## King Fahd University of Petroleum \& Minerals Computer Engineering Dept

COE 402 - Computer Systems Performance Evaluation
Term 043
Dr. Ashraf S. Hasan Mahmoud
Rm 22-148-3
Ext. 1724
Email: ashraf@ccse.kfupm.edu.sa

## Queuing Model

- Consider the following system:

$$
\mathrm{A}(\mathrm{t}) \quad \mathrm{n}(\mathrm{t})=\mathrm{A}(\mathrm{t})-\mathrm{D}(\mathrm{t}) \quad \mathrm{D}(\mathrm{t})
$$

| ith customer arrives at time $\mathrm{S}_{\mathrm{i}}$ | Queueing System | ith customer departs at time $\mathrm{D}_{\mathrm{i}}$ |
| :---: | :---: | :---: |
|  | $\mathrm{r}_{\mathrm{i}}=\mathrm{D}_{\mathrm{i}}-\mathrm{A}_{\text {i }}$ | $\mathrm{w}_{\mathrm{i}}=\mathrm{r}_{\mathrm{i}}-\mathrm{S}_{\mathrm{i}}$ |
| $\mathrm{A}(\mathrm{t})$ - number of arrivals in (0,t] |  | $=D_{i}-A_{i}-S_{i}$ |
| $\mathrm{D}(\mathrm{t})$ - number of departures in $(0, \mathrm{t}]$ |  |  |
| $\mathrm{n}(\mathrm{t})$ - number of customers in system in (0,t] |  |  |
| $\mathrm{r}_{\mathrm{i}}-$ duration of time spent in system for $\mathrm{r}^{\text {ith }}$ customer - textbook call this response time $(r)$$\mathrm{w}_{\mathrm{i}}-$ duration of time spent waiting for service for $\mathrm{i}^{\text {th }}$ customer |  |  |
|  |  |  |

## Queuing Model

- The model is characterized using the following quantities:
$\lambda$ - mean arrival rate of customers $=1 / E[\tau]$ (remember $\tau$ is the interarrival time)
- $s=$ service time per job
- $\quad \mu=$ mean service rate $=1 / E[s]$
- $n=$ number of job in system $\rightarrow E[n]$
- $n_{q}=$ number of jobs in buffer $\rightarrow E\left[n_{q}\right]$
- $\quad \mathrm{ns}=$ number of jobs in server $\rightarrow \mathrm{E}[\mathrm{ns}]$
- $\quad r=$ response time or total time in system $\rightarrow \mathrm{E}[r]$
- $w=$ waiting time $\rightarrow E[w]$



## Example: Queueing System

Problem: A data communication line delivers a block of information every 10 microseconds. A decoder check each block for errors and corrects the errors if necessary. It takes 1 microsecond to determine whether the block has any errors. If the block has one error it takes 5 microseconds to correct it and it has more than 1 error it takes $\mathbf{2 0}$ microseconds to correct the error. Blocks wait in the queue when the decoder falls behind. Suppose that the decoder is initially empty and that the number of errors in the first 10 blocks are: $0,1,3,1,0,4,0,1,0,0$.
a) Plot the number of blocks in the decoder as a function of time.
b) Find the mean number of blocks in the decoder
c) What percent of the time is the decoder empty?

## Example: Queueing System - cont'd

## Solution:

Interarrival time $=10 \boldsymbol{\mu s e c}$
Service time $=1 \quad$ if no errors
$1+5$ if 1 error
$1+20$ if more than 1 error
The queue parameters ( $A, D, S$, and $W$ ) are shown below:

| Block \#: | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Arrivals: | 10 | 20 | 30 | 40 | 50 | 60 | 70 | 80 | 90 | 100 |
| Errors: | 0 | 1 | 3 | 1 | 0 | 4 | 0 | 1 | 0 | 0 |
| Service: | 1 | 6 | 21 | 6 | 1 | 21 | 1 | 6 | 1 | 1 |
| Departs: | 11 | 26 | 51 | 57 | 58 | 81 | 82 | 88 | 91 | 101 |
| Waiting: | 0 | 0 | 0 | 11 | 7 | 0 | 11 | 2 | 0 | 0 |

## Example: Queueing System - cont'd

Solution:
Using the previous results and knowing that

$$
n(t)=A(t)-D(t)
$$

One can produce the following results

Average no of customers in system $=0.950$
Average customer waiting time $=3.100 \mathrm{microsec}$
Maximum simulation time $\quad=101.000$ microsec
Duration server busy = 65.000
Server utilization $=0.6436$
Server idle
$=0.3564$

