

HW#1

$$Q1. \quad S = (ab + cd) (\bar{c} + \bar{d} + \bar{e}) (efh)$$

The set of all test vectors that detect the fault h stuck-at-1 is defined by

$$\bar{h} \cdot (S_h \oplus S_{\bar{h}}) = 1$$

$$S_h = (ab + cd) (\bar{c} + \bar{d} + \bar{e}) (ef)$$

$$S_{\bar{h}} = 0$$

$$\begin{aligned} S_h \oplus S_{\bar{h}} &= (ab + cd) (\bar{c} + \bar{d} + \bar{e}) (ef) \\ &= (ab + cd) (\bar{c}ef + \bar{d}ef) \\ &= (ab\bar{c}ef + ab\bar{d}ef) \end{aligned}$$

Thus, the set of test vectors that detect the fault h stuck-at-1 are:

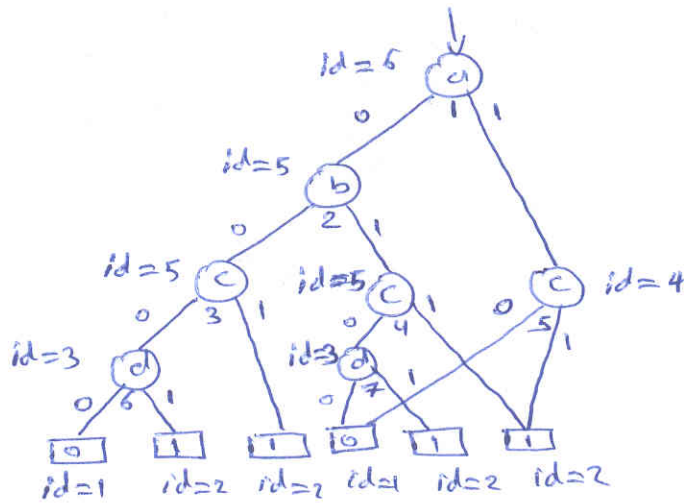
$$ab\bar{c}ef\bar{h} + ab\bar{d}ef\bar{h} = 1$$

which can be represented as:

$$\begin{aligned} \{a, b, c, d, e, f, h\} &= \{110x110, 11x0110\} \\ &= \{1100110, 1101110, 1110110\} \end{aligned}$$

Thus, there are only 3 test vectors that can detect the fault.

Q2.



First, we set $id(v) = 1$ for all leaf vertices with value 0 and $id(v) = 2$ for all leaf vertices with value 1. We initialize ROBDD with two vertices for 0 and 1. Then, we process vertices at level 4, i.e. nodes with index = d.

$$V = \{6, 7\}$$

Since none of these vertices has $id(\text{low}(v)) = id(\text{high}(v))$ none of these vertices is dropped. Keys are assigned to vertices; $key(6) = (1, 2)$, $key(7) = (1, 2)$.

$$\text{oldkey} = (0, 0).$$

We process again the set $V = \{6, 7\}$.

$$v = \{6\}$$

$$\text{next id} = 3, \quad id(6) = 3, \quad \text{oldkey} = (1, 2)$$

We add $v = \{6\}$ to the ROBDD.

$$v = \{7\}$$

$$\text{Since } key(7) = \text{oldkey}, \quad id(7) = 3$$

Next, we process vertices at level 3, i.e. nodes with index = c.

$$V = \{3, 4, 5\}$$

None of the nodes is dropped.

$$key(3) = (3, 2), \quad key(4) = (3, 2), \quad key(5) = (1, 2)$$

$$\text{oldkey} = (0, 0)$$

$v = \{5\}$
 $nextid = 4$, $oldkey = (1, 2)$, $id(5) = 4$
 $v = \{5\}$ is added to the ROBDD.

$v = \{3\}$
 $nextid = 5$, $id(3) = 5$, $oldkey = (3, 2)$
 $v = \{3\}$ is added to the ROBDD.

