MIPS Functions and the Runtime Stack

COE 233

Logic Design and Computer Organization

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Presentation Outline

- **Functions**
- **Function Call and Return**
- The Stack Segment
- Preserving Registers
- Examples: Bubble Sort and Recursion

Functions

- A function (or a procedure) is a block of instructions that can be called at several different points in the program
  - Allows the programmer to focus on just one task at a time
  - Allows code to be reused
- The function that initiates the call is known as the **caller**
- The function that receives the call is known as the **callee**
- When the callee finishes execution, control is transferred back to the caller function.
- A function can receive **parameters** and return **results**
- The function parameters and results act as an interface between a function and the rest of the program

Function Call and Return

- To execution a function, the **caller** does the following:
  - Puts the parameters in a place that can be accessed by the callee
  - Transfer control to the callee function
- To return from a function, the **callee** does the following:
  - Puts the results in a place that can be accessed by the caller
  - Return control to the caller, next to where the function call was made
- Registers are the fastest place to pass parameters and return results. The MIPS architecture uses the following:
  - **`$a0-$a3:`** four argument registers in which to pass parameters
  - **`$v0-$v1:`** two value registers in which to pass function results
  - **`$ra:`** return address register to return back to the caller

Function Call and Return Instructions

- **JAL** (**Jump-and-Link**) is used to call a function
  - Save return address in **$31 = PC+4** and jump to function
  - Register **$31** (**$ra**) is used by **JAL** as the **return address**
- **JR** (**Jump Register**) is used to return from a function
  - Jump to instruction whose address is in register Rs (PC = Rs)
- **JALR** (**Jump-and-Link Register**)
  - Save return address in Rd = PC+4, and
  - Call function whose address is in register Rs (PC = Rs)
  - Used to call functions whose addresses are known at runtime

Instruction	Meaning	Format					
`jal label`	`$31 = PC+4, j Label`	`Op=3`	`26-bit address`				
`jr Rs`	`PC = Rs`	`Op=0`	`Rs`	`0`	`0`	`0`	`8`
`jalr Rd, Rs`	`Rd = PC+4, PC = Rs`	`Op=0`	`Rs`	`0`	`Rd`	`0`	`9`

Example

- Consider the following `swap` function (written in C)
- Translate this function to MIPS assembly language

```c
void swap(int v[], int k)
{   int temp;
    temp = v[k]
    v[k] = v[k+1];
    v[k+1] = temp;
}
```

Parameters:

`$a0` = Address of `v[]`
`$a1` = `k`, and
Return address is in `$ra`

```
swap:
 sll $t0,$a1,2       # $t0=k*4
 add $t0,$t0,$a0     # $t0=v+k*4
 lw  $t1,0($t0)      # $t1=v[k]
 lw  $t2,4($t0)      # $t2=v[k+1]
 sw  $t2,0($t0)      # v[k]=$t2
 sw  $t1,4($t0)      # v[k+1]=$t1
 jr  $ra             # return
```

Call / Return Sequence

- Suppose we call function swap as: **swap(a,10)**
  - Pass **address** of array **a** and **10** as arguments
  - Call the function swap saving **return address** in **$31 = $ra**
  - Execute function swap
  - Return control to the point of origin (return address)

Registers

	. . .
$a0=$4	addr a
$a1=$5	10
	. . .
$ra=$31	ret addr

Caller

```
  la    $a0, a
  li    $a1, 10
  jal   swap
# return here
  . . .
```

```
swap:
 sll $t0,$a1,2
 add $t0,$t0,$a0
 lw  $t1,0($t0)
 lw  $t2,4($t0)
 sw  $t2,0($t0)
 sw  $t1,4($t0)
 jr  $ra
```