The following Matlab code can be used to solve this queue system (Note the code is general - it solves any system provided The Arrivals vector $A$, and the service vector $\mathbf{S}$ )


## Example: Queueing System - cont'd

| 0001 \% | $0033 \%$ Compute $\mathrm{N}(\mathrm{t})$ |
| :---: | :---: |
| 0002 \% Problem 9.3-Leon Garcia's book | $0034 \mathrm{~T}=[] ; \%$ time axis |
| 0093 clear all | ${ }^{0035} \mathbf{T}(1)=0 ; \%$ time origin |
| $0094 \mathrm{~A}=$ [10:10:100]; | ${ }_{0}^{0036} \mathbf{N}$ (1) $=[] ;$ \% number of cutomers |
|  | $0037 \mathrm{~N}(1)=0$; initial condition |
| $0006 \mathrm{~S}=$ zeros(size(A)); | $0038 \mathrm{k}=2 ; \%$ place for next insert |
| $0007 \mathrm{D}=$ zeros(size(A)); | 0039 A_max $=$ A(length(A)); \% last arrival instant |
|  | $0040 \mathrm{i}=1 ; \quad$ \% index for arrivals |
| $0009 \%$ this loop to computes service times | $0041 \mathrm{j}=1$ 1; \% index for departures |
| 0010 for $\mathrm{i}=1$ : length( A$)$; | $0042 \mathrm{t}=0$; \% system time |
| 0011 if (Errors(i)==0) $\mathrm{S}(\mathrm{i})=1$; | ${ }^{0043}$ |
| 0012 else | 0044 while ( t < A max) |
| 0013 if (Errors(i)==1) S(i) = 6; | $0045 \quad \mathrm{t}=\min (\mathrm{A}(\mathrm{i}), \mathrm{D}(\mathrm{j})$ ); |
| 0014 else | 0046 if ( $\mathrm{t}=\mathrm{=A}(\mathrm{i})$ ) |
| 0015 S(i) $=21$; | $0047 \quad \mathrm{~N}(\mathrm{k})=\mathrm{N}(\mathrm{k}-1)+1$; |
| 0016 end | 0048 T(k) = t; |
| 0017 end | ${ }^{0049} \mathrm{k}=\mathrm{k}+1$; |
| 0018 \% | ${ }^{0050}$ i ${ }^{\text {i }}$ = $\mathbf{i}+1$; \% get next arrival |
| 0019 \% this section computes the departure time for | 0051 else \% departure occurs |
| the ith user | $0052 \mathrm{~N}(\mathrm{k})=\mathrm{N}(\mathrm{k}-1)-1$; |
| 0020 if (i>1) \% this is not the first user |  |
| 0021 if ( $D(i-1)<A(i)) D(i)=A(i)+S(i) ;$ |  |
| 0022 else | 0055 j $\quad$ j ${ }^{0}+1 ; \%$ get next departure |
| ${ }^{0023}$ 203 $D(i)=D(i-1)+S(i) ;$ | ${ }^{0056}$ end |
| ${ }^{0024}$ end end | 0057 end |
| 0025 else | 0058 \% |
| ${ }^{0026} \mathrm{D}(\mathrm{i})=\mathrm{A}(\mathrm{i})+\mathrm{S}(\mathrm{i})$; | 0059 \% record remaining departure instants |
| 0027 end | 0060 for i=j:1:length( D ) |
| 0028 \% | $0061 \mathrm{t}^{\text {d }}$ (k) $\mathrm{D}(\mathrm{i})$; |
| 0029 \% compute waiting time | ${ }^{0062} \quad N(k)=N(k-1)-1 ;$ |
| ${ }_{\text {®031 }}^{\text {日030 }}$ end $\mathrm{W}(\mathrm{i})=\mathrm{D}(\mathrm{i})-\mathrm{A}(\mathrm{i})-\mathrm{S}(\mathrm{i})$; |  |
| ${ }^{0031}$ end ${ }^{\text {ens }}$ (i) |  |
|  | $\begin{aligned} & 0065 \text { end } \\ & { }_{0066} \end{aligned}$ |
|  | $0067 \mathrm{k}=\mathrm{k}-1$; \% decrement k to get real size of N and T 0068 \% |
|  | 0069 \% compute means |
|  | 0070 MeanW $=\operatorname{mean}(W) ;$ <br> 0071 T_Intervales $=T(2: k)-T(1: k-1) ;$ |
|  |  |
| 8/20/2005 Dr. Ashraf S. | 0073 IdleDurationsIndex $=$ find $(\mathbb{N}(1: \mathrm{k}-1)=0$ ) ; |
|  |  |

