

Section 14.2 and 14.4





Lecture outline

- Multidimensional Data
- Grid Files
- Bitmap Indexes
- Summary
- Reading List

