

ICS 233
**Computer Architecture &
Assembly Language**

**MIPS PROCESSOR
INSTRUCTION SET**

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ICS 233
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Lecture 9

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Lecture Outline

- ❑ **Translating IF Statement**
- ❑ **Translating WHILE loop**

Translating an IF Statement

- Consider the following IF statement:
`if (a == b) c = d + e; else c = d - e;`
Assume that a, b, c, d, e are in \$s0, \$s1, \$s2, \$s3, \$s4 respectively
- How to translate the above IF statement?

```
        bne    $s0, $s1, else
        addu   $s2, $s3, $s4
        j     exit
else:    subu   $s2, $s3, $s4
exit:    . . .
```

Compound Expression with AND

- Programming languages use **short-circuit evaluation**
- If first expression is **false**, second expression is **skipped**

```
if (($s1 > 0) && ($s2 < 0)) {$s3++;}
```

```
# One Possible Implementation ...
    bgtz    $s1, L1      # first expression
    j      next        # skip if false
L1: bltz    $s2, L2      # second expression
    j      next        # skip if false
L2: addiu   $s3,$s3,1    # both are true
next:
```

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Better Implementation for AND

```
if (($s1 > 0) && ($s2 < 0)) {$s3++;}
```

The following implementation uses less code

Reverse the relational operator

Allow the program to **fall through** to the second expression

Number of instructions is reduced from 5 to 3

```
# Better Implementation ...
    blez    $s1, next    # skip if false
    bgez    $s2, next    # skip if false
    addiu   $s3,$s3,1    # both are true
next:
```

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Compound Expression with OR

- Short-circuit evaluation for logical OR
- If first expression is true, second expression is skipped

```
if (($s1 > $s2) || ($s2 > $s3)) {$s4 = 1;}
```

- Use fall-through to keep the code as short as possible

```
bgt $s1, $s2, L1    # yes, execute if part
ble $s2, $s3, next  # no: skip if part
L1: ori $s4, $0, 1   # set $s4 to 1
next:
```

- bgt, ble are pseudo-instructions
 - Translated by the assembler to real instructions

TRY THIS

- Translate the IF statement to assembly language
- \$s1 and \$s2 values are unsigned

```
if( $s1 <= $s2 ) {
    $s3 = $s4
}
```

```
bgtu $s1, $s2, next
or $s3, $s4, $0
next:
```

- \$s3, \$s4, and \$s5 values are signed

```
if (($s3 <= $s4) &&
    ($s4 > $s5)) {
    $s3 = $s4 + $s5
}
```

```
bgt $s3, $s4, next
ble $s4, $s5, next
addu $s3, $s4, $s5
next:
```

Translating a WHILE Loop

- Consider the following WHILE statement:

```
i = 0; while (A[i] != k) i = i+1;
```

Where *A* is an array of integers (4 bytes per element)

Assume address *A*, *i*, *k* in *\$s0*, *\$s1*, *\$s2*, respectively

Memory	
...	
A[i]	A+4 <i>i</i>
...	
A[2]	A+8
A[1]	A+4
A[0]	A
...	

- How to translate above WHILE statement?

```

xor    $s1, $s1, $s1    # i = 0
or     $t0, $s0, $0     # $t0 = address A
loop:  lw    $t1, 0($t0) # $t1 = A[i]
      beq   $t1, $s2, exit # exit if (A[i]== k)
      addiu $s1, $s1, 1   # i = i+1
      sll  $t0, $s1, 2    # $t0 = 4*i
      addu $t0, $s0, $t0 # $t0 = address A[i]
      j    loop
exit:  . . .

```

Using Pointers to Traverse Arrays

- Consider the same WHILE loop:

```
i = 0; while (A[i] != k) i = i+1;
```

Where address of *A*, *i*, *k* are in *\$s0*, *\$s1*, *\$s2*, respectively

- We can use a **pointer** to traverse array *A*

Pointer is incremented by 4 (faster than indexing)

```

or     $t0, $s0, $0     # $t0 = $s0 = addr A
j     cond              # test condition
loop:  addiu $s1, $s1, 1 # i = i+1
      addiu $t0, $t0, 4 # point to next
cond:  lw    $t1, 0($t0) # $t1 = A[i]
      bne   $t1, $s2, loop # loop if A[i]!= k

```

- Only 4 instructions (rather than 6) in loop body

Copying a String

The following code copies source string to target string
Address of source in \$s0 and address of target in \$s1
Strings are terminated with a null character (C strings)

```
i = 0;  
do {target[i]=source[i]; i++;} while (source[i]!=0);
```

```
or    $t0, $s0, $0    # $t0 = pointer to source  
or    $t1, $s1, $0    # $t1 = pointer to target  
L1: lb    $t2, 0($t0)    # load byte into $t2  
sb    $t2, 0($t1)    # store byte into target  
addiu $t0, $t0, 1    # increment source pointer  
addiu $t1, $t1, 1    # increment target pointer  
bne   $t2, $zero, L1 # loop until NULL char
```

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Summing an Integer Array

```
sum = 0;  
for (i=0; i<n; i++) sum = sum + A[i];
```

Assume \$s0 = array address, \$s1 = array length = n

```
or    $t0, $s0, $0    # $t0 = address A[i]  
xor   $t1, $t1, $t1    # $t1 = i = 0  
xor   $s2, $s2, $s2    # $s2 = sum = 0  
L1: lw    $t2, 0($t0)    # $t2 = A[i]  
addu  $s2, $s2, $t2    # sum = sum + A[i]  
addiu $t0, $t0, 4    # point to next A[i]  
addiu $t1, $t1, 1    # i++  
bne   $t1, $s1, L1    # loop if (i != n)
```

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