<code>lui \$t5, upper16</code> <code>ori \$t5, \$t5, lower16</code>
<code>addi \$t5, \$t3, imm32</code>	<code>lui \$at, upper16</code> <code>ori \$at, \$at, lower16</code> <code>add \$t5, \$t3, \$at</code>
<code>beq \$t5, imm32, Label</code>	<code>lui \$at, upper16</code> <code>ori \$at, \$at, lower16</code> <code>beq \$t5, \$at, Label</code>
<code>ble \$t5, \$t3, Label</code>	<code>slt \$at, \$t3, \$t5</code> <code>beq \$at, \$zero, Label</code>
<code>bgt \$t5, \$t3, Label</code>	<code>slt \$at, \$t3, \$t5</code> <code>bne \$at, \$zero, Label</code>
<code>bge \$t5, \$t3, Label</code>	<code>slt \$at, \$t5, \$t3</code> <code>beq \$at, \$zero, Label</code>

4. (2 pts) Translate the following statements into MIPS assembly language. Assume that *a*, *b*, *c*, and *d* are allocated in `$s0`, `$s1`, `$s2`, and `$s3`. All values are signed 32-bit integers.

a) `if ((a > b) || (b > c)) {d = 1;}`

Solution:

```

bgt $s0, $s1, L1
ble $s1, $s2, next
L1:
ori $s3, $zero, 1
next:

```

b) `if ((a <= b) && (b > c)) {d = 1;}`

Solution:

```

bgt $s0, $s1, next
ble $s1, $s2, next
ori $s3, $zero, 1
next:

```

5. (3 pts) Consider the following fragment of C code:

```
for (i=0; i<=100; i=i+1) { a[i] = b[i] + c; }
```

Assume that a and b are arrays of words and the base address of a is in \$a0 and the base address of b is in \$a1. Register \$t0 is associated with variable i and register \$s0 with c. Write the code in MIPS.

Solution:

```

        addu $t0, $zero, $zero      # i = 0
        addu $t1, $a0, $zero       # $t1 = address a[i]
        addu $t2, $a1, $zero       # $t2 = address b[i]
        addiu $t3, $zero, 101      # $t3 = 101 (max i)
loop:   lw    $t4, 0($t2)           # $t4 = b[i]
        addu $t5, $t4, $s0         # $t5 = b[i] + c
        sw    $t5, 0($t1)         # a[i] = b[i] + c
        addiu $t0, $t0, 1         # i++
        addiu $t1, $t1, 4         # address of next a[i]
        addiu $t2, $t2, 4         # address of next b[i]
        bne  $t0, $t3, loop       # exit if (i == 101)

```

6. (3 pts) Add comments to the following MIPS code and describe in one sentence what it computes. Assume that \$a0 is used for the input and initially contains n, a positive integer. Assume that \$v0 is used for the output.

```

begin:   addi $t0, $zero, 0        # $t0 = sum = 0
        addi $t1, $zero, 1        # $t1 = i = 1
loop:   slt  $t2, $a0, $t1        # (n<i)? or (i>n)?
        bne  $t2, $zero, finish   # exit loop if (i>n)
        add  $t0, $t0, $t1        # sum = sum + i
        addi $t1, $t1, 2         # i = i + 2
        j    loop                 # repeat loop
finish:  add  $v0, $t0, $zero     # result = sum

```

Result \$v0 is the sum of the odd positive integers 1 + 3 + 5 + ... which are less than or equal to n.

7. (4 pts) The following code fragment processes an array and produces two important values in registers \$v0 and \$v1. Assume that the array consists of 5000 words indexed 0 through 4999, and its base address is stored in \$a0 and its size (5000) in \$a1. Describe in one sentence what this code does. Specifically, what will be returned in \$v0 and \$v1?

```

                add  $a1, $a1, $a1      # $a1 = 5000 * 2
                add  $a1, $a1, $a1      # $a1 = 5000 * 4
                add  $v0, $zero, $zero  # $v0 = 0
                add  $t0, $zero, $zero  # $t0 = 0
outer:          add  $t4, $a0, $t0      # $t4 = address A[i]
                lw   $t4, 0($t4)        # $t4 = A[i]
                add  $t5, $zero, $zero  # $t5 = count = 0
                add  $t1, $zero, $zero  # $t1 = 0
inner:          add  $t3, $a0, $t1      # $t3 = address A[j]
                lw   $t3, 0($t3)        # $t3 = A[j]
                bne  $t3, $t4, skip     # if (A[i]!=A[j]) skip
                addi $t5, $t5, 1        # count++
skip:           addi $t1, $t1, 4        # j = j+4
                bne  $t1, $a1, inner    # inner loop = 5000
                slt  $t2, $t5, $v0      # if (count < $v0)
                bne  $t2, $zero, next   # then goto next
                add  $v0, $t5, $zero    # $v0 = count
                add  $v1, $t4, $zero    # $v1 = A[i]
next:           addi $t0, $t0, 4        # i = i+4
                bne  $t0, $a1, outer    # outer loop = 5000

```

This code compares every element in the array against all elements for identical matches. It counts the frequency of occurrence of each value in the array. The *count* of the most frequently used value is returned in \$v0 and the *value* itself is returned in \$v1.