Details of JAL and JR

Address	Instructions	Assembly Language
00400020	lui $1, 0x1001	la $a0, a
00400024	ori $4, $1, 0	
00400028	ori $5, $0, 10	ori $a1,$0,10
0040002C	jal 0x10000f	jal swap
00400030	. . .	# return here
		swap:
0040003C	sll $8, $5, 2	sll $t0, $a1, 2
00400040	add $8, $8, $4	add $t0, $t0, $a0
00400044	lw $9, 0($8)	lw $t1, 0($t0)
00400048	lw $10,4($8)	lw $t2, 4($t0)
0040004C	sw $10,0($8)	sw $t2, 0($t0)
00400050	sw $9, 4($8)	sw $t1, 4($t0)
00400054	jr $31	jr $ra

Pseudo-Direct Addressing

```
PC = imm26<<2
0x10000f << 2
= 0x0040003C
```

$31 | 0x00400030

Register **$31** is the return address register

Second Example

- Function **tolower** converts a capital letter to lowercase
- If parameter **ch** is not a capital letter then return **ch**

```
char tolower(char ch) {
  if (ch>='A' && ch<='Z')
    return (ch + 'a' - 'A');
  else
    return ch;
}
```

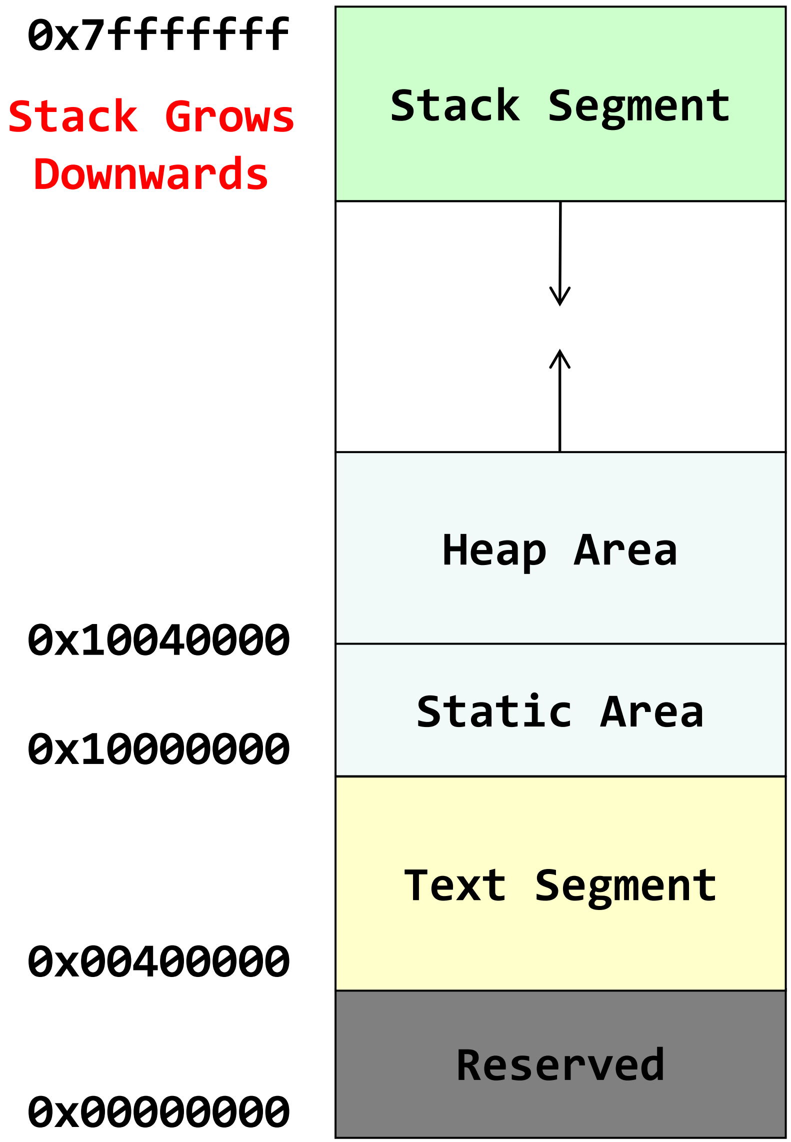
```
tolower:                            # $a0 = parameter ch
  blt   $a0, 'A', else              # branch if $a0 < 'A'
  bgt   $a0, 'Z', else              # branch if $a0 > 'Z'
  addi  $v0, $a0, 32                # 'a' – 'A' == 32
  jr    $ra                         # return to caller
else:
  move  $v0, $a0                    # $v0 = ch
  jr    $ra                         # return to caller
```

Next . . .