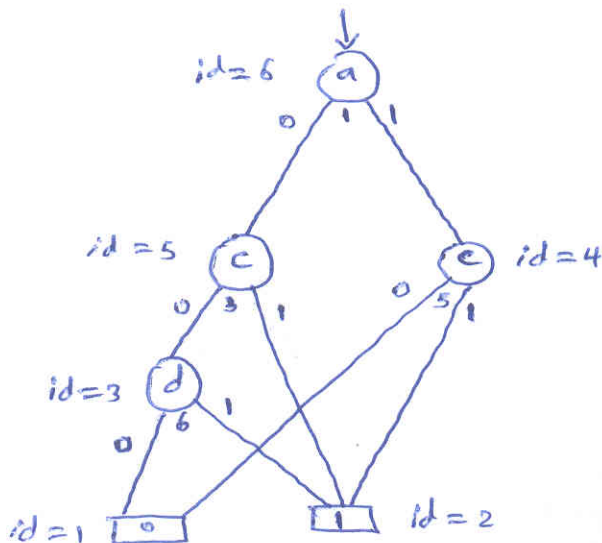
$v = \{4\}$
 Since $key(4) = oldkey$, $id(4) = 5$

Then, we process vertices at level 2, i.e., nodes with index = b.

$V = \{2\}$
 Since $id(low(2)) = id(high(2)) = 5$, $id(2) = 5$
 and the vertex is dropped from V .

Finally, we process vertices at level 1, i.e., nodes with index = a.

$V = \{1\}$
 $key(1) = (5, 4)$.
 $oldkey = (0, 0)$
 $nextid = 6$, $id(1) = 6$, $oldkey = (5, 4)$.
 $v = \{1\}$ is added to the ROBDD.
 Thus, the reduced OBDD is as follows:



$$Q_3. \quad f = a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) a_2 \bar{b}_2 + (\bar{a}_3 \oplus b_3)(\bar{a}_2 \oplus b_2) a_1 \bar{b}_1$$

(i) implementation using 2x1 MUXs.

We perform Shannon expansion on variable a_3 .

$$f = \bar{a}_3 [\bar{b}_3 a_2 \bar{b}_2 + \bar{b}_3 (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1] \\ + a_3 [\bar{b}_3 + b_3 a_2 \bar{b}_2 + b_3 (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1]$$

Then, we expand on variable b_3 .

$$f = \bar{a}_3 [\bar{b}_3 [a_2 \bar{b}_2 + (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1] \\ + b_3 [0]] \\ + a_3 [\bar{b}_3 [1] \\ + b_3 [a_2 \bar{b}_2 + (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1]]$$

Note that the cofactor $a_2 \bar{b}_2 + (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1$ is shared and is implemented once.

We next expand on a_2

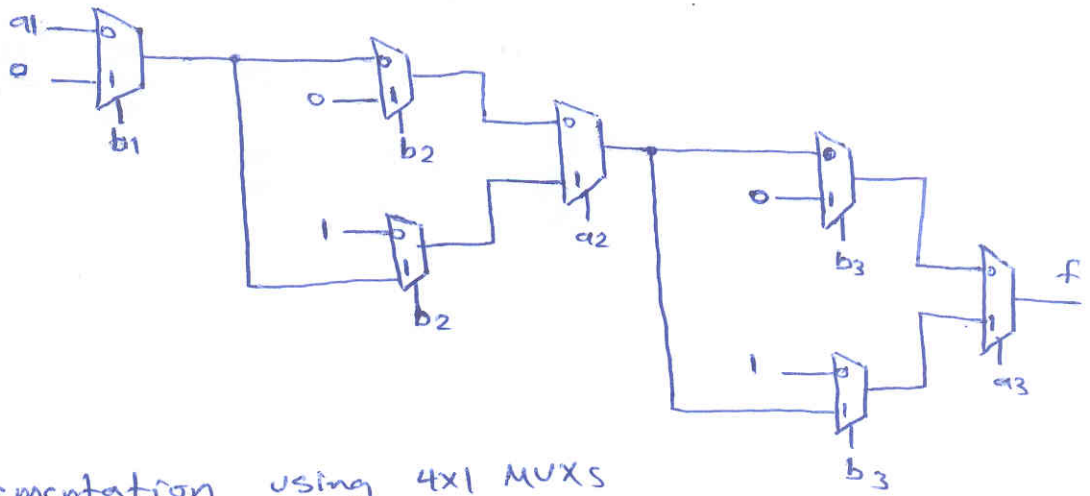
$$a_2 \bar{b}_2 + (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1 = \bar{a}_2 [\bar{b}_2 a_1 \bar{b}_1] \\ + a_2 [\bar{b}_2 + b_2 a_1 \bar{b}_1]$$

Next, we expand on b_2

$$= \bar{a}_2 [\bar{b}_2 [a_1 \bar{b}_1] + b_2 [0]] \\ + a_2 [\bar{b}_2 [1] + b_2 [a_1 \bar{b}_1]]$$

Also, $a_1 \bar{b}_1$ is shared and is implemented once. We expand on b_1 .

