

# Minimum Spanning Tree

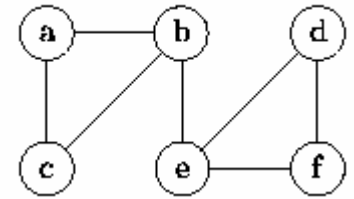
- What is a Minimum Spanning Tree.
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# What is a Minimum Spanning Tree.

- Let  $G = (V, E)$  be a simple, connected, undirected graph that is not edge-weighted.
- A spanning tree of  $G$  is a free tree (i.e., a tree with no root) with  $|V| - 1$  edges that connects all the vertices of the graph.
- Thus a minimum spanning tree for  $G$  is a graph,  $T = (V', E')$  with the following properties:
  - $V' = V$
  - $T$  is connected
  - $T$  is acyclic.
- A spanning tree is called a tree because every acyclic undirected graph can be viewed as a general, unordered tree. Because the edges are undirected, any vertex may be chosen to serve as the root of the tree.

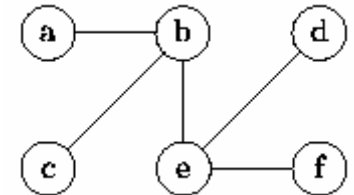
# Constructing Minimum Spanning Trees

- Any traversal of a connected, undirected graph visits all the vertices in that graph. The set of edges which are traversed during a traversal forms a spanning tree.



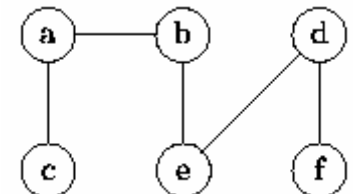
(a) Graph G

- For example, Fig:(b) shows the spanning tree obtained from a breadth-first traversal starting at vertex b.



(b) Breadth-first spanning tree of G rooted at b

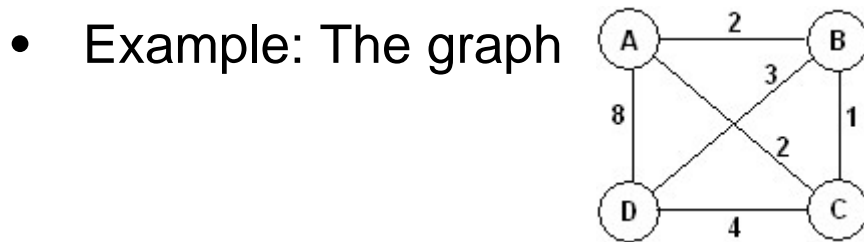
- Similarly, Fig:(c) shows the spanning tree obtained from a depth-first traversal starting at vertex c.



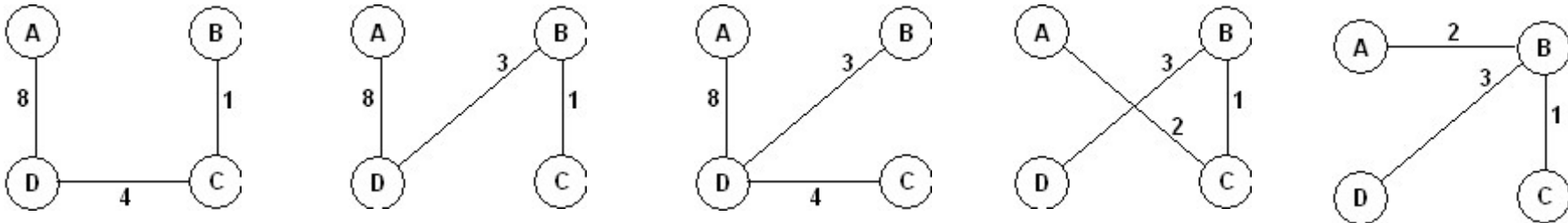
(c) Depth-first spanning tree of G rooted at c

# What is a Minimum-Cost Spanning Tree

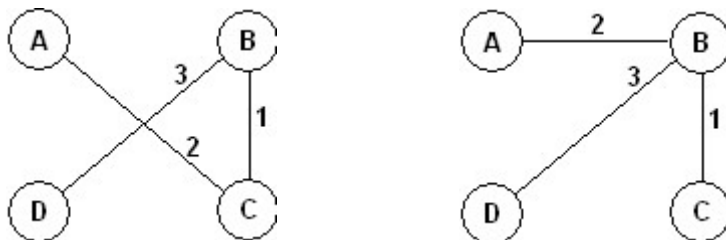
- For an edge-weighted, connected, undirected graph,  $G$ , the total cost of  $G$  is the sum of the weights on all its edges.
- A minimum-cost spanning tree for  $G$  is a minimum spanning tree of  $G$  that has the least total cost.