- Functions
- Function Call and Return
- **The Stack Segment**
- **Preserving Registers**
- Examples: Bubble Sort and Recursion

The Stack Segment

- Every program has 3 segments when loaded into memory:
  - **Text segment**: stores machine instructions
  - **Data segment**: area used for static and dynamic variables
  - **Stack segment**: area that can be allocated and freed by functions
- The program uses only logical (virtual) addresses
- The actual (physical) addresses are managed by the OS

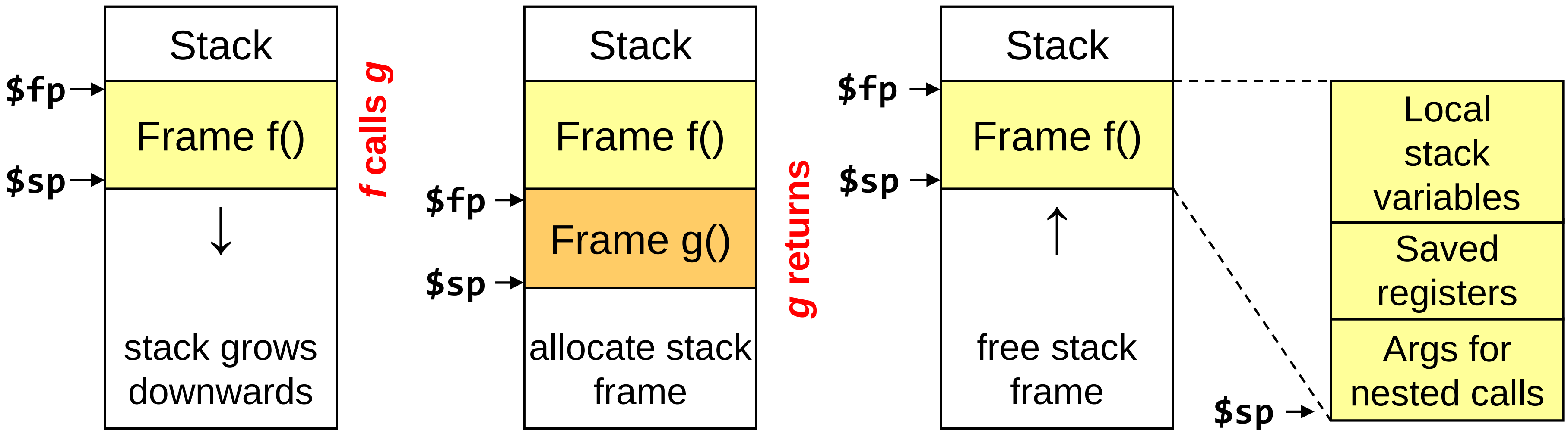


The Stack Segment (cont'd)

- The stack segment is used by functions for:
  - Passing parameters that cannot fit in registers
  - Allocating space for local variables
  - Saving registers across function calls
  - Implement recursive functions
- The stack segment is implemented via software:
  - The **`Stack Pointer $sp = $29`** (points to the top of stack)
  - The **`Frame Pointer $fp = $30`** (points to a stack frame)
- The stack pointer **`$sp`** is initialized to the base address of the stack segment, just before a program starts execution
- The MARS tool initializes register **`$sp`** to **`0x7fffeffc`**

