

3rd April
→ 15:30 - 17:00Primal

$$\min c^t x \quad (3)$$

s.t. $Ax = b$
 $x \geq 0$

Dual

$$\max \lambda^t b \quad (4)$$

$$\text{s.t. } \lambda^t A \leq c^t$$

$$\lambda \geq 0$$

Duality theorem

Lemma If x and λ are feasible for (3) and (4), respectively, then $c^t x \geq \lambda^t b$.

If either of (3) or (4) has a finite optimal solution, so does the other, and the corresponding values of the objective functions are equal. If either problem has unbounded objective, the other problem has no feasible solution.

- For (3) suppose that $x = (x_B, 0)$ with basis B is a feasible basic optimal solution.

We know that $x_B = \bar{B}^{-1}b$ and reduced cost vector is $r_D^t = g_D^t - g_B^t \bar{B}^{-1} D$
 $(r_N^t = g_N^t - g_B^t \bar{B}^{-1} N)$

Since $r_N^t \geq 0 \rightarrow g_B^t \bar{B}^{-1} N \leq g_N^t$.

Let us suppose that $\lambda = g_B^t \bar{B}' \rightarrow \lambda^t A = [\lambda^t B, \lambda^t N] = [g_B^t; g_B^t \bar{B}^{-1} N] \leq [g_B^t; g_D^t] = c^t$
 $\rightarrow \lambda^t A \leq c^t$, so λ is feasible for the dual.

and $\lambda^t b = g_B^t \bar{B}' b = g_B^t x_B \rightarrow$ this establishes the optimality of λ for the dual.

Duality

Primal

obj max

 \leq constraint \geq " $=$ "

positive variable

negative variable

free variable

Dual

obj min

positive variable

negative variable

free variable

 \leq constraint \geq constraint $=$ constraint

Ex

$$\max 2x + 3y + 5z$$

s.t.

$$\begin{cases} x + 2y + 3z \leq 3 \\ 3x + y + z \leq 4 \\ x + y + 2z \leq 5 \end{cases}$$

$$x, y, z \geq 0$$

=

$$\min 3\alpha + 4\beta + 5\gamma$$

s.t.

$$\begin{cases} \alpha + 3\beta + \gamma \geq 2 \\ 2\alpha + \beta + \gamma \geq 3 \\ \alpha + \beta + 2\gamma \geq 5 \end{cases}$$

$$\alpha, \beta, \gamma \geq 0$$

Ex 1:

$$\max 2x + y$$

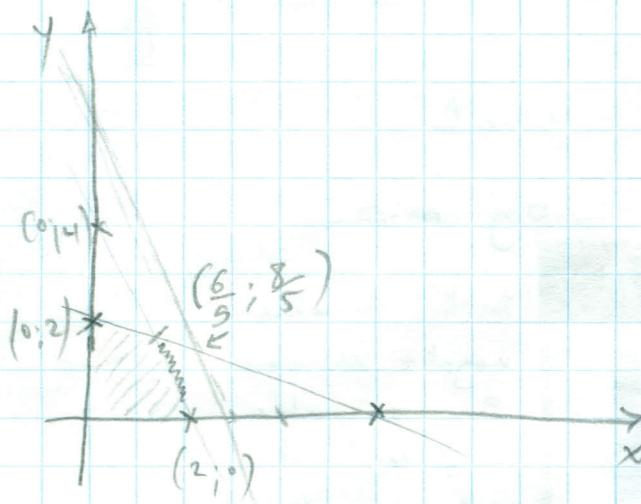
$$\text{s.t. } \begin{cases} x + 3y \leq 6 \\ 2x + y \leq 4 \\ x, y \geq 0 \end{cases}$$

$$\rightarrow \min 6\alpha + 4\beta$$

s.t

$$\begin{cases} \alpha + 2\beta \geq 2 \\ 3\alpha + \beta \geq 1 \end{cases}$$

$$\alpha, \beta \geq 0$$



$$\max 2x + y$$

s.t.

$$\begin{cases} x + 3y \leq 6 \\ 2x + y \leq 4 \end{cases}$$

$$x, y \geq 0$$

$$x + 3y = 6 \rightarrow x = 6 - 3y$$

$$2x + y = 4 \quad 12 - 6y + y = 4 \rightarrow 8 = 5y \Rightarrow y = \frac{8}{5} \text{ and } x = 6 - \frac{24}{5} = \frac{6}{5}$$

$$z(2;0) = 4$$

$$z(0,2) = 2$$

$$z\left(\frac{6}{5}, \frac{8}{5}\right) = \frac{12+8}{5} = \frac{20}{5} = 4$$

$$x^* = \frac{6}{5}; y^* = \frac{8}{5} \rightarrow z^* = 4$$

$$\min 6x + 4y$$

$$\begin{cases} x + 2y \geq 2 \\ 3x + y \geq 1 \end{cases}$$

$$\alpha, \beta \geq 0; \alpha = 0; \beta = 1$$

$$z^* = 4$$

• what happens if the rhs of const 1 is reduced to 5?

• what happens if the rhs of const 2 is increased by 2?

• what happens if the obj. coef of y is increased?

• what can you say about the dual optimal solution?