## Example: Queueing System - cont'd

```
0076% Display results
    0077 fprintf('Block #: '); fprintf('%3d ', [1:1:length(A)]); fprintf('\n');
    0078 fprintf('Arrivals
    0079 fprintf('Errors:
    fprintf('service
    fprintf('Departs
    0 0 8 2 \text { fprintf('Waiting}
    * 0084 fprintf('Average no of customers in system = %7.3f\n', MeanN);
    0086 fprintf('Maximum simulation time 
    0088 sum(T Intervales(IdleDurationsIndex)));
```



```
    lol}\begin{array}{ll}{0089 fprintf('Server utilization }&{=%7.4f\n', Utilization);}\\{009 fprintf('Server idle }&{=%7.4f\n',1.0-Utilization}
    0092 % Plot results
    0094 h= stairs(T,N); gri
    0095 set(h, 'LineWidth',
    0097 ylabel('No of customers in system, N(t)');
    0098
    0099 figure(2)
    0100 [AT, AA] = stairs(A, cumsum(ones(size(A))))
    0101 [DT, DD] = stairs(D, cumsum(ones(size(D))))
```



```
    0104 set(h, 'LineWidth', '3);
    0106 ylabel('No of customers')
    0107 xlabel('Time');
    0108 legend('A(t)',''D(t)', 'N(t)', 0);
    0119
    0111 h = stem(w); grid
    0112 set(h, 'LineWidth', 3);
    0113 ylabel('Waiting time');
    0114 xlabel('Customer index');;
- Blue curve: A(t)
- Red curve: D(t)
- Total time spent in the system for all customers = area in between two curves

\section*{Little's Formula}
- Little's formula:
\[
E[n]=\lambda E[r]
\]

Holds for many service disciplines and for systems with arbitrary number of servers. It holds for many interpretations of the system as well

\section*{Example 1:}
- Problem: Let \(\mathbf{n}_{\mathbf{s}}(\mathrm{t})\) be the number of customers being served at time \(t\), and let \(\tau\) denote the service time. If we designate the set of servers to be the "system" then Little's formula becomes:
\[
E\left[n_{s}\right]=\lambda E[s]
\]

Where \(E\left[n_{s}\right]\) is the average number of busy servers for a system in the steady state.

\section*{Example 1: cont'd}

Note: for a single server \(n_{s}(t)\) can be either 0 or \(1 \rightarrow E\left[n_{s}\right]\) represents the portion of time the server is busy. If \(\mathbf{p}_{0}=\) \(\operatorname{Prob}\left[n_{s}(t)=0\right]\), then we have
\[
\begin{aligned}
1-p_{0} & =E\left[n_{s}\right]=\lambda E[s], O r \\
p_{0} & =1-\lambda E[s]
\end{aligned}
\]

The quantity \(\lambda E[s]\) is defined as the utilization, \(U\), for a single server. Usually, it is given the symbol \(\rho\)
\[
\rho=\lambda E[s]
\]

For a c-server system, we define the utilization (the fraction of busy servers) to be
\[
\rho=\lambda E[s] / c
\]

\section*{The M/M/1 Queue}
- Consider a single server system where customers arrive according to a Poisson process of rate \(\boldsymbol{\lambda}\)
- \(\quad \rightarrow\) inter-arrival times are iid exponential r.v. with mean \(\mathrm{E}[\tau]=1 / \lambda\)
- Assume the service times are iid exponential r.v. with mean \(E[s]=1 / \mu\)
- Assume the inter-arrival times and service times are independent
- Assume the system can accommodate unlimited number of customers

\section*{The M/M/1 Queue - cont'd}
- What is the steady state pmf of \(n(t)\), the number of customers in the system?
- What is the PDF of \(\boldsymbol{r}\), the total customer delay in the system or the response time (as used in the textbook)?