Has 16 spanning trees. Some are:



The graph has two minimum-cost spanning trees, each with a cost of 6:



# Applications of Minimum-Cost Spanning Trees

Minimum-cost spanning trees have many applications. Some are:

- Building cable networks that join  $n$  locations with minimum cost.
- Building a road network that joins  $n$  cities with minimum cost.
- Obtaining an independent set of circuit equations for an electrical network.
- In pattern recognition minimal spanning trees can be used to find noisy pixels.

# Prim's Algorithm

- Prim's algorithm finds a minimum cost spanning tree by selecting edges from the graph one-by-one as follows:
- It starts with a tree,  $T$ , consisting of the starting vertex,  $x$ .
- Then, it adds the shortest edge emanating from  $x$  that connects  $T$  to the rest of the graph.
- It then moves to the added vertex and repeats the process.

Consider a graph  $G=(V, E)$ ;

Let  $T$  be a tree consisting of only the starting vertex  $x$ ;

while ( $T$  has fewer than  $|V|$  vertices)

{

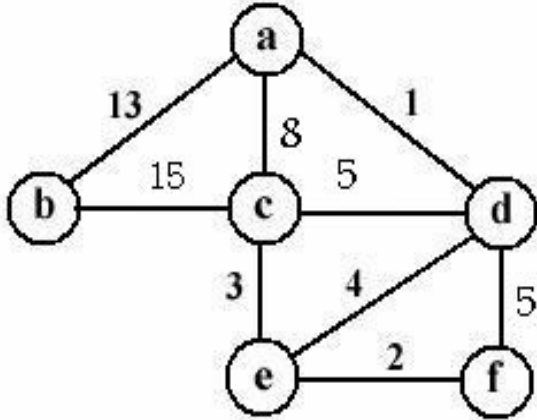
    find a smallest edge connecting  $T$  to  $G-T$ ;

    add it to  $T$ ;

}

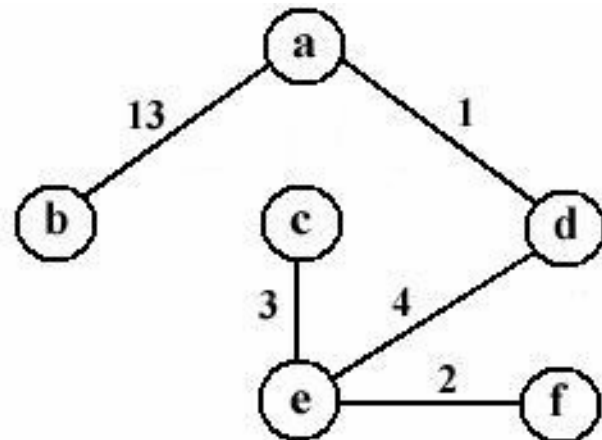
# Example

Trace Prim's algorithm starting at vertex a:



Pass:	initially	1	2	3	4	5	6	weight	vl
Active vertex:		a	d	e	f	c	b		
a	0							0	-
b	$\infty$	13	13	13	13	13		13	a
c	$\infty$	8	5	3	3			3	e
d	$\infty$	1						1	a
e	$\infty$	$\infty$	4					4	d
f	$\infty$	$\infty$	5	2				2	e

The resulting minimum-cost spanning tree is:



# Implementation of Prim's Algorithm.

- Prim's algorithm can be implemented similar to the Dijkstra's algorithm as shown below:

```
public static Graph primAlgorithm(Graph g, Vertex start){
    int n = g.getNumberOfVertices();
    Entry table[] = new Entry[n];
    for(int v = 0; v < n; v++)
        table[v] = new Entry();

    table[g.getIndex(start)].distance = 0;
    PriorityQueue queue = new BinaryHeap(g.getNumberOfEdges());
    queue.enqueue(new Association(new Integer(0), start));
    while(!queue.isEmpty()) {
        Association association = (Association)queue.dequeueMin();
        Vertex v1 = (Vertex) association.getValue();
        int n1 = g.getIndex(v1);
        if(!table[n1].known){
            table[n1].known = true;
            Iterator p = v1.getEmanatingEdges();
            while (p.hasNext()){
                Edge edge = (Edge) p.next();
                Vertex v2 = edge.getMate(v1);
                int n2 = g.getIndex(v2);
                Integer weight = (Integer) edge.getWeight();
                int d = weight.intValue();
```