$$a_1 \bar{b}_1 = \bar{b}_1 [a] + b_1 [0]$$



(ii) implementation using 4x1 MUXS

We perform expansion on variables a_3 & b_3

$$f = \bar{a}_3 \bar{b}_3 [a_2 \bar{b}_2 + (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1]$$

$$+ \bar{a}_3 b_3 [0] + a_3 \bar{b}_3 [1]$$

$$+ a_3 b_3 [a_2 \bar{b}_2 + (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1]$$

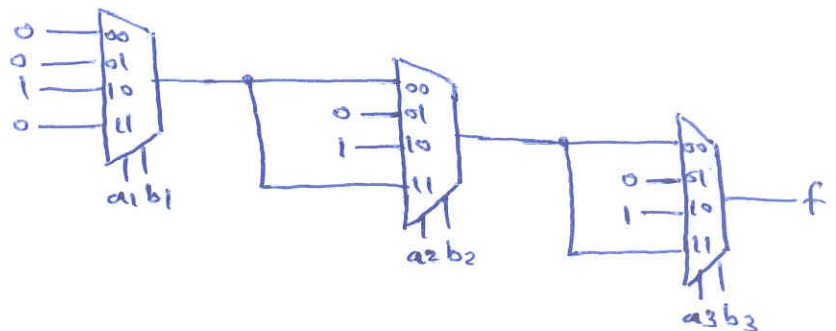
The cofactor $a_2 \bar{b}_2 + (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1$ is shared and implemented once.

We next expand on variables a_2 & b_2 .

$$a_2 \bar{b}_2 + (\bar{a}_2 \oplus b_2) a_1 \bar{b}_1 = \bar{a}_2 \bar{b}_2 [a_1 \bar{b}_1] + \bar{a}_2 b_2 [0]$$

$$+ a_2 \bar{b}_2 [1] + a_2 b_2 [a_1 \bar{b}_1]$$

The cofactor $a_1 \bar{b}_1$ is shared and implemented once.



$$\begin{aligned}
 \text{(iii)} \quad f &= \bar{a}_1 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) a_2 \bar{b}_2] \\
 &+ a_1 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) a_2 \bar{b}_2 + (\bar{a}_3 \oplus b_3)(\bar{a}_2 \oplus b_2) \bar{b}_1] \\
 &= \bar{a}_1 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) a_2 \bar{b}_2] \\
 &+ a_1 [\bar{b}_1 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) a_2 \bar{b}_2 + (\bar{a}_3 \oplus b_3)(\bar{a}_2 \oplus b_2)] \\
 &\quad + b_1 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) a_2 \bar{b}_2]] \\
 &= \bar{a}_1 [\bar{a}_2 [a_3 \bar{b}_3] + a_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) \bar{b}_2]] \\
 &+ a_1 [\bar{b}_1 [\bar{a}_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) \bar{b}_2] \\
 &\quad + a_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) \bar{b}_2 + (\bar{a}_3 \oplus b_3) b_2]] \\
 &\quad + b_1 [\bar{a}_2 [a_3 \bar{b}_3] + a_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) \bar{b}_2]]] \\
 &= \bar{a}_1 [\bar{a}_2 [a_3 \bar{b}_3] + a_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) \bar{b}_2]] \\
 &+ a_1 [\bar{b}_1 [\bar{a}_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) \bar{b}_2] \\
 &\quad + a_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3)]] \\
 &\quad + b_1 [\bar{a}_2 [a_3 \bar{b}_3] + a_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) \bar{b}_2]]]]
 \end{aligned}$$

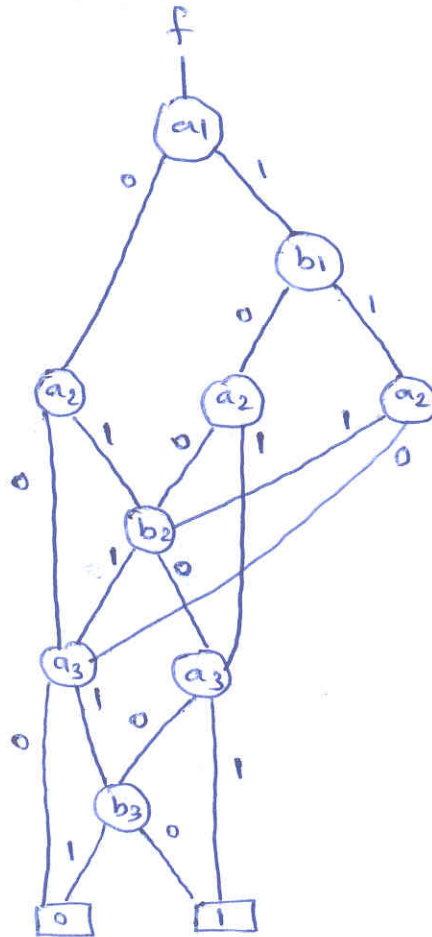
The common cofactors will be shared.