\section*{The M/M/1 Queue - cont'd}
- Consider the transition rate diagram for M/M/1 system

- Note:
- System state - number of customers in systems
- \(\boldsymbol{\lambda}\) is rate of customer arrivals
- \(\mu\) is rate of customer departure

\section*{The M/M/1 Queue - Distribution of Number of Customers}
- Writing the global balance equations for this Markov chain and solving for
\(\operatorname{Prob}[n(t)=j\) ], yields (refer to previous example)
\[
\begin{aligned}
p_{j} & =\operatorname{Prob}[n(t)=j] \\
& =(1-\rho) \rho^{j} \text { for } j=0,1,2, \ldots,
\end{aligned}
\]
for \(\rho=\lambda / \mu<1\)

Note that for \(\rho=1 \rightarrow\) arrival rate \(\boldsymbol{\lambda}=\) service rate \(\mu\)

\section*{The M/M/1 Queue - Expected Number of Customers}
- The mean number of customers (in the whole system = buffer + server) is given by
\[
\begin{aligned}
E[n] & =\sum_{j} j \operatorname{Prob}[N(t)=j] \\
& =\rho /(1-\rho) \quad \text { for } \rho<1
\end{aligned}
\]
- You can also show that the variance is equal to
\[
\begin{aligned}
\operatorname{Var}[n] & =\sum_{j}(j-E[n])^{2} \operatorname{Prob}[n(t)=j] \\
& =\rho /(1-\rho)^{2} \quad \text { for } \rho<1
\end{aligned}
\]

\section*{The M/M/1 Queue - Mean Customer Delay}
- The mean total response time for the system, \(E[r]\), is found using Little's formula
\[
\begin{aligned}
E[r] & =E[n] / \lambda \\
& =(1 / \mu) /(1-\rho) \\
& =1 /(\mu-\lambda)
\end{aligned}
\]
- Actually, the response time CDF can be shown to be
\[
F(r)=1-e^{-r \mu(1-\rho)} \quad r \geq 0
\]
- Can you fine \(\mathrm{F}^{-1}(\mathbf{r})\) for this CDF?

\section*{The M/M/1 Queue - Mean Queueing Time}
- The mean waiting time in queue, \(\mathrm{E}[\mathrm{w}]\), is given by
\[
\begin{aligned}
E[w] & =E[r]-E[s] \\
& =\rho /(1-\rho) E[s]
\end{aligned}
\]
- The CDF for the waiting time can be shown to be
\[
F(w)=1-\rho e^{-w \mu(1-\rho)} \quad w \geq 0
\]
- Can you fine \(\mathrm{F}^{-1}(\mathbf{w})\) for this CDF?

\section*{The M/M/1 Queue - Mean Number in \\ Queue}
- Again we employ Little's formula:
\[
\begin{aligned}
E\left[n_{q}\right] & =\lambda E[w] \\
& =\rho^{2} /(1-\rho)
\end{aligned}
\]

Another way of finding \(\mathrm{E}\left[n_{q}\right]\) is using
\(E\left[n_{q}\right]=\sum_{n=1}^{\infty}(n-1) p_{n}=\sum_{n=1}^{\infty}(n-1)(1-\rho) \rho^{n}\)

Remember:
\[
\text { server utilization } \rho=\lambda / \mu=1-p_{0}
\]

All previous quantities \(E[n], E[r], E[w]\), and
\[
\mathrm{E}\left[\mathbf{n}_{\mathbf{q}}\right] \rightarrow \infty \text { as } \rho \rightarrow \mathbf{1}
\]

\section*{Busy Period}
- Def: The time interval between two successive idle intervals
\[
1
\]
- Mean busy period = \(\qquad\)
\(\mu(\mathbf{1 - p})\)
- The textbook provides lots of other formulas in regard to the busy period.

\section*{Scaling Effect for M/M/1 Queues}
- Consider a queue of arrival rate \(\lambda\) whose service rate is \(\mu\)
- \(\rho=\lambda / \mu\),
- The expected delay \(\mathrm{E}[\mathrm{r}]\) is given by
\[
E[r]=(1 / \mu) /(1-\rho)
\]
- If the arrival rate increases by a factor of
\(K\), then we either
1. Have \(K\) queueing systems, each with a server of rate \(\mu\)
2. Have one queueing system with a server of rate \(\mathrm{K} \mu\)