Stack Frame

- **Stack frame** is an area of the stack containing …
  - Saved arguments, registers, local arrays and variables (if any)
- Called also the **activation frame**
- Frames are pushed and popped by adjusting …
  - Stack pointer `$sp = $29` (and sometimes frame pointer `$fp = $30`)
  - Decrement `$sp` to allocate stack frame, and increment to free



Leaf Function

- A leaf function does its work without calling any function
- Example of leaf functions are: **`swap`** and **`tolower`**
- A leaf function can freely modify some registers:
  - Argument registers: **`$a0 - $a3`**
  - Result registers: **`$v0 - $v1`**
  - Temporary registers: **`$t0 - $t9`**
  - These registers can be modified without saving their old values
- A leaf function does not need a stack frame if …
  - Its variables can fit in temporary registers
- A leaf function allocates a stack frame only if …
  - It requires additional space for its local variables

Non-Leaf Function

- A non-leaf function is a function that calls other functions
- A non-leaf function must allocate a stack frame
- Stack frame size is computed by the programmer (compiler)
- To allocate a stack frame of **N** bytes …
  - Decrement **$sp** by **N** bytes: **$sp = $sp – N**
  - **N** must be multiple of **4** bytes to have registers aligned in memory
  - In our examples, only register **$sp** will be used (**$fp** is not needed)
- Must save register **$ra** before making a function call
  - Must save **$s0-$s7** if their values are going to be modified
  - Other registers can also be preserved (if needed)
  - Additional space for local variables can be allocated (if needed)

Steps for Function Call and Return

- To make a function call …
  - Make sure that register `$ra` is saved before making a function call
  - Pass arguments in registers `$a0` thru `$a3`
  - Pass additional arguments on the stack (if needed)
  - Use the `JAL` instruction to make a function call (`JAL` modifies `$ra`)
- To return from a function …
  - Place the function results in `$v0` and `$v1` (if any)
  - Restore all registers that were saved upon function entry
    - Load the register values that were saved on the stack (if any)
  - Free the stack frame: `$sp = $sp + N` (stack frame = `N` bytes)
  - Jump to the return address: `jr $ra` (return to caller)

Preserving Registers

- The MIPS software specifies which registers must be preserved across a function call, and which ones are not

Must be Preserved	Not preserved
Return address: **$ra**	Argument registers: **$a0** to **$a3**
Stack pointer: **$sp**	Value registers: **$v0** and **$v1**
Saved registers: **$s0** to **$s7** and **$fp**	Temporary registers: **$t0** to **$t9**
Stack above the stack pointer	Stack below the stack pointer

- Caller saves register **$ra** before making a function call
- A callee function must preserve **$sp**, **$s0** to **$s7**, and **$fp**.
- If needed, the caller can save argument registers **$a0** to **$a3**. However, the callee function is free to modify them.

Example on Preserving Register

- A function **f** calls **g** twice as shown below. We don't know what **g** does, or which registers are used in **g**.
- We only know that function **g** receives two integer arguments and returns one integer result. Translate **f**:

```
int f(int a, int b) {
  int d = g(b, g(a, b));
  return a + d;
}
```

Translating Function f

```
int f(int a, int b) {
  int d = g(b, g(a, b)); return a + d;
}
```

```
f:  addiu   $sp, $sp, -12      # allocate frame = 12 bytes
    sw      $ra, 0($sp)        # save $ra
    sw      $a0, 4($sp)        # save a (caller-saved)
    sw      $a1, 8($sp)        # save b (caller-saved)
    jal     g                  # call g(a,b)
    lw      $a0, 8($sp)        # $a0 = b
    move    $a1, $v0           # $a1 = result of g(a,b)
    jal     g                  # call g(b, g(a,b))
    lw      $a0, 4($sp)        # $a0 = a
    addu    $v0, $a0, $v0      # $v0 = a + d
    lw      $ra, 0($sp)        # restore $ra
    addiu   $sp, $sp, 12       # free stack frame
    jr      $ra                # return to caller
```

Next . . .