$$a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) \bar{b}_2 = \bar{b}_2 [a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3)] + b_2 [a_3 \bar{b}_3]$$

$$a_3 \bar{b}_3 + (\bar{a}_3 \oplus b_3) = \bar{a}_3 [\bar{b}_3] + a_3 [1]$$

$$a_3 \bar{b}_3 = \bar{a}_3 [0] + a_3 [\bar{b}_3]$$

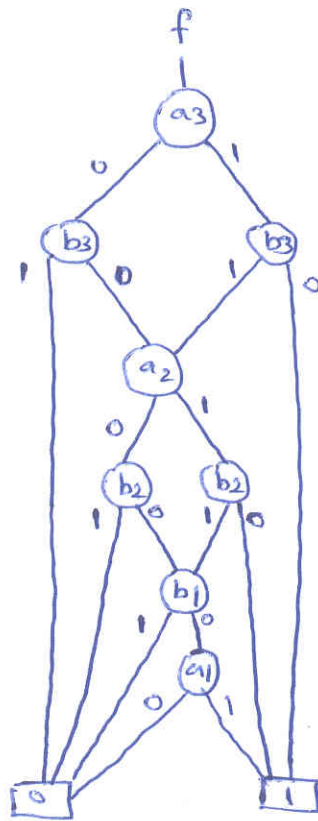
So, the ROBDD using the variable ordering $\{a_1, b_1, a_2, b_2, a_3, b_3\}$ is :



This variable ordering has resulted in an ROBDD with 9 vertices.

Based on the variable ordering $\{a_3, b_3, a_2, b_2, a_1, b_1\}$ is used, we will get an ROBDD with 8 vertices as can be seen from the solution in (1).

The ROBDD based on the variable ordering $\{a_3, b_3, a_2, b_2, a_1, b_1\}$ is shown below:



This ROBDD has 8 vertices, so, this is a better variable ordering,

$$\text{Q4. } f = a \oplus b \oplus c$$

$$g = ac + \bar{a}\bar{b} + b\bar{c}$$

(i) We first expand both functions on the variables a & b and then we perform the required operations

$$f = \bar{a}\bar{b}[c] + \bar{a}b[\bar{c}] + a\bar{b}[\bar{c}] + ab[c]$$

$$g = \bar{a}\bar{b}[1] + \bar{a}b[\bar{c}] + a\bar{b}[c] + ab[1]$$

$$f \cdot g = \bar{a}\bar{b}[c] + \bar{a}b[\bar{c}] + a\bar{b}[0] + ab[c]$$

$$f + g = \bar{a}\bar{b}[1] + \bar{a}b[\bar{c}] + a\bar{b}[1] + ab[1]$$

$$= \bar{a}\bar{b} + \bar{a}b\bar{c} + a\bar{b} + ab$$

$$= \bar{a}\bar{b} + \bar{a}b\bar{c} + a$$

$$= \bar{b} + b\bar{c} + a$$

$$= \bar{b} + \bar{c} + a$$

$$f \oplus g = \bar{a}\bar{b}[\bar{c}] + \bar{a}b[0] + a\bar{b}[1] + ab[\bar{c}]$$

$$= \bar{a}\bar{b}\bar{c} + a\bar{b} + ab\bar{c}$$

$$= \bar{b}\bar{c} + a\bar{b} + a\bar{c}$$

(ii) ITE DAG for $f \oplus g$

$$f \oplus g = \text{ITE}(f, \bar{g}, g)$$

$$\bar{g} = \bar{a}\bar{b}[0] + \bar{a}b[c] + a\bar{b}[\bar{c}] + ab[0]$$

$$= \bar{a}b\bar{c} + a\bar{b}\bar{c}$$

$$f \oplus g = \text{ITE}(a \oplus b \oplus c, \bar{a}b\bar{c} + a\bar{b}\bar{c}, ac + \bar{a}\bar{b} + b\bar{c})$$

We assume the variable ordering $\{a, b, c\}$.