- Which of the two options will perform 820/2better?

\section*{Scaling Effect for M/M/1 Queues cont'd}
- Example: K = 2: M/M/1 and M/M/2 systems with the same arrival rate and the same maximum processing rate

Case 2


Case 1


\section*{Scaling Effect for M/M/1 Queues cont'd}
- Case 1: \(K\) queueing systems
- Identical systems
- \(E[r]\) is the same for all \(-E[r]=(1 / \mu) /(1-\rho)\)
- Case 2: 1 queueing system with server of rate K \(\mu\)
- \(\rho\) for this system \(=(K \lambda) /(K \mu)=\lambda / \mu\) - same as the original system
- \(\quad E\left[r^{\prime}\right]=(1 /(K \mu)) /(1-\rho)=(1 / K) E[r]\)
- Therefore, the second option will provide a less total delay figure - significant delay performance improvement!

\section*{Example 31.1}
- Problem: On a network gateway, measurements show that the packets arrive ata mean rate of \(\mathbf{1 2 5}\) packets per second (pps) and the gateway take about 2 milliseconds to forward them. Using an M/M/1 model:
(A) Analyze the gateway.
(B) What is the probability of buffer overflow had only 13 buffers?
(C) How many buffers do we need to keep packet loss below one packet per million?

\section*{Example 31.1 - cont'd}
- Solution:
A) Analyzing the gateway:

Arrival rate, \(\lambda=125\) pps
Service rate, \(\mu=1 / 0.002=500 \mathrm{pps}\)
\(\rightarrow\) Gateway utilization, \(\rho=\lambda / \mu=0.25\)
\(\operatorname{Prob}\left[n\right.\) packets in gateway] \(=p_{n}=(1-\rho) \rho^{n}\) \(=(0.75)(0.25)^{n}\)

\section*{Note- \\ (1) the solution in the textbook \\ has typos in parts (B) and (C) \\ (2) The solution for parts (B) and (C) \\ are approximate - the more accurate model should be M/M/1/B - see the solution for this example after the \\ M/M/1/B material}

Mean \# of packets in gateway, \(\mathrm{E}[\mathrm{n}]=\rho /(1-\rho)\)
\[
=0.25 / 0.75=0.33
\]

Mean time in gateway, \(\mathrm{E}[\mathrm{r}]=1 /(\mu-\lambda)=1 /(500-125)=2.66 \mathrm{msec}\)
B) Prob[buffer overflow] = Prob[more than 13 packets in gateway]
\[
\begin{aligned}
& =p_{14}+p_{15}+p_{16}+\ldots \\
& =(1-\rho)\left\{\rho^{14}+\rho^{15}+\rho^{16}+\ldots\right\} \\
& =(1-\rho) \rho^{14}\left\{1+\rho+\rho^{2}+\ldots\right\} \\
& =\rho^{14}=3.7 \times 10^{-9}
\end{aligned}
\]
C) To limit the probability of loss to less than \(10^{-6} \rightarrow\) Using (B)

Prob[buffer overflow] \(\leq \mathbf{1 0}^{-6}\)
\(\rightarrow\) Prob[more than \(n\) packets in gateway] \(\leq 10^{-6}\)
\(\rightarrow \mathrm{p}^{\mathrm{n+1}} \leq 10^{-6} \boldsymbol{\rightarrow} \mathrm{n} \leq \log \left(10^{-6}\right) / \log (\rho)-1=9.97-1 \rightarrow\) choose \(\mathrm{n}=9\) buffers

\section*{M/M/1/B - Finite Capacity Queue}
- Consider an M/M/1 with finite capacity B < \(\infty\)
- For this queue - there can be at most B customers in the system
- 1 being served
- B-1 waiting
- A customer arriving while the system has B customers is BLOCKED (does not wait)!

\section*{M/M/1/B - Finite Capacity Queue cont'd}
- Transition rate diagram for this queueing system is given by:
- \(\mathrm{n}(\mathrm{t})\) - A continuous-time Markov chain which takes on the values from the set \(\{0,1, \ldots, B\}\)


\section*{M/M/1/B - Finite Capacity Queue -}

\section*{cont'd}
- The global balance equations:
\(\lambda \quad p_{0}=\mu p_{1}\)