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- The Stack Segment
- Preserving Registers
- **Examples: Bubble Sort and Recursion**

Bubble Sort (Leaf Function)

```
void bubbleSort (int A[], int n) {
  int swapped, i, temp;
  do {
    n = n-1;
    swapped = 0;                 // false
    for (i=0; i<n; i++) {
      if (A[i] > A[i+1]) {
        temp = A[i];             // swap A[i]
        A[i] = A[i+1];           // with A[i+1]
        A[i+1] = temp;
        swapped = 1;             // true
      }
    }
  } while (swapped);
}
```

Worst case Performance	$O(n^2)$
Best case Performance	$O(n)$

Translating Function Bubble Sort

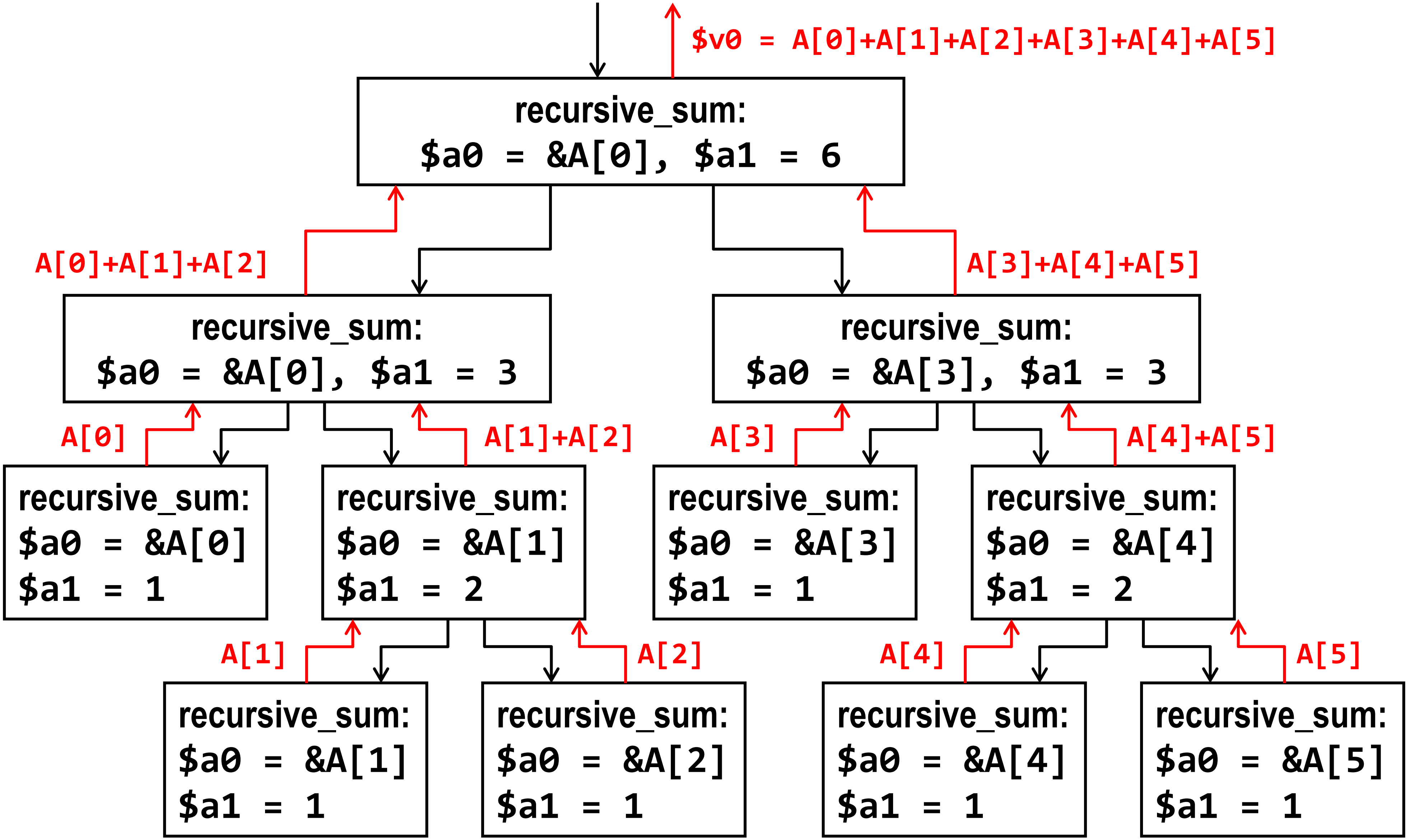
```
bubbleSort:                      # $a0 = &A, $a1 = n
do:   addiu  $a1, $a1, -1        # n = n-1
      blez   $a1, L2             # branch if (n <= 0)
      move   $t0, $a0            # $t0 = &A
      li     $t1, 0              # $t1 = swapped = 0
      li     $t2, 0              # $t2 = i = 0
for:  lw     $t3, 0($t0)         # $t3 = A[i]
      lw     $t4, 4($t0)         # $t4 = A[i+1]
      ble    $t3, $t4, L1        # branch if (A[i] <= A[i+1])
      sw     $t4, 0($t0)         # A[i] = $t4
      sw     $t3, 4($t0)         # A[i+1] = $t3
      li     $t1, 1              # swapped = 1
L1:   addiu  $t2, $t2, 1         # i++
      addiu  $t0, $t0, 4         # $t0 = &A[i]
      bne    $t2, $a1, for       # branch if (i != n)
      bnez   $t1, do             # branch if (swapped)
L2:   jr     $ra                 # return to caller
```

