## COMPUTER ENGINEERING DEPARTMENT

COE 205

## COMPUTER ORGANIZATION \& ASSEMBLY PROGRAMMING

## Major Exam I

Second Semester (092)
Time: 8:00-10:00 PM

Student Name : _KEY
Student ID. $\qquad$

| Question | Max Points | Score |
| :---: | :---: | :---: |
| Q1 | $\mathbf{7 6}$ |  |
| Q2 | $\mathbf{1 2}$ |  |
| Q3 | $\mathbf{1 2}$ |  |
| Total | $\mathbf{1 2 0}$ |  |

(Q1) Fill the blank in each of the following:
(1) The advantages and disadvantages of programming in assembly language are Advantages: accessibility to system hardware and space and time efficiency.
Disadvantages: Programs are not portable and program development and maintenance is more difficult than using high level language.
(2) The advantages and disadvantages of DRAM are

Advantages: smaller in size and hence denser and cheaper. Disadvantages: slower as it needs refreshing.
(3) The instruction set architecture of a processor is composed of the instruction set, programmer-accessible registers, and memory.
(4) In protected mode, the logical address consists of a 16-bit segment selector and a 32-bit offset.
(5) In protected mode, the linear address is computed based on adding a 32 -bit offset with a 32 -bit base address obtained from descriptor table indexed by segment selector.
(6) Given that variables I and J are Dword variables, their content can be swapped using the following instructions:
mov eax, I
xchg eax, J
mov I, eax
(7) Given that variable I is defined as Qword, the content of I is incremented using the following code:
add Dword PTR I, 1
adc Dword PTR I+4, 0
(8) Given a magnetic disk with the following properties: Rotation Speed $=5000$ RPM (rotations per minute), Average Seek $=4 \mathrm{~ms}$, Sector $=512$ bytes, Track $=250$ sectors. The average time to access a block of 50 consecutive sectors is Average access time $=$ Seek Time + Rotation Latency + Transfer Time Rotations per second $=5000 / 60=83.33$ RPS Rotation time in milliseconds $=1000 / 83.33=12 \mathrm{~ms}$ Time to transfer 50 sectors $=(50 / 250) * 12=2.4 \mathrm{~ms}$ Average access time $=4+6+2.4=12.4 \mathrm{~ms}$.
(9) The integer number -2000 is represented in hexadecimal using 16-bit 2's complement representation as F830.
(10) Assuming 16-bit 2's complement representation, the hexadecimal number EC10 represents the decimal number $\underline{-5104}$.
(11) Assuming 4-bit 2's complement representation, the largest number that can be stored is $\underline{+7}$ in decimal and 0111 in binary and the smallest number that can be stored is $\underline{8}$ in decimal and $\underline{1000}$ in binary.
(12) Given that the number 88 h is represented using 8 -bit 2 's complement representation, the equivalent number represented using 16-bit 2's complement representation is FF88.
(13) Given that register $\mathrm{AL}=\mathrm{C} 4$ stores an ASCII character, then the stored character is $\underline{D}$ and the used parity is odd. Note that ' $A$ ' $=41 \mathrm{~h}$ and ' $a$ ' $=61 \mathrm{~h}$.
(14) The EIP register holds the address of the next instruction to be fetched from memory.
(15) Given that the instruction ADD AX, I (having the machine code 03060000) is stored at address 0000002 C , then the address of the next instruction to be fetched from memory is $0000002 \mathrm{C}+4=00000030$.
(16) Given a processor with an 8 -stage pipeline and clock frequency of 2 GHZ , the time that will be required to execute a program of 4 billion instructions assuming that there will be no pipeline stalls is nearly $4 \times 10^{9} /\left(2 \times 10^{9}\right)=2$ seconds.
(17) Assume that the range of addresses from 00000 to 00 A 1 A is used by another program. Given that a program has a code segment of 8 Kbyte and a data segment of 3 Kbyte , the code segment number allocated is 00 A 2 and the data segment number allocated is 02 A 2 .