# Implementation of Prim's Algorithm Cont'd

```
        if(!table[n2].known && table[n2].distance > d){
            table[n2].distance = d; table[n2].predecessor = v1;
            queue.enqueue(new Association(new Integer(d), v2));
        }
    }
}
}
GraphAsLists result = new GraphAsLists(false);
Iterator it = g.getVertices();
while (it.hasNext()){
    Vertex v = (Vertex) it.next();
    result.addVertex(v.getLabel());
}
it = g.getVertices();
while (it.hasNext()){
    Vertex v = (Vertex) it.next();
    if (v != start){
        int index = g.getIndex(v);
        String from = v.getLabel();
        String to = table[index].predecessor.getLabel();
        result.addEdge(from, to, new Integer(table[index].distance));
    }
}
return result;
}
```

# Kruskal's Algorithm.

- Kruskal's algorithm also finds the minimum cost spanning tree of a graph by adding edges one-by-one.

enqueue edges of  $G$  in a queue in increasing order of cost.

$T = \phi$  ;

while(queue is not empty){

    dequeue an edge  $e$ ;

    if( $e$  does not create a cycle with edges in  $T$ )

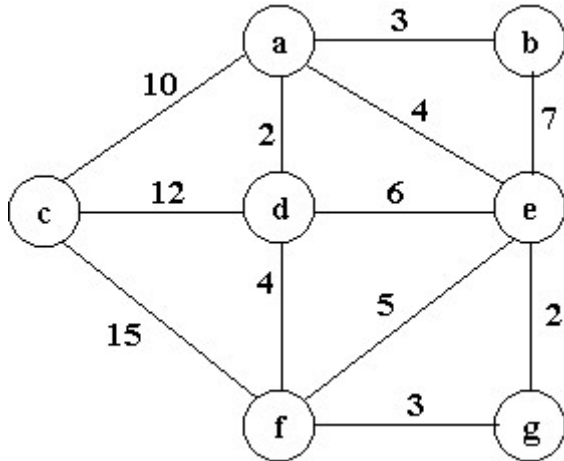
        add  $e$  to  $T$ ;

}

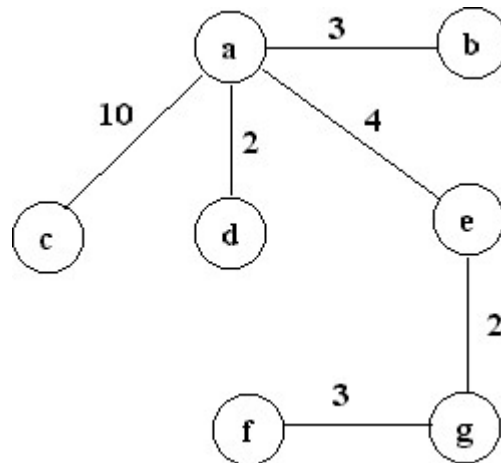
return  $T$ ;

# Example for Kruskal's Algorithm.

Trace Kruskal's algorithm in finding a minimum-cost spanning tree for the undirected, weighted graph given below:



edge	ad	eg	ab	fg	ae	df	ef	de	be	ac	cd	cf
weight	2	2	3	3	4	4	5	6	7	10	12	15
insertion status	√	√	√	√	√	x	x	x	x	√	x	x
insertion order	1	2	3	4	5					6		



The minimum cost is: 24

# Implementation of Kruskal's Algorithm

```
public static Graph kruskalsAlgorithm(Graph g){
    Graph result = new GraphAsLists(false);
    Iterator it = g.getVertices();
    while (it.hasNext()){
        Vertex v = (Vertex)it.next();
        result.addVertex(v.getLabel());
    }
    PriorityQueue queue = new BinaryHeap(g.getNumberOfEdges());
    it = g.getEdges();
    while(it.hasNext()){
        Edge e = (Edge) it.next();
        if (e.getWeight()==null)
            throw new IllegalArgumentException("Graph is not weighted");
        queue.enqueue(e);
    }

    while (!queue.isEmpty()){
        Edge e = (Edge) queue.dequeueMin();
        String from = e.getFromVertex().getLabel();
        String to = e.getToVertex().getLabel();
        if (!result.isReachable(from, to))
            result.addEdge(from,to,e.getWeight());
    }
    return result;
}
```