\((\lambda+\mu) \mathbf{p}_{j}=\lambda \mathbf{p}_{j-1}+\mu \mathbf{p}_{j+1} \quad\) for \(\mathbf{j}=1,2, \ldots, B-1\)
\(\mu \quad p_{B}=\lambda p_{B-1}\)
\(\rightarrow \operatorname{Prob}[n(t)=j]=p_{j} \quad j=0,1, \ldots, B ; \rho<1\) \(=(1-\rho) \rho^{j} /\left(1-\rho^{B+1}\right)\)
When \(\rho=1, p_{j}=1 /(B+1)\) (all states are equiprobable)

> Will this system become
> unstable for \(\rho=1\) ? Why?
- Two important numbers: \(\mathbf{p}_{\mathbf{0}}\) and \(\mathbf{p}_{\mathrm{B}}\)
- \(p_{0}\) is the probability of the server being idle \(-p_{0}=(1-\rho) /(1-\) \(\mathrm{P}^{\mathrm{B+1}}\) )
- \(\quad p_{B}\) is the probability of an arrival being blocked (or system overflow \()-p_{B}=(1-\rho) \rho^{B} /\left(1-\rho^{B+1}\right)\)

\section*{M/M/1/B - Mean Number of}

\section*{Customers}
- Mean number of customers, \(\mathrm{E}[\mathrm{n}]\) is given by:
\[
\begin{aligned}
E[n] & =\sum_{j=0}^{B} j \operatorname{Pr}[n(t)=j] \\
& = \begin{cases}\frac{\rho}{1-\rho}-\frac{(B+1) \rho^{B+1}}{1-\rho^{B+1}} & \rho<1 \\
B / 2 & \rho=1\end{cases}
\end{aligned}
\]

\section*{M/M/1/B - Blocking Rate}
- A customer arriving while the system is in state K is BLOCKED (does not wait)!
- Therefore, rate of blocking, \(\lambda_{b}\) is given by
\[
\lambda_{b}=\lambda p_{b}
\]
- The actual arrival rate into the system is \(\lambda_{a}\) given

\section*{M/M/1/B - Blocking Rate - cont'd}


\section*{M/M/1/B - Mean Response Time}
- The mean total response time, \(E[r]\) is given by
\[
E[r]=E[n] / \lambda_{a}
\]

\section*{Example 31.1 - Redoing Part (B) and (C)}
- Problem: On a network gateway, measurements show that the packets arrive ata mean rate of \(\mathbf{1 2 5}\) packets per second (pps) and the gateway take about 2 milliseconds to forward them.
(A) Analyze the gateway.
(B) What is the probability of buffer overflow had only 13 buffers?
(C) How many buffers do we need to keep packet loss below one packet per million?

\section*{Example 31.1 - cont'd}
- Solution:
B) \(\operatorname{Prob}\left[\right.\) buffer overflow] \(=\mathbf{p}_{13}\)
\[
\begin{aligned}
& =(1-\rho) \rho^{13} /\left(1-\rho^{13+1}\right) \\
& =1.1 \times 10^{-8}
\end{aligned}
\]
C) To limit the probability of loss to less than \(10^{-6} \rightarrow\) Using (B) Prob[buffer overflow] \(\leq \mathbf{1 0}^{-6}\)
\(\rightarrow \mathrm{p}_{\mathrm{B}} \leq 10^{-6}\)
\(\rightarrow(1-\rho) \rho^{B} /\left(1-\rho^{B+1}\right) \leq 10^{-6}\)
\(\rightarrow \rho^{B} /\left(1-\rho^{B+1}\right) \leq 10^{-6} /(1-\rho)=1.33 \times 10^{-6}\)
The above can be solved numerically for the value of \(B\)

\section*{Multi-Server Systems: M/M/m}
- The transition rate diagram for a multiserver M/M/m queue is as follows:
- Departure rate \(=\mathbf{k} \mu\) when \(k\) servers are busy


\((\mathrm{m}-1) \mu \quad \mathrm{m} \mu\)

\(\mathrm{m} \mu \quad \mathrm{m} \mu \quad \mathrm{m} \mu\)

\section*{Multi-Server Systems: M/M/m cont'd}
- Writing the global balance equations:
\(\lambda \quad p_{0}=\mu p_{1}\)
\(\mathrm{j} \mu \quad \mathrm{p}_{\mathrm{j}}=\boldsymbol{\lambda} \mathrm{p}_{\mathrm{j}-1}\) for \(\mathrm{j}=1,2, \ldots, \mathrm{~m}\)
\(m \mu p_{j}=\lambda p_{j-1}\) for \(j=m, m+1, \ldots\)
\(\rightarrow\)
\[
\begin{array}{ll}
