King Fahd University of Petroleum and Minerals College of Computer Science and Engineering Computer Engineering Department

COE 301 COMPUTER ORGANIZATION ICS 233: COMPUTER ARCHITECTURE & ASSEMBLY LANGUAGE Term 161 (Fall 2016-2017) Major Exam 2 Saturday Dec. 10, 2016

Time: 150 minutes, Total Pages: 13

 Name:
 ID:
 Section:

Notes:

- Do not open the exam book until instructed
- Answer all questions
- All steps must be shown
- Any assumptions made must be clearly stated

Question	Max Points	Score
Q1	20	
Q2	10	
Q3	17	
Q4	23	
Total	70	

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[20 Points]

(Q1) Write MIPS programs with <u>minimal</u> used instructions. Use <u>MIPS programming</u> <u>convention</u> in saving and restoring registers in procedures.

- (i) [4 points] Write a procedure GetAscii that receives a single hexadecimal digit in register \$a0 and returns the ASCII code of that digit in register \$v0. For example, if \$a0=0x9 the procedure will return 0x39 in \$v0 and if \$a0=0xA, the procedure will return 0x41 in \$a0. Assume the use of capital letters for the digits A to F.
- (ii) [11 points] Write a procedure **DispHex** that receives a number in register \$a0 and displays the hexadecimal representation of that number. Only significant hexadecimal digits need to be displayed. For example, if \$a0=0x1E, the procedure will display 1E. Your DisHex procedure should utilize the GetAscii procedure.
- (iii) [5 points] Write a MIPS program that asks the user to enter a decimal number and displays its hexadecimal content using the **DispHex** procedure. Two sample runs of the program are given below:

Enter a decimal number: 260 Your number in hexadecimal is: 0x104

Enter a decimal number: 0 Your number in hexadecimal is: 0x0 Page 3 of 13

(Q2)

(i) [4 Points] Given that Multiplicand=0111 and Multiplier=1011 are signed 2's complement numbers, show the signed multiplication of Multiplicand by Multiplier. The result of the multiplication should be an 8 bit signed number in HI and LO registers. Show the steps of your work.

[

Ite	eration	Multiplicand	Sign	Product =HI,LO		
0	Initialize					
1						
2						
3						
4						

(ii) [6 Points] Given that Dividend=0111 and Divisor=1011 are signed 2's complement numbers, show the signed division of Dividend by Divisor. The result of division should be stored in the Remainder and Quotient registers. Show the steps of your work, and show the final result.

Ite	eration	Remainder (HI)	Quotient (LO)	Divisor	Difference
0	Initialize				
1					
2					
3					
4					
Fin	nal Result				

(Q3)

1. [2 Points] Find the decimal value of the following single precision float:

 $[0,\,1000\,\,1000,\,0000\,\,0100\,\,1100\,\,0000\,\,0000\,\,000]$

2. [2 Points] Find the decimal value of the following single precision float:

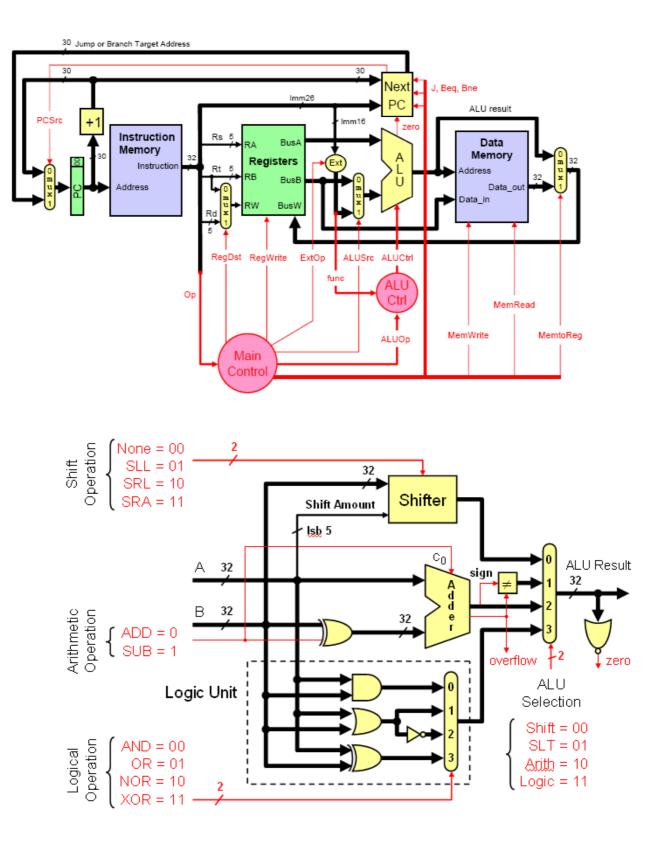
3. **[3 Points]** Find the normalized single precision representation of –59.625.

4. **[4 Points]** Round the given single precision float with the given GRS bits using the following rounding modes showing the resulting normalized number:

	+1.111	1111	1111	1111	1111	1111	GRS 100	x	2 ⁻¹²⁷	
Zero:		[]
+infinity:		[]
-infinity:	l	[]
Nearest E	Even:	[]

5. **[6 Points]** Find the normalized <u>difference</u> between **A** and **B** (i.e., A-B) by using rounding to <u>+infinity</u>. Perform the operation using **guard**, **round** and **sticky** bits.