adds an edge only, if it does not create a cycle

# Implementation of Kruskal's Algorithm – Cont'd

```
public abstract class AbstractGraph implements Graph {
    public boolean isReachable(String from, String to){
        Vertex fromVertex = getVertex(from);
        Vertex toVertex = getVertex(to);
        if (fromVertex == null || toVertex==null)
            throw new IllegalArgumentException("Vertex not in the graph");
        PathVisitor visitor = new PathVisitor(toVertex);
        this.preorderDepthFirstTraversal(visitor, fromVertex);
        return visitor.isReached();
    }

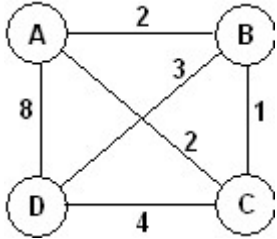
    private class PathVisitor implements Visitor {
        boolean reached = false;
        Vertex target;
        PathVisitor(Vertex t){target = t;}

        public void visit(Object obj){
            Vertex v = (Vertex) obj;
            if (v.equals(target)) reached = true;
        }
        public boolean isDone(){return reached;}
        boolean isReached(){return reached;}
    }
}
```

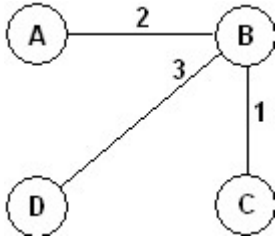
# Prim's and Kruskal's Algorithms

Note: It is not necessary that Prim's and Kruskal's algorithm generate the same minimum-cost spanning tree.

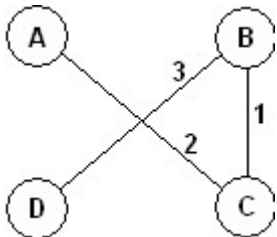
For example for the graph:



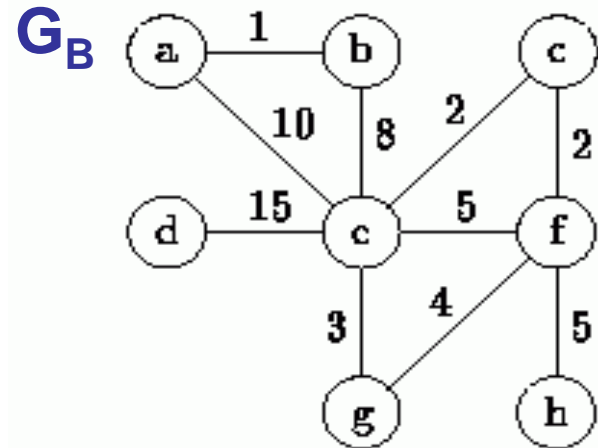
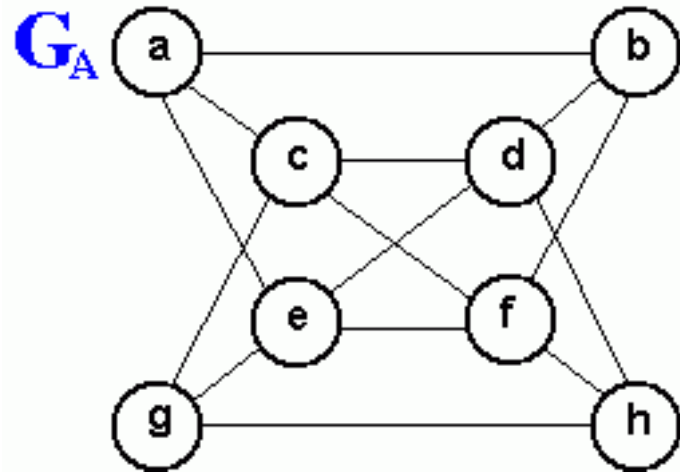
Kruskal's algorithm (that imposes an ordering on edges with equal weights) results in the following minimum cost spanning tree:



The same tree is generated by Prim's algorithm if the start vertex is any of: A, B, or D; however if the start vertex is C the minimum cost spanning tree is:



# Review Questions



1. Find the breadth-first spanning tree and depth-first spanning tree of the graph  $G_A$  shown above.
2. For the graph  $G_B$  shown above, trace the execution of Prim's algorithm as it finds the minimum-cost spanning tree of the graph starting from vertex  $a$ .
3. Repeat question 2 above using Kruskal's algorithm.