p_{j}=a^{j} / j!p_{0} & (\text { for } j=1,2, \ldots, m-1) \text { and } \\
p_{j}=\rho^{j-m} / m!a^{m} p_{0} & (\text { for } j=m, m+1, \ldots)
\end{array}
\]
where \(a=\lambda / \mu\) and \(\rho=a / m\)

\section*{Multi-Server Systems: M/M/m cont'd}
- To find \(p_{0}\), we resort to the fact that \(\Sigma p_{j}=1\)

The probability that an arriving customer has to wait
\(\operatorname{Prob}[W>0]=\operatorname{Prob}[\mathrm{N} \geq \mathrm{m}]\)

In the textbook, this
quantity is denoted
by
8/20/2005
\[
\begin{aligned}
& =p_{m}+p_{m+1}+p_{m+2}+\ldots \\
& =p_{m} /(1-\rho) \\
& =p_{0} a^{m} /\{m!(1-\rho)\}
\end{aligned}
\]

Multi-Server Systems: M/M/m cont'd
- The mean number of customers in queue (waiting):
\[
\begin{aligned}
E\left[n_{q}\right] & =\sum_{j=m}^{\infty}(j-m) \operatorname{Pr}[n(t)=j] \\
& =\sum_{j=m}^{\infty}(j-m) \rho^{j-m} p_{m} \\
& =\frac{\rho}{(1-\rho)^{2}} p_{m} \\
& =\frac{\rho}{1-\rho} \operatorname{Pr}[W>0]
\end{aligned}
\]

\section*{Multi-Server Systems: M/M/m cont'd}
- The mean waiting time in queue:
\[
E[w]=E\left[n_{q}\right] / \lambda
\]
- The mean total response time in system:
\[
\begin{aligned}
E[r] & =E[w]+E[s] \\
& =E[w]+1 / \mu
\end{aligned}
\]
- The mean number of customers in system:
\[
\begin{aligned}
E[n] & =\lambda E[r] \\
& =E\left[n_{q}\right]+a
\end{aligned}
\]

Why?
- In addition to the formulas above the textbook gives the corresponding formulas for the variances too.

\section*{Example 2:}
- A company has a system with four private telephone lines connecting two of its sites. Suppose that requests for these lines arrive according to a Poisson process at rate of one call every 2 minutes, and suppose that call durations are exponentially distributed with mean 4 minutes. When all lines are busy, the system delays (i.e. queues) call requests until a line becomes available. Find the probability of having to wait for a line.

\section*{Example 2: cont'd}
- Solution:
\[
\lambda=1 / 2,1 / \mu=4, m=4 \rightarrow a=\lambda / \mu=2
\]
\[
\rightarrow \rho=a / m=1 / 2
\]
\(\begin{aligned} p_{0} & =\left\{1+2+2^{2} / 2!+2^{3} / 3!+2^{4} / 4!(1 /(1-\rho))\right\}^{-1} \\ & =3 / 23\end{aligned}\)
\(=3 / 23\)
\(p_{m}=a^{m} / m!p_{0}\)
\(=2^{4} / 4!\times 3 / 23\)

Prob[W > 0] \(=p_{m} /(1-\rho)\)
\(=2^{4} / 4!\times 3 / 23 \times 1 /(1-1 / 2)\)
\(=4 / 23\)
\(\approx 0.17\)

\section*{Waiting/Response Time Distribution for M/M/m}
- See textbook
\[
F(w)=1-\frac{p_{c}}{1-\rho} e^{-m \mu(1-\rho) w} \quad w>0
\]
- The q-percentile can be computed as follows:
\[
\begin{aligned}
w_{q} & =\max \left\{0, \frac{1}{m \mu(1-\rho)} \ln \left(\frac{100 * \operatorname{Pr}[W>0]}{100-q}\right)\right\} \\
& =\max \left\{0, \frac{E[w]}{\substack{\operatorname{Pr}[W>0] \\
\text { Dr. Ashraf s. Hasan Mahmoud }}} \ln \left(\frac{100 * \operatorname{Pr}[W>0]}{100-q}\right)\right\}
\end{aligned}
\]

\section*{Example 13.2:}
- Problem: Students arrive at the university computer center in a Poisson manner at an average rate of 10 per hour. Each student spends an average of \(\mathbf{2 0}\) minutes at the terminal, and the time can be assumed to be exponentially distributed. The center currently has five terminals. Some students have been complaining that waiting times are too long.
- Analyze the center using a queueing model.

\section*{\(p_{0}=\left\{\sum_{j=0}^{m-1} \frac{a^{j}}{j!}+\frac{a^{m}}{m!} \frac{1}{1-\rho}\right\}^{-1}\)}
- Solution:

The center model: M/M/5
Arrival rate, \(\lambda=1 / 6\) student/min
Service rate, \(\mu=1 / 20\) student/min
\(\rightarrow\) Center utilization: \(a=\lambda / \mu=3.3333\)
\[
\rho=a / m=3.3333 / 5=0.6667 \leftarrow a v g \text { server }
\]
utilization
\(p_{0}=\left\{1+10 / 3+(10 / 3)^{\wedge} 2 / 2!+(10 / 3)^{\wedge} 3 / 3!+\right.\)
\(\left.(10 / 3)^{\wedge} 4 / 4!+(10 / 3)^{\wedge} 5 / 5!/(1-2 / 3)\right\}^{-1}\)
\(=\{31.4938\}^{-1}=0.0318\)
\(p_{m}=\mathbf{a}^{m} / \mathrm{m}!\mathrm{p}_{0}\)
\(=(10 / 3)^{\wedge} 5 / 5!* 0.0318\)
\(=0.1091\)
\(\operatorname{Prob}[W>0]=p_{m} /(1-\rho)\)
\[
=0.1091 /(1-2 / 3)
\]
\(=0.3271 \leftarrow\) This is the probability that a student has to
wait!

\section*{Example 13.2: cont'd}
\[
E\left[n_{q}\right]=\frac{\rho}{1-\rho} \operatorname{Pr}[W>0]
\]
- Solution:

Avg \# of students waiting in the center, \(E\left[n_{q}\right]\)
\(E\left[n_{q}\right]=\rho /(1-\rho) \operatorname{Prob}[W>0]\)
\[
=(2 / 3) /(1-2 / 3) * 0.3271
\]
\[
=0.6543 \text { students }
\]

Avg waiting time for a student, \(\mathrm{E}[\mathrm{w}]\)
\(E[w]=E\left[n_{q}\right] / \lambda\)
\[
=0.6543 /(1 / 6)
\]
\[
=3.9259 \approx 4 \text { minutes }
\]

Avg time spent in the centre, \(\mathrm{E}[\mathrm{r}]\)
\(E[r]=E[w]+E[s]\)
\(=4+1 /(1 / 20)\)
\(=24\) minutes
Avg/variance \# of students in the center, \(\mathrm{E}[\mathrm{n}]\) \(E[n]=\lambda E[r]\)
\[
=(1 / 6) * 24=4 \text { students }
\]
\(\operatorname{Var}[n]=\ldots=479\) students \(^{2}\)

\section*{}
- Solution:

Avg \# of students using the terminals, \(\mathrm{E}\left[\mathrm{n}_{\mathrm{s}}\right]\)
\[
\begin{aligned}
E\left[n_{s}\right] & =E[n]-E[n q] \\
& =4-0.6543 \\
& =3.35 \text { students }
\end{aligned}
\]

The 90 -percentile of the waiting time, \(\mathbf{w}_{\mathbf{0 . 9}}\)
w0.9 = \(\max \{0,4 / 0.3271 \ln (10 * 0.3271)\}\)
= 14 minutes
THUS, only \(10 \%\) of the students have to wait more than 14 minutes!!

\section*{Multi-Server Systems: M/M/m/m}
- The transition rate diagram for a multiserver with no waiting room (M/M/m/m) queue is as follows:
- Departure rate \(=\mathbf{k} \mu\) when \(\mathbf{k}\) servers are busy

(m-1) \(\mu \quad \mathrm{m} \mu\)

\section*{PMF for Number of Customers for M/M/m/m}
- Writing the global balance equations, one can show:
\[
p_{j}=a^{j} / j!p_{0} \quad(\text { for } j=0,1, \ldots, m)
\]
where \(\mathbf{a}=\lambda / \mu\) (the offered load)
- To find \(p_{0}\), we resort to the fact that \(\boldsymbol{\Sigma} p_{j}\) \(=1\)
\[
p_{0}=\left\{\sum_{j=0}^{m} \frac{a^{j}}{j!}\right\}^{-1}
\]

\section*{Erlang-B Formula}
- Erlang-B formula is defined as the probability that all servers are busy:
\[
\begin{aligned}
\operatorname{Pr}[n=m]= & p_{m} \\
& =\frac{a_{m} / m!}{1+a+a^{2} / 2!+\ldots+a^{m} / m!}
\end{aligned}
\]

Expected Number of customers in \(\mathrm{M} / \mathrm{M} / \mathrm{m} / \mathrm{m}\)
- The actual arrival rate into the system:
\[
\lambda_{a}=\lambda\left(1-p_{m}\right)
\]
- Average total delay figure:
\[
E[r]=E[s]=1 / \mu
\]

Why?
- Average number of customers:
\[
E[n]=\lambda_{a} E[s]=\lambda_{a} / \mu
\]```