 (Q4) Consider the single-cycle datapath and control given below along with ALU design for the MIPS processor implementing a subset of the instruction set:



(i) Show the control signals generated for the execution of the following instructions by filling the table given below: (5 points)

Op	RegDst	RegWrite	ExtOp	ALUSrc	ALUOp	Beq	Bne	J	MemRead	MemWrite	MemtoReg
R-type											
slti											
SW											
beq											
j											

(ii) Excluding the ALUOp, Beq, Bne and J signals, show the design of the control unit for the control signals given in the table above based on the given instructions. Assume that the opcode of these instructions is a 6-bit opcode such that the opcode for R-type instructions is 0, the opcode for slti is 1, the opcode for sw is 2, and so on for the rest of the instructions. (5 points)

(iii) Show the design of the Next PC block. (4 points)

- (iv) We wish to add the following instructions to the MIPS single-cycle datapath. Add any necessary datapath modifications and control signals needed for the implementation of these instructions. Show only the <u>modified</u> and <u>added</u> components to the datapath.
 - a. sra (3 points)

	Instruction	Meaning	Format					
sra	rd, rt, imm ⁵	rd= rt>>imm ¹⁶	$Op^6 = 0$	0	rt ⁵	rd ⁵	Imm ⁵	f ⁵ =3

b. jr (3 points)

Instruction	Meaning	Format						
jr rs	PC=rs	$op^6 = 0$	rs ⁵	0	0	0	8	

- (v) Assume that the propagation delays for the major components used in the datapath are as follows:
 - Instruction and data memories: 120 ps
 - ALU and adders: 30 ps
 - Register file access (read or write): 14 ps
 - Main control: 8 ps
 - ALU control: 7 ps

Ignore the delays in the multiplexers, PC access, extension logic, and wires. What is the cycle time for the single-cycle datapath given above? (**3 points**)

Syscall Services:

Service	\$v0	Arguments / Result
Print Integer	1	<pre>\$a0 = integer value to print</pre>
Print Float	2	<pre>\$f12 = float value to print</pre>
Print Double	3	<pre>\$f12 = double value to print</pre>
Print String	4	<pre>\$a0 = address of null-terminated string</pre>
Read Integer	5	Return integer value in <mark>\$v0</mark>
Read Float	6	Return float value in <mark>\$f0</mark>
Read Double	7	Return double value in <mark>\$f0</mark>
Read String	8	<pre>\$a0 = address of input buffer \$a1 = maximum number of characters to read</pre>
Exit Program	10	
Print Char	11	<pre>\$a0 = character to print</pre>
Read Char	12	Return character read in \$v0

MIPS Instructions:

Inct	ruction	Meaning				РТ	ne F	orma	+	
add		Meaning \$s1 = \$s2 + \$s3	00 -	0	c – ¢c		ype F			f = 0x20
		\$s1 = \$s2 + \$s3	-	_		_			sa = 0	
sub		\$s1 = \$s2 - \$s3		_		_				f = 0x21
		\$s1 = \$s2 - \$s3	· ·			_				f = 0x22
			- 40		5 - ψ.					1 - 0/20
	ruction	Meaning					ype F			
		\$s1 = \$s2 & \$s3				_				f = 0x24
		\$s1 = \$s2 \$s3	· ·	_		_			-	f = 0x25
		\$s1 = \$s2 ^ \$s3								f = 0x26
nor S	\$s1, \$s2, \$s3	\$s1 = ~(\$s2 \$s3)	op =	: 0 r:	s = \$s	2 rt =	\$s3 ro	1 = \$s′	1 sa = 0	f = 0x27
Inst	ruction	Meaning					ype F			
sll	\$s1,\$s2,10	\$s1 = \$s2 << 10	op :		rs = 0	_			sa = 10	
srl	\$s1,\$s2,10	\$s1 = \$s2>>>10	op :		rs = 0	_	_		sa = 10	-
sra	\$s1, \$s2, 10	\$s1 = \$s2 >> 10	· ·		rs = 0	_			sa = 10	
sllv	\$s1,\$s2,\$s3	\$s1 = \$s2 << \$s3	<u> </u>	_		_	_		sa = 0	f = 4
srlv	\$s1,\$s2,\$s3	\$s1 = \$s2>>>\$s3	<u> </u>			_			sa = 0	f = 6
srav	\$s1,\$s2,\$s3	\$s1 = \$s2 >> \$s3	op :	= 0	rs = \$	s3 rt =	\$s2 rd	= \$s1	sa = 0	f = 7
Inst	ruction	Meaning				I-T	ype F	orma	it	
addi	\$s1, \$s2, 10	\$s1 = \$s2 + 10	ор	= 0x	8 rs	= \$s2	rt = \$	s1	imm ¹⁶	= 10
addiu	ı \$s1, \$s2, 10	\$s1 = \$s2 + 10	_			= \$s2			imm ¹⁶ = 10	
andi	\$s1, \$s2, 10					= \$s2			imm ¹⁶ = 10	
ori	\$s1, \$s2, 10		·			= \$s2	rt = \$		imm ¹⁶	
xori	\$s1, \$s2, 10				_	= \$s2	rt = \$		imm ¹⁶	
lui	\$s1, 10	\$s1 = 10 << 16	op	= 0x	.I	0	rt = \$		imm ¹⁶	= 10
Inst	ruction	Meaning					For	mat		
j	label	jump to label		op ⁶	= 2			imm	26	
beq	rs, rt, label	branch if (rs ==	rt)	op ⁶	= 4	rs ⁵	rt ⁵		imm ¹⁰	6
bne	rs, rt, label	branch if (rs !=	rt)	op ⁶	= 5	rs ⁵	rt ⁵		imm ¹⁰	6
blez	rs, label	branch if (rs<=0))	op ⁶	= 6	rs ⁵	0		imm ¹⁰	6
bgtz	rs, label	branch if (rs > 0))	op ⁶	= 7	rs ⁵	0		imm ¹⁰	6
bltz	rs, label	branch if (rs < 0))	op ⁶	= 1	rs ⁵	0		imm ¹⁰	6
bgez	rs, label	branch if (rs>=0	· ·				6			
Instr	ruction	Meaning					For	mat		
slt	rd, rs, rt	rd=(rs <rt?1:0< td=""><td>))</td><td>op⁶</td><td>= 0</td><td>rs⁵</td><td>rt⁵</td><td>rd⁵</td><td>0</td><td>0x2a</td></rt?1:0<>))	op ⁶	= 0	rs ⁵	rt ⁵	rd ⁵	0	0x2a
sltu	rd, rs, rt	rd=(rs <rt?1:0< td=""><td></td><td>· ·</td><td>= 0</td><td>rs⁵</td><td>rt⁵</td><td>rd⁵</td><td></td><td>0x2b</td></rt?1:0<>		· ·	= 0	rs ⁵	rt ⁵	rd ⁵		0x2b
slti	rt, rs, imm ¹⁶		·	· ·	xa	rs ⁵	rt ⁵	1	imm ¹	6
sltiu	rt, rs, imm ¹⁶	-		0	xb	rs ⁵	rt ⁵		imm ¹	
			-							

Instruction			I-Typ	e Forr	nat		
lb rt, imm ¹⁶ (rs)	rt = MEM[rs+imn	n ¹⁶] 0x2	0 rs	⁵ rt	5	imn	n ¹⁶
Ih rt, imm ¹⁶ (rs)	rt = MEM[rs+imn	n ¹⁶] 0x2	1 rs	⁵ rt	5	imn	n ¹⁶
lw rt, imm ¹⁶ (rs)	rt = MEM[rs+imn	n ¹⁶] 0x2	3 rs	⁵ rt	5	imn	n ¹⁶
Ibu rt, imm16(rs)	rt = MEM[rs+imn	n ¹⁶] 0x2	4 rs	⁵ rt	5	imn	n ¹⁶
Ihu rt, imm16(rs)	rt = MEM[rs+imn	n ¹⁶] 0x2	5 rs	⁵ rt	5	imn	n ¹⁶
sb rt, imm16(rs)	MEM[rs+imm ¹⁶]	= rt 0x2	8 rs	⁵ rt	5	imn	n ¹⁶
sh rt, imm16(rs)	MEM[rs+imm ¹⁶]	= rt 0x2	9 rs	⁵ rt	5	imn	n ¹⁶
sw rt, imm16(rs)	MEM[rs+imm ¹⁶]	= rt 0x2	b rs	⁵ rt	5	imn	n ¹⁶
Instruction	Meaning			For	mat		
jal label S	\$31=PC+4, jump	op ⁶ = 3			_imm ²	26	
jr Rs	PC = Rs	op ⁶ = 0	rs ⁵	0	0	0	8
jalr Rd, Rs F	d=PC+4, PC=Rs	op ⁶ = 0	rs ⁵	0	rd ⁵	0	9
Instruction	Meaning			For	mat		
mult Rs, Rt	Hi, Lo = <u>Rs</u> × <u>Rt</u>	op ⁶ = 0	Rs⁵	Rt⁵	0	0	0x18
multu Rs, Rt	Hi, Lo = <u>Rs</u> × <u>Rt</u>	op ⁶ = 0	Rs⁵	Rt⁵	0	0	0x19
mul Rd, Rs, Rt	Rd = <u>Rs</u> × <u>Rt</u>	0x1c	Rs⁵	Rt⁵	Rd⁵	0	0x02
div <u>Rs</u> , <u>Rt</u>	Hi, Lo = Rs / Rt	op ⁶ = 0	Rs⁵	Rt⁵	0	0	0x1a
divu Rs, Rt	Hi, Lo = Rs / Rt	op ⁶ = 0	Rs⁵	Rt⁵	0) 0 0x1b	
mfhi Rd	Rd = Hi	op ⁶ = 0	0	0	Rd⁵	0	0x10
mflo Rd	Rd = Lo	op ⁶ = 0	0	0	Rd⁵	0	0x12