- $x = a$

$$t = \text{ITE}(\bar{b} \oplus c, \bar{b}\bar{c}, c + b\bar{c})$$

$$e = \text{ITE}(b \oplus c, bc, \bar{b} + b\bar{c})$$

- $\text{ITE}(\bar{b} \oplus c, \bar{b}\bar{c}, c + b\bar{c})$

- $x = b$

$$t = \text{ITE}(c, 0, 1) \Rightarrow \text{trivial case} = \bar{c}$$

$$\Rightarrow t = 3$$

$$e = \text{ITE}(\bar{c}, \bar{c}, c) \Rightarrow \text{trivial case} = 1$$

$$\Rightarrow e = 2$$

since $t \neq e$, identifier 4 will be returned for $\text{ITE}(\bar{b} \oplus c, \bar{b}\bar{c}, c + b\bar{c})$

- $\text{ITE}(b \oplus c, bc, \bar{b} + b\bar{c})$

- $x = b$

$$t = \text{ITE}(\bar{c}, c, \bar{c}) \Rightarrow \text{trivial case} = 0$$

$$\Rightarrow t = 1$$

$$e = \text{ITE}(c, 0, 1) \Rightarrow \text{trivial case} = \bar{c}$$

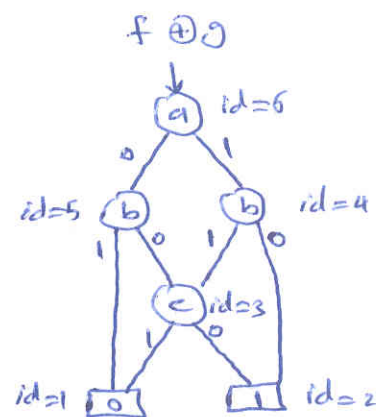
$$\Rightarrow e = 3$$

since $t \neq e$, identifier 5 will be assigned to $\text{ITE}(b \oplus c, bc, \bar{b} + b\bar{c})$.

since $t \neq e$, identifier 6 will be assigned.

The constructed unique table and the corresponding ITE DAG is shown below:

id	var	left child	right child
3	c	2	1
4	b	2	3
5	b	3	1
6	a	5	4



Q5. Minimum cover using Exact-Cover procedure.

$$A = \begin{matrix} & c_1 & c_2 & c_3 & c_4 & c_5 & c_6 & c_7 \\ r_1 & \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

First, we check if there are essential columns. We can see that c_7 is essential and hence it is selected and rows r_5 , r_7 , and r_{10} are removed.

Thus, $x = [0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1]$ and the reduced matrix is:

$$A = \begin{matrix} & c_1 & c_2 & c_3 & c_4 & c_5 & c_6 \\ r_1 & \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 & 0 \end{bmatrix} \end{matrix}$$

Next, we check row dominance and column dominance.

- c_1 dominates c_6 , so c_6 is removed
- c_4 dominates c_5 , so c_5 is removed
- r_8 dominates r_1 and r_3 , so r_8 is removed
- r_9 dominates r_4 , so r_9 is removed

Thus, the reduced matrix is :

$$A = \begin{matrix} & c_1 & c_2 & c_3 & c_4 \\ r_1 & \left[\begin{array}{cccc} 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \end{array} \right] \end{matrix}$$

The matrix cannot be reduced further.

Let us assume a branching column c_1 .

If we select c_1 , then $x = [1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1]$, and

the matrix \tilde{A} becomes

$$\tilde{A} = \begin{matrix} & c_2 & c_3 & c_4 \\ r_4 & \left[\begin{array}{ccc} 0 & 1 & 1 \\ 1 & 0 & 1 \end{array} \right] \\ r_5 & \end{matrix}$$

- c_4 dominates c_2 , so c_2 is removed
- c_4 dominates c_3 , so c_3 is removed

Then, c_4 becomes essential and is selected

$$x = [1 \ 0 \ 0 \ 1 \ 0 \ 0 \ 1]$$

$$b = [1 \ 0 \ 0 \ 1 \ 0 \ 0 \ 1]$$

if c_1 is not selected, then $x = [0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1]$

and the matrix \tilde{A} becomes

$$\tilde{A} = \begin{matrix} & c_2 & c_3 & c_4 \\ r_1 & \left[\begin{array}{ccc} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{array} \right] \end{matrix}$$

Then, c_2, c_3 and c_4 become essential columns and are selected.

$$\text{Thus, } x = [0 \ 1 \ 1 \ 1 \ 0 \ 0 \ 1]$$

Since the cost of this solution is larger than the best solution so far, the returned best solution is $x = [1 \ 0 \ 0 \ 1 \ 0 \ 0 \ 1]$.