Example of a Recursive Function

```
int recursive_sum (int A[], int n) {
  if (n == 0) return 0;
  if (n == 1) return A[0];
  int sum1 = recursive_sum (&A[0], n/2);
  int sum2 = recursive_sum (&A[n/2], n – n/2);
  return sum1 + sum2;
}
```

- Two recursive calls
  - First call computes the sum of the first half of the array elements
  - Second call computes the sum of the 2nd half of the array elements
- How to translate a recursive function into assembly?

Illustrating Recursive Calls



Translating a Recursive Function

```
recursive_sum:                       # $a0 = &A, $a1 = n
    bnez    $a1, L1                  # branch if (n != 0)
    li      $v0, 0
    jr      $ra                      # return 0
L1: bne     $a1, 1, L2               # branch if (n != 1)
    lw      $v0, 0($a0)              # $v0 = A[0]
    jr      $ra                      # return A[0]
L2: addiu   $sp, $sp, -12            # allocate frame = 12 bytes
    sw      $ra, 0($sp)              # save $ra
    sw      $s0, 4($sp)              # save $s0
    sw      $s1, 8($sp)              # save $s1
    move    $s0, $a0                 # $s0 = &A (preserved)
    move    $s1, $a1                 # $s1 = n  (preserved)
    srl     $a1, $a1, 1              # $a1 = n/2
    jal     recursive_sum            # first recursive call
```

Translating a Recursive Function (cont'd)

```
srl     $t0, $s1, 1          # $t0 = n/2
sll     $t1, $t0, 2          # $t1 = (n/2) * 4
addu    $a0, $s0, $t1        # $a0 = &A[n/2]
subu    $a1, $s1, $t0        # $a1 = n – n/2
move    $s0, $v0             # $s0 = sum1 (preserved)
jal     recursive_sum        # second recursive call
addu    $v0, $s0, $v0        # $v0 = sum1 + sum2
lw      $ra, 0($sp)          # restore $ra
lw      $s0, 4($sp)          # restore $s0
lw      $s1, 8($sp)          # restore $s1
addiu   $sp, $sp, 12         # free stack frame
jr      $ra                  # return to caller
```

❖ **$ra**, **$s0**, and **$s1** are preserved across recursive calls