# King Fahd University of Petroleum and Minerals <br> College of Computer Sciences and Engineering <br> Department of Computer Engineering 

## SEC 521 -Network Security (T151)

## Homework \# 01 (due date \& time: Sunday 04/10/2015 during class period)

Problem \# 1: Solve problem 2.2 of the $4^{\text {th }}$ edition of William Stallings textbook.
Problem \# 2: Use the A5/1 algorithm. Suppose that, after a particular step, the values in the registers are

$$
\begin{array}{ll}
X=\left(x_{0}, x_{1}, \ldots, x_{18}\right) & =(1010101010101010110) \\
Y=\left(y_{0}, y_{1}, \ldots, y_{21}\right) & =(1100110001101100010011) \\
Z=\left(z_{0}, z_{1}, \ldots, z_{22}\right) & =(11100101110000011000011)
\end{array}
$$

List the next 4 keystream bits and give the contents of $X, Y$, and $Z$ after the generation of each of these 4 bits.

Problem \# 3: Consider a Feistel cipher with three rounds. Then the plaintext is denoted as $P=$ ( $L_{0}, R_{0}$ ) and the corresponding ciphertext is $C=\left(L_{3}, R_{3}\right)$. What is the simplest form of the ciphertext $C$, in terms of $L_{0}, R_{0}$, and the subkey, for each of the following round functions?
a. $\quad F\left(R_{i-1}, K_{i}\right)=\overline{R_{i-1}}$, where $\overline{R_{i-1}}$ is the logical complement of $R_{i-1}$
b. $\quad F\left(R_{i-1}, K_{i}\right)=R_{i-1} \oplus K_{i}$

Problem \# 4: Solve problem 2.16 (only part b) of the $4^{\text {th }}$ edition of William Stallings textbook.
Problem \# 5: Use the "Repeated Squaring" method on p. 104 of the "Public-Key Cryptography" slides to compute $9^{25} \bmod 15$. Show the power groupings and the steps.

Problem \# 6: Solve problem 3.14 (parts d and e) of the $4^{\text {th }}$ edition of William Stallings textbook.
Problem \# 7: Solve problem 3.21 of the $4^{\text {th }}$ edition of William Stallings textbook.
Problem \# 8: Suppose that Bob uses the following variant of RSA. He first chooses $N$, then he finds two encryption exponents, $e_{0}$ and $e_{1}$, and the corresponding decryption exponents $d_{0}$ and $d_{1}$. He asks Alice to encrypt her message $M$ to him by first computing $C_{0}=M^{e 0} \bmod N$, then encrypting $C_{0}$ to obtain the ciphertext, $C_{1}=C_{0}{ }^{e 1}$ $\bmod N$. Alice then sends $C_{1}$ to Bob. Does this double encryption increase the security as compared to a single RSA encryption? Why or why not?