The used address space for the code segment is $00 \mathrm{~A} 20+1 \mathrm{FFF}=2 \mathrm{~A} 1 \mathrm{~F}$. Thus, the next available address is 2A20. Thus, the next available segment is 2 A 2 .
(18) Assume that $\mathrm{DS}=00 \mathrm{EF}, \mathrm{CS}=013 \mathrm{~A}, \mathrm{ES}=0112$, $\mathrm{SS}=0 \mathrm{FEC}, \mathrm{IP}=00 \mathrm{FF}, \mathrm{BX}=309 \mathrm{~A}$, and $\mathrm{SP}=01 \mathrm{FC}$. Based on 16 -bit real-mode addressing, the linear address of the next instruction to be fetched from memory is $013 \mathrm{~A} 0+000 \mathrm{FF}=0149 \mathrm{~F}$.
(19) Assume that $\mathrm{DS}=00 \mathrm{EF}, \mathrm{CS}=013 \mathrm{~A}, \mathrm{ES}=0112$, $\mathrm{SS}=0 \mathrm{FEC}, \mathrm{IP}=00 \mathrm{FF}, \mathrm{BX}=309 \mathrm{~A}$, and $\mathrm{SP}=01 \mathrm{FC}$. Based on 16 -bit real addressing mode, the linear address of the source operand in the instruction MOV AX, [BX+5] is $00 \mathrm{EF} 0+0309 \mathrm{~A}+00005=03 \mathrm{~F} 8 \mathrm{~F}$.
(20) The addressing mode of the source operand in the instruction MOV EAX, [MSG+1] is direct addressing mode.
(21) Assume that $\mathrm{AX}=00 \mathrm{FFh}$. Executing the instruction INC AL produces the result $\mathrm{AX}=\underline{0000}$.
(22) The addressing mode of the source operand in the instruction MOV SI, [EBX] is register indirect addressing mode.
(23) The assembler allocates $\underline{4 *}^{*}\left(5^{*}(1+30)\right)=620$ bytes for the variable Array defined below:

Array Dword $5 \operatorname{dup}(30,30 \operatorname{dup}(0))$
(24) The content of register EAX after executing the following instructions is_$\underline{5+2+10=7}$.

## I=5

J EQU 10
MOV EAX, I-J
I=I-3
ADD EAX, I+J
(25) Assuming the following data segment and assuming that variable X is given the linear address 00404000 h , then the linear address for variables Y and Z will be $\underline{00404008 \mathrm{~h}}$ and $\underline{00404010 \mathrm{~h}}$.
.DATA
X BYTE 1, 2, 3, 4, 5
ALIGN 4
Y DWORD 4, 5
ALIGN 2
Z WORD 7, 8,9
(26) Assuming the following data segment and assuming that variable X is given the linear address 00404000 h , then the content of register EAX after executing the instruction MOV EAX, Y-5 is $\underline{14000 A 00}$.

```
.DATA
X BYTE "EXAM I",0
    WORD 10,20
Y DWORD 30,40
```

(27) Assuming the following data segment and assuming that variable X is given the linear address 00404000 h , after executing the code given below, the content of register $E A X=\underline{4}$ and $E B X=\underline{00404004 h}$.

```
.DATA
X WORD 10, 20,30
Y DWORD 30, 40,50
.CODE
MOV EAX, TYPE Y
MOV EBX, OFFSET Y-2
```

(28) After executing the code given below, the content of registers EAX and EBX will be $\underline{00000006}$ and $\underline{24 d=00000018 \mathrm{~h}}$.
.DATA

| ARRAY DWORD | $10,20,30$, |
| :--- | :--- |
| $40,50,60$ |  |

.CODE
MOV EAX, LENGTHOF ARRAY
MOV EBX, SIZEOF ARRAY
(29) After executing the code given below, the content of register EAX will be 28001E00.
.DATA
ARRAY WORD 10, 20, 30, 40, 50, 60
CODE
MOV EAX, DWORD PTR ARRAY+3
(30) Assuming that variable ARRAY is defined as shown below:

ARRAY DWORD 1, 2, 3, 4, 5, 6
The content of register AX after executing the instruction MOV EAX, ARRAY+2 will be 0000 .
(31) Assume that AL=93h. Executing the instruction MOVSX EBX, AL produces the result EBX=FFFFFF93.
(32) Assume that $\mathrm{AX}=\mathrm{A} 100 \mathrm{~h}$. Executing the instruction $N E G A X$ produces the following results: $\mathrm{AX}=\underline{5 \mathrm{~F} 00}$ overflow flag $=\underline{0}$, sign flag $=\underline{0}$, zero flag=0, carry flag $=\underline{1}$, auxiliary flag $=\underline{0}$ and parity flag $=\underline{1}$.
(33) Assume that $\mathrm{AX}=\mathrm{ABCDh}$ and $\mathrm{BX}=8876$. Executing the instruction $A D D A X$, $B X$ produces the following results: $\mathrm{AX}=\underline{3443}$, overflow flag $=\underline{1}$, sign flag $=\underline{0}$, zero flag $=\underline{0}$, carry flag $=\underline{1}$, auxiliary flag= $\underline{1}$ and parity flag= $\underline{0}$.
(34) Assume that $\mathrm{AX}=98 \mathrm{~A} 0 \mathrm{~h}$ and $\mathrm{BX}=$ FFDAh. Executing the instruction $S U B$ $A X, B X$ produces the following results: $\mathrm{AX}=\underline{98 \mathrm{C} 6}$, overflow flag= $\underline{0}$, sign flag $=\underline{1}$, zero flag $=\underline{0}$, carry flag $=\underline{1}$, auxiliary flag $=\underline{1}$ and parity flag $=\underline{1}$.
(35) The content of register EAX after executing the instructions below will be 0000000 C .
.DATA
ARRAY DWORD 1, 2, 3, 4
DWORD 5, 6, 7, 8
DWORD 9, 10, 11, 12
RS EQU SIZEOF ARRAY
.CODE
MOV ESI, 2*RS
MOV EDI, 3
MOV EAX, ARRAY[ESI+EDI*TYPE ARRAY]
(Q2) Consider a program that has the following data segment assuming a flat memory model:

| $X$ | $E Q U$ | 16 |
| :--- | :--- | :--- |
| $Y$ | BYTE | 17 |
| $Z$ | WORD | 18 |
| $W$ | DWORD | 19 |

Indicate whether the following are valid IA-32 instructions or not. If invalid, give the reason:

1. MOV EAX, W-1

Valid.
2. MOV Z, Word PTR Y

Invalid. Both operands are memory operands.
3. MOV DS, X

Invalid. We can't move a constant to a segment register.
4. MOV Z, X

Valid.
5. MOV AX, OFFSET Z

Invalid. There is size mismatch because the address is 32 -bit.
6. MOVSX EAX, X

Invalid. The source operand can't be a constant in a MOVSX instruction.

## 7. MOV W, Dword PTR AX

Invalid. PTR operator can't be used with a register.
8. INC [EBX]

Invalid. There is ambiguity in the size of the operand to be incremented.
(Q3) Suppose that the following directives are declared in the data segment with a starting linear address of 00404000. Show the linear addresses of allocated memory and their corresponding content in hexadecimal. Note that the ASCII code for character ' $a$ ' is 61 h and that of character ' $A$ ' is 41 h . The ASCII code of character ' 0 ' is 30 h .

| I | BYTE | $10, " 10 ", 0$ |
| :--- | :--- | :--- |
|  | WORD | $10,-10$ |
| J | DWORD | $112,-112$ |
| K | EQU | 100 |

L BYTE $\mathrm{K}+20$
BYTE $\quad 2,2 \operatorname{dup}(1,-1)$

| Variable | Linear Address (Hex.) | Content (Hex.) |
| :---: | :---: | :---: |
| I | 00404000 | 0 A |
|  | 00404001 | 31 |
|  | 00404002 | 30 |
|  | 00404003 | 00 |
|  | 00404004 | 0 A |
|  | 00404005 | 00 |
|  | 00404006 | F6 |
|  | 00404007 | FF |
|  | 00404008 | 70 |
| J | 00404009 | 00 |
|  | 0040400 A | 00 |
|  | 0040400 B | 00 |
|  | 0040400 C | 90 |
|  | 0040400 D | FF |
|  | 0040400 E | FF |
|  | 0040400 F | FF |
|  | 00404010 | 78 |
|  | 00404011 | 02 |
|  | 00404012 | 01 |
|  | 00404013 | FF |
|  | 00404014 | 01 |
|  | 00404015 | FF |
|  |  |  |
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