COE 405 Design & Modeling of Digital Systems

Course Project – Term 181 Design and Modeling of a RISC Processor Due on: Wednesday Dec. 12, 2018

In this project, you will design, model in Verilog and synthesize a 16-bit MIPS-like processor. The details about the processor are given below.

Instruction Set Architecture

The 16-bit MIPS-like processor has seven 16-bit general-purpose registers: R1 through R7. R0 is hardwired to zero and cannot be written. There is also one special-purpose 12-bit register, which is the program counter (PC). All instructions are 16 bits and there are three instruction formats: R-type, I-type, and J-type as shown below:

R-type format

4-bit opcode (Op), 3-bit register numbers (Rs, Rt, and Rd), and 3-bit function field (funct)

Op ⁴ Rs ³ Rt ³ Rd ³ funct ³
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I-type format

4-bit opcode (Op), 3-bit register numbers (Rs and Rt), and 6-bit immediate constant

Op^4	Rs ³	Rt ³	Immediate ⁶
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J-type format

4-bit opcode (Op) and 12-bit immediate constant

Op ⁴ Immediate ¹²

For R-type instructions, Rs and Rt specify the two source register numbers, and Rd specifies the destination register number. The function field can specify at most eight functions for a given opcode. Opcodes 0 and 1 are reserved for R-type instructions.

For I-type instructions, Rs specifies a source register number, and Rt can be a second source or a destination register number. The immediate constant is only 6 bits because of the fixed-size nature of the instruction. The 6-bit immediate constant is assumed to be sign-extended for all instructions except the logical instructions (i.e., ANDI, ORI).

For J-type, a 12-bit immediate constant is used for J (jump), JAL (jump-and-link), and LUI (load upper immediate) instructions.

Instruction Encoding

Sixteen R-type instructions, eleven I-type instructions, and three J-type instructions are defined. These instructions, their meaning, and their encoding are shown below:

Instr	Meaning	Encoding						
AND	Reg(Rd) = Reg(Rs) & Reg(Rt)	Op = 0000	Rs	Rt	Rd	f = 000		
OR	Reg(Rd) = Reg(Rs) Reg(Rt)	Op = 0000	Rs	Rt	Rd	f = 001		
NOR	$\operatorname{Reg}(\operatorname{Rd}) = \sim (\operatorname{Reg}(\operatorname{Rs}) \operatorname{Reg}(\operatorname{Rt}))$	Op = 0000	Rs	Rt	Rd	f = 010		
XOR	$Reg(Rd) = Reg(Rs) \wedge Reg(Rt)$	Op = 0000	Rs	Rt	Rd	f = 011		
SLL	$Reg(Rd) = Reg(Rs) \ll Reg(Rt)$	Op = 0000	Rs	Rt	Rd	f = 100		
SRL	Reg(Rd) = Reg(Rs) zero >> Reg(Rt)	Op = 0000	Rs	Rt	Rd	f = 101		
SRA	Reg(Rd) = Reg(Rs) sign >> Reg(Rt)	Op = 0000	Rs	Rt	Rd	f = 110		
ROL	$Reg(Rd) = Reg(Rs)$ rotate $\leq Reg(Rt)$	Op = 0000	Rs	Rt	Rd	f = 111		
ADD	$\operatorname{Reg}(\operatorname{Rd}) = \operatorname{Reg}(\operatorname{Rs}) + \operatorname{Reg}(\operatorname{Rt})$	Op = 0001	Rs	Rt	Rd	f = 000		
SUB	$\operatorname{Reg}(\operatorname{Rd}) = \operatorname{Reg}(\operatorname{Rs}) - \operatorname{Reg}(\operatorname{Rt})$	Op = 0001	Rs	Rt	Rd	f = 001		
SLT	Reg(Rd) = Reg(Rs) signed < Reg(Rt)	Op = 0001	Rs	Rt	Rd	f = 010		
SLTU	Reg(Rd) = Reg(Rs) unsigned $< Reg(Rt)$	Op = 0001	Rs	Rt	Rd	f = 011		
DIV	Reg(Rd) = Quot(Reg(Rs) / Reg(Rt))	Op = 0001	Rs	Rt	Rd	f = 100		
REM	$\operatorname{Reg}(\operatorname{Rd}) = \operatorname{Rem}(\operatorname{Reg}(\operatorname{Rs}) / \operatorname{Reg}(\operatorname{Rt}))$	Op = 0001	Rs	Rt	Rd	f = 101		
MUL	$\operatorname{Reg}(\operatorname{Rd}) = \operatorname{Reg}(\operatorname{Rs}) * \operatorname{Reg}(\operatorname{Rt})$	Op = 0001	Rs	Rt	Rd	f = 110		
JR	PC = lower 12 bits of Reg(Rs)	Op = 0001	Rs	000	000	f = 111		
LW	$\operatorname{Reg}(\operatorname{Rt}) = \operatorname{Mem}(\operatorname{Reg}(\operatorname{Rs}) + \operatorname{ext}(\operatorname{im}^6))$	Op = 0010	Rs	Rt	Imn	nediate ⁶		
SW	$Mem(Reg(Rs) + ext(im^6)) = Reg(Rt)$	Op = 0011	Rs	Rt	Immediate ⁶			
ANDI	$\operatorname{Reg}(\operatorname{Rt}) = \operatorname{Reg}(\operatorname{Rs}) \& \operatorname{ext}(\operatorname{im}_6)$	Op = 0110	Rs	Rt	Immediate ⁶			
ORI	$\operatorname{Reg}(\operatorname{Rt}) = \operatorname{Reg}(\operatorname{Rs}) \operatorname{ext}(\operatorname{im}_6)$	Op = 0111	Rs	Rt	Immediate ⁶			
ADDI	$\operatorname{Reg}(\operatorname{Rt}) = \operatorname{Reg}(\operatorname{Rs}) + \operatorname{ext}(\operatorname{im}^6)$	Op = 1000	Rs	Rt	Immediate ⁶			
BEQ	Branch if $(\text{Reg}(\text{Rs}) == \text{Reg}(\text{Rt}))$	Op = 0100	Rs	Rt	Immediate ⁶			
BNE	Branch if (Reg(Rs) != Reg(Rt))	Op = 0101	Rs	Rt	Immediate ⁶			
BLTZ	Branch if $(\text{Reg}(\text{Rs}) < 0)$	Op = 1100	Rs	Rt	Immediate ⁶			
BLEZ	Branch if $(\text{Reg}(\text{Rs}) \le 0)$	Op = 1101	Rs	Rt	Immediate ⁶			
BGTZ	Branch if $(\text{Reg}(\text{Rs}) > 0)$	Op = 1110	Rs	Rt	Immediate ⁶			
BGEZ	Branch if $(\text{Reg}(\text{Rs}) \ge 0)$	Op = 1111	Rs	Rt	Immediate ⁶			
J	$PC = Immediate^{12}$	Op = 1001	Immediate ¹²					
JAL	$R7 = PC + 1$, $PC = Immediate^{12}$	Op = 1011						
LUI	$R1 = Immediate^{12} \ll 4$	Op = 1010		Imn	nediate	Immediate ¹²		

There are three shift and one rotate instructions. For shift and rotate instructions, the least significant 4 bits of register Rt are used as the shift/rotate amount. There is only one rotate left (ROL) instruction. To rotate right by *n* bits, you can rotate left by 16 - n bits, because registers are 16 bits. The Load Upper Immediate (LUI) is of the J-type to have a 12-bit immediate constant loaded into the upper 12 bits of register R1. The LUI can be combined with ORI (or ADDI) to load any 16-bit constant into a register. Although the instruction set is reduced, it is still rich enough to write useful programs. We can have procedure calls and returns using the JAL and JR instructions.

Memory

Your processor will have separate instruction and data memories with $2^{12} = 4096$ words each. Each word is 16 bits or 2 bytes. Memory is *word addressable*. Only words (not bytes) can be read and written to memory, and each address is a word address. This will simplify the processor implementation. The PC contains a word address (not a byte address). Therefore, it is sufficient to increment the PC by 1 (rather than 2) to point to the next instruction in memory. Also, the Load and Store instructions can only load and store words. There is no instruction to load or store a byte in memory.

Addressing Modes

For branch instructions (BEQ, BNE, BLTZ, BLEZ, BGTZ and BGEZ), PC-relative addressing mode is used: $PC = PC + sign-extend(immediate^6)$. For jump instructions (J and JAL), direct addressing is used: $PC = Immediate^{12}$. For LW and SW instructions, base-displacement addressing mode is used. The base address in register Rs is added to the sign-extended immediate⁶ to compute the memory address.

