

*King Fahd University of Petroleum and Minerals*  
*College of Computer Science and Engineering*  
*Computer Engineering Department*

**COE 301 COMPUTER ORGANIZATION**  
**ICS 233: COMPUTER ARCHITECTURE & ASSEMBLY LANGUAGE**  
**Term 171 (Fall 2017-2018)**  
**Major Exam 1**  
**Saturday Oct. 21, 2017**

**Time: 120 minutes, Total Pages: 10**

**Name: \_KEY\_\_\_\_\_ ID: \_\_\_\_\_ Section: \_\_\_\_\_**

**Notes:**

- Do not open the exam book until instructed
- Answer all questions
- All steps must be shown
- Any assumptions made must be clearly stated
- No calculators are allowed to be used in the exam

<b>Question</b>	<b>Max Points</b>	<b>Score</b>
<b>Q1</b>	<b>28</b>	
<b>Q2</b>	<b>11</b>	
<b>Q3</b>	<b>17</b>	
<b>Total</b>	<b>56</b>	

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[28 Points]

(Q1) Fill in the blank in each of the following questions:

- (1) Assuming 12-bit unsigned number representation, the binary number 1111 1111 0000 is equal to the decimal number 4080.
- (2) Assuming 16-bit signed 2`s complement representation, the hexadecimal number FEA0 is equal to the decimal number -352.
- (3) Two advantages of programming in assembly language are space and time efficiency and accessibility to system hardware.
- (4) Two advantages of programming in high-level language are programs are portable and program development and maintenance are faster.
- (5) The instruction set architecture of a processor consists of the instruction set, memory and programmer accessible registers.
- (6) With a 24-bit address bus and 32-bit data bus, the maximum memory size (assuming byte addressable memory) that can be accessed by a processor is  $2^{24}=16$  MB and the maximum number of bytes that can be read or written in a single cycle is  $32/8=4$ .
- (7) The advantage of static RAM over dynamic RAM is that it is faster but the disadvantage is that it is less dense and more expensive.
- (8) Given a magnetic disk with the following properties:
  - Time of one rotation is 8 ms
  - Average seek = 8 ms, Sector = 512 bytes, Track = 200 sectors

The average time to access a block of 100 consecutive sectors is  
 $8 \text{ ms} + 0.5 * 8 \text{ ms} + 100/200 * 8 \text{ ms} = 16 \text{ ms}$ .

(9) Assuming variable Array is defined as shown below:

```
Array: .word 10  
       .half 11, 12  
       .byte 13, 14, 15, 16
```

The content of register \$t1 (in hexadecimal) after executing the following sequence of instructions is 0x000c000b.

```
la $t0, Array  
lw $t1, 4($t0)
```

(10) The pseudo instruction *bgt \$s2, 10, Next* is implemented by the following minimum MIPS instructions:

```
slti $at, $s2, 11  
bne $at, $0, Next
```

(11) The pseudo instruction *li \$t0, 0x12345678* is implemented by the following minimum MIPS instructions:

```
lui $t0, 0x1234  
ori $t0, $t0, 0x5678
```

(12) The pseudo instruction *rol \$s0, \$s0, 4* (\$s0 is rotated to the left by 4 bits and stored in \$s0) is implemented by the following minimum MIPS instructions:

```
srl $at, $s0, 28  
sll $s0, $s0, 4  
or $s0, $s0, $at
```

(13) Assuming that \$a0 contains an Alphabetic character, the instruction andi \$a0, \$a0, 0xDF will make the character stored in \$a0 always upper case. Note that the ASCII code of character 'A' is 0x41 while that of character 'a' is 0x61.

(14) Assume that the instruction *bne \$t0, \$t1, NEXT* is at address 0x00400040 in the text segment, and the label NEXT is at address 0x00400028. Then, the value stored in the assembled instruction for the label NEXT is (0x00400028-0x00400044)/4=FFF9.