Program Execution

The program will be loaded and will start at address 0 in the instruction memory. The data segment will be loaded and will start also at address 0 in the data memory. To terminate the execution of a program, the last instruction in the program can jump or branch to itself indefinitely.

Calculator Application

You need to develop a calculator application for performing addition, subtraction, multiplication and division of any entered two integer 16-bit numbers. You need to interface with the LCD screen and keypad to display a message on the LCD screen asking the user to enter the first number and then read the input from the keypad. Then, display the entered number in decimal. After that, display a message asking the user to enter the second number, read it and display it in decimal. Finally, ask the user to enter the requested operation, read the request, perform the operation and display the result on the LCD screen in decimal.

Project Report

The report document must contain sections highlighting the following:

1 – Design and Implementation

- Specify clearly the design giving detailed description of the datapath, its components, control, and the implementation details (highlighting the design choices you made and why, and any notable features that your processor might have.) Document clearly design alternatives explored and why a given design is selected.
- Provide drawings of the component circuits and the overall datapath.
- Provide a complete description of the control logic and the control signals. Provide a table giving the control signal values for each instruction.
- Provide a complete description of the forwarding logic, the cases that were handled, and the cases that stall the pipeline
- Use a hierarchical Verilog modeling style when modeling your processor. Your CPU should be composed of a control unit and datapath. The datapath should be composed of ALU, Register file, Instruction Memory, Data Memory and other necessary components.
- Provide list of sources for any parts of your design that are not entirely yours (if any).
- Carry out the design and implementation with the following aspects in mind:
 - Consider alternative design solutions and justify your design selection

- Correctness of the individual components
- Correctness of the overall design when wiring the components together
- Completeness: all instructions were implemented properly, detecting dependences and forwarding was handled properly, and stalling the pipeline was handled properly for all cases.

2 – Simulation and Testing

- Carry out the simulation of the processor developed using Modelsim or Isim.
- Test each of the components individually and demonstrate its correct operation including the ALU and register file.
- Describe the test programs that you used to test your design with enough comments describing the program, its inputs, and its expected output. List all the instructions that were tested and work correctly. List all the instructions that do not run properly.
- Also provide snapshots of the Simulator window with your test program loaded and showing the simulation output results.
- Synthesize the processor on FPGA and demonstrate its correct functionality by correct implementation of the calculator application.

3 – Teamwork

- This project is a team work project with **up to four** students per team. Make sure to write the names of all the group members on the project report title page.
- Each group should assign a group leader that leads the conduction of the project, divide the project tasks among the team members.
- Project tasks should be divided among the group members so that each group member contributes equally in the project and everyone is involved in all the following activities:
 - Design and Implementation
 - Simulation and Testing
 - Synthesis and FPGA implementation
 - Interfacing with keypad
 - Interfacing with LCD screen
 - Design and results reporting
- Come up with a project plan detailing the tasks to be performed in the project, their planned start and finish dates and the team member primarily responsible for performing the task. Submit the project plan by **Thursday, Oct. 25**.
- Clearly show in the report the work done by each group member, and how the work deviated from the proposed plan.
- Each group member will be evaluated based on his contribution in the project. Thus, it is expected that each group member could have a different mark in the project.
- Students who <u>help</u> other team members should mention that to earn credit for that.

Submission Guidelines

All submissions will be done through WebCT.

Attach one zip file containing all Verilog files used in your design, a video demo of your project, as well as the report document.

Submit also a hard copy of the report during the class lecture.

Students are expected to report project progress every two weeks.

The deadlines for submissions are given below:

Task	Deadline
Formation of project groups	Oct. 18, 2018
Submission of project plan	Oct. 25, 2018
Submission of Progress Report	Nov. 24, 2018
Submission of Final Report	Dec. 12, 2018

Grading policy

Grading Criteria	Mark
Demonstration of correct functionality of	60
components and whole processor design by	
simulation	
Demonstration of correct functionality of	25
Calculator Application on FPGA	
Project Documentation and Report	15
Organization	
Total	100

The project will be evaluated based on the final report, project demonstration and oral project evaluation with all team members.