- (15) Assuming that variable Array is defined as shown below:

```
Array2: .half -2,-3, 4, 5
```

After executing the following sequence of instructions, the content of the two registers (in hexadecimal) is \$t1=000000FF and \$t2=FFFFFFFD.

```
la $t0, Array2
lbu $t1, 1($t0)
lh $t2, 2($t0)
```

- (16) Assuming the following data segment, and assuming that the first variable X is given the address **0x10010000**, then the addresses for variables Y and Z will be 0x10010006 and 0x10010010.

```
.data
X:   .byte 10, 11, 12, 13, 14
Y:   .half 15, 16, 17, 18
Z:   .word 19, 20
```

- (17) To multiply the **signed** content of register \$t0 by 112 without using multiplication instructions, we use the following minimum MIPS instructions (HINT:  $112=16*7$ ):

```
sll $t1, $t0, 4
sll $t0, $t1, 3
sub $t0, $t0, $t1
```

**(Q2) Answer each of the following questions. Show how you obtained your answer:**

**(i)** Given that **TABLE** is defined as: **TABLE: .asciiz "Aiman El-Maleh"**

Determine the content of register **\$t0** after executing the following code:

```

        xor $t0, $t0, $t0
        la $t1, TABLE
        li $t2, 'a'
Next:   lbu $t3, ($t1)
        beq $t3, $zero, ENL
        ori $t3, $t3, 0x20
        addi $t1, $t1, 1
        bne $t2, $t3, Next
        addi $t0, $t0, 1
        j Next
ENL:

```

The content of register  $\$t0=3$  as the program counts the number of characters equal to 'A' or 'a' in TABLE.

**(ii)** Determine the content of register  $\$t1$  after executing the following code:

```

        li $t0, 0x1234
        xor $t1, $t1, $t1
AGAIN:  andi $t2, $t0, 0xf
        add $t1, $t1, $t2
        srl $t0, $t0, 4
        bne $t0, $zero, AGAIN

```

The content of register  $\$t1=0xA$  as the program computes the sum of the hexadecimal digits in register  $\$t0$ .

(iii) Given that **TABLE** is defined as: **TABLE: .word 90, 70, 80, 60, 100**

Determine the content of register **\$v0** after executing the following code:

```
        la    $a0, TABLE
        addi  $a1, $a0, 16
        lw   $v0, 0($a0)
loop:   addi  $a0, $a0, 4
        lw   $t1, 0($a0)
        bge  $t1, $v0, skip
        move $v0, $t1
skip:   bne  $a0, $a1, loop
```

The content of register  $\$v0=0x3C=60$  as the program computes the minimum of the numbers stored in **TABLE**.

[17 points]

(Q3) Write separate MIPS assembly code fragments with minimum instructions to implement each of the given requirements. You can use pseudo instructions in your solution.

- (i) [10 points] Write a MIPS code fragment that returns the maximum integer value found in a user-specified row number of a  $32 \times 32$  matrix **A** of 32-bit signed integers. The program should read the desired row number from the user and check that it is in the range between 0 and 31. If not, the program should display the error message “**Row number is out of range.**” and terminate. Otherwise, the program should display the message “**Maximum integer in the row is** ” and the value of the maximum integer found in the specified row, and then terminate. Assume that matrix **A** is already stored in memory.

```
.data
prompt:      .asciiz      "Please enter a row number between 0 and 31: "
outofrange: .asciiz      "Row number is out of range.\n"
outmsg:      .asciiz      "Maximum integer in the row is "
```

```
.text
.globl main
main:
    la    $a0,prompt    # display prompt string
    li    $v0,4
    syscall
    li    $v0,5         # read row number into $t0
    syscall
    move  $t0,$v0
    bltz  $t0,error     # check row boundary
    addiu $t1,$t0,-31   # If $t0 > 31, then result of ($t0-31) > 0
    bgtz  $t1,error
    la    $t1,A         # compute starting location of 1st element in desired row
    sll   $t2,$t0,5     # $t2 = i*32 (ixCOL+0)
    sll   $t2,$t2,2     # $t2 = i*32*4 (ixCOL+0)x(int size)
    addu  $t2,$t1,$t2   # $t2 = address of 1st element in desired row
    li    $t3,31       # max j = 31
    lw    $t4,0($t2)    # read 1st element of desired row & set as maximum
loop:
    addiu $t2,$t2,4     # increment index to point to next row element
    lw    $t5,0($t2)    # read next element of desired row
    ble   $t5,$t4,next  # next element ($t5) <= current max ($t4)?
    move  $t4,$t5      # No -> set max ($t4) = next element ($t5)
next:
    addiu $t3,$t3,-1    # prepare for next row element
    bgtz  $t3,loop
    la    $a0,outmsg    # display prompt string
    li    $v0,4
    syscall
    move  $a0,$t4       # output $t4 = maximum in desired row
    li    $v0,1
    syscall
    j     exit
error:
    la    $a0,outofrange
    li    $v0,4
    syscall
exit:
    li    $v0,10        # exit
    syscall
```

- (ii) [7 points] Given two arrays **A** and **B**, write the smallest MIPS assembly fragment for the following computation. Assume that register **\$s0** will be used to store **cnt** and assume that the following registers have the mentioned values: register **\$s1** = number of elements, **N**, in each array, register **\$s2** = base address of the array **A**, and register **\$s3** = base address of the array **B**. Each array element is a 32-bit signed integer. Assume that **N > 0**. Insert comments to clarify the meaning of instructions and the use of registers.

```
int cnt = 0;
for (i=0; i != N; i++) {
    if (((A[i] - B[i]) > 5) || ((B[i] - A[i]) > 5)) cnt = cnt + 1;
}
```

```
li    $s0,0          # $s0 = cnt = 0
loop:
lw    $t0,0($s2)     # $t0 = A[i]
lw    $t1,0($s3)     # $t1 = B[i]
addiu $t2,$t0,5      # $t2 = A[i]+5
addiu $t3,$t1,5      # $t3 = B[i]+5
bgt   $t0,$t3,incr   # Check if (A[i]-B[i]>5)
ble   $t1,$t2,done   # Check if (B[i]-A[i]>5)
incr:
addiu $s0,$s0,1      # cnt++
done:
addiu $s2,$s2,4      # point to A[i+1]
addiu $s3,$s3,4      # point to B[i+1]
addiu $s1,$s1,-1     # decrement loop index
bne   $s1,$0,loop
```



**MIPS Instructions:**

Instruction	Meaning	R-Type Format						
add \$s1, \$s2, \$s3	\$s1 = \$s2 + \$s3	op = 0	rs = \$s2	rt = \$s3	rd = \$s1	sa = 0	f = 0x20	
addu \$s1, \$s2, \$s3	\$s1 = \$s2 + \$s3	op = 0	rs = \$s2	rt = \$s3	rd = \$s1	sa = 0	f = 0x21	
sub \$s1, \$s2, \$s3	\$s1 = \$s2 - \$s3	op = 0	rs = \$s2	rt = \$s3	rd = \$s1	sa = 0	f = 0x22	
subu \$s1, \$s2, \$s3	\$s1 = \$s2 - \$s3	op = 0	rs = \$s2	rt = \$s3	rd = \$s1	sa = 0	f = 0x23	

Instruction	Meaning	R-Type Format						
and \$s1, \$s2, \$s3	\$s1 = \$s2 & \$s3	op = 0	rs = \$s2	rt = \$s3	rd = \$s1	sa = 0	f = 0x24	
or \$s1, \$s2, \$s3	\$s1 = \$s2   \$s3	op = 0	rs = \$s2	rt = \$s3	rd = \$s1	sa = 0	f = 0x25	
xor \$s1, \$s2, \$s3	\$s1 = \$s2 ^ \$s3	op = 0	rs = \$s2	rt = \$s3	rd = \$s1	sa = 0	f = 0x26	
nor \$s1, \$s2, \$s3	\$s1 = ~( \$s2   \$s3 )	op = 0	rs = \$s2	rt = \$s3	rd = \$s1	sa = 0	f = 0x27	

Instruction	Meaning	R-Type Format						
sll \$s1, \$s2, 10	\$s1 = \$s2 << 10	op = 0	rs = 0	rt = \$s2	rd = \$s1	sa = 10	f = 0	
srl \$s1, \$s2, 10	\$s1 = \$s2 >>> 10	op = 0	rs = 0	rt = \$s2	rd = \$s1	sa = 10	f = 2	
sra \$s1, \$s2, 10	\$s1 = \$s2 >> 10	op = 0	rs = 0	rt = \$s2	rd = \$s1	sa = 10	f = 3	
sllv \$s1, \$s2, \$s3	\$s1 = \$s2 << \$s3	op = 0	rs = \$s3	rt = \$s2	rd = \$s1	sa = 0	f = 4	
srlv \$s1, \$s2, \$s3	\$s1 = \$s2 >>> \$s3	op = 0	rs = \$s3	rt = \$s2	rd = \$s1	sa = 0	f = 6	
srav \$s1, \$s2, \$s3	\$s1 = \$s2 >> \$s3	op = 0	rs = \$s3	rt = \$s2	rd = \$s1	sa = 0	f = 7	

Instruction	Meaning	I-Type Format				
addi \$s1, \$s2, 10	\$s1 = \$s2 + 10	op = 0x8	rs = \$s2	rt = \$s1	imm <sup>16</sup> = 10	
addiu \$s1, \$s2, 10	\$s1 = \$s2 + 10	op = 0x9	rs = \$s2	rt = \$s1	imm <sup>16</sup> = 10	
andi \$s1, \$s2, 10	\$s1 = \$s2 & 10	op = 0xc	rs = \$s2	rt = \$s1	imm <sup>16</sup> = 10	
ori \$s1, \$s2, 10	\$s1 = \$s2   10	op = 0xd	rs = \$s2	rt = \$s1	imm <sup>16</sup> = 10	
xori \$s1, \$s2, 10	\$s1 = \$s2 ^ 10	op = 0xe	rs = \$s2	rt = \$s1	imm <sup>16</sup> = 10	
lui \$s1, 10	\$s1 = 10 << 16	op = 0xf	0	rt = \$s1	imm <sup>16</sup> = 10	

Instruction	Meaning	Format				
j label	jump to label	op <sup>6</sup> = 2	imm <sup>26</sup>			
beq rs, rt, label	branch if (rs == rt)	op <sup>6</sup> = 4	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>	
bne rs, rt, label	branch if (rs != rt)	op <sup>6</sup> = 5	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>	
blez rs, label	branch if (rs <= 0)	op <sup>6</sup> = 6	rs <sup>5</sup>	0	imm <sup>16</sup>	
bgtz rs, label	branch if (rs > 0)	op <sup>6</sup> = 7	rs <sup>5</sup>	0	imm <sup>16</sup>	
bltz rs, label	branch if (rs < 0)	op <sup>6</sup> = 1	rs <sup>5</sup>	0	imm <sup>16</sup>	
bgez rs, label	branch if (rs >= 0)	op <sup>6</sup> = 1	rs <sup>5</sup>	1	imm <sup>16</sup>	

Instruction	Meaning	Format						
slt rd, rs, rt	rd=(rs<rt?1:0)	op <sup>6</sup> = 0	rs <sup>5</sup>	rt <sup>5</sup>	rd <sup>5</sup>	0	0x2a	
sltu rd, rs, rt	rd=(rs<rt?1:0)	op <sup>6</sup> = 0	rs <sup>5</sup>	rt <sup>5</sup>	rd <sup>5</sup>	0	0x2b	
slti rt, rs, imm <sup>16</sup>	rt=(rs<imm?1:0)	0xa	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>			
sltiu rt, rs, imm <sup>16</sup>	rt=(rs<imm?1:0)	0xb	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>			

Instruction		Meaning	I-Type Format			
lb	rt, imm <sup>16</sup> (rs)	rt = MEM[rs+imm <sup>16</sup> ]	0x20	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>
lh	rt, imm <sup>16</sup> (rs)	rt = MEM[rs+imm <sup>16</sup> ]	0x21	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>
lw	rt, imm <sup>16</sup> (rs)	rt = MEM[rs+imm <sup>16</sup> ]	0x23	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>
lbu	rt, imm <sup>16</sup> (rs)	rt = MEM[rs+imm <sup>16</sup> ]	0x24	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>
lhu	rt, imm <sup>16</sup> (rs)	rt = MEM[rs+imm <sup>16</sup> ]	0x25	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>
sb	rt, imm <sup>16</sup> (rs)	MEM[rs+imm <sup>16</sup> ] = rt	0x28	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>
sh	rt, imm <sup>16</sup> (rs)	MEM[rs+imm <sup>16</sup> ] = rt	0x29	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>
sw	rt, imm <sup>16</sup> (rs)	MEM[rs+imm <sup>16</sup> ] = rt	0x2b	rs <sup>5</sup>	rt <sup>5</sup>	imm <sup>16</sup>

## Syscall Services:

Service	\$v0	Arguments / Result
Print Integer	1	\$a0 = integer value to print
Print Float	2	\$f12 = float value to print
Print Double	3	\$f12 = double value to print
Print String	4	\$a0 = address of null-terminated string
Read Integer	5	Return integer value in \$v0
Read Float	6	Return float value in \$f0
Read Double	7	Return double value in \$f0
Read String	8	\$a0 = address of input buffer \$a1 = maximum number of characters to read
Exit Program	10	
Print Char	11	\$a0 = character to print
Read Char	12	Return character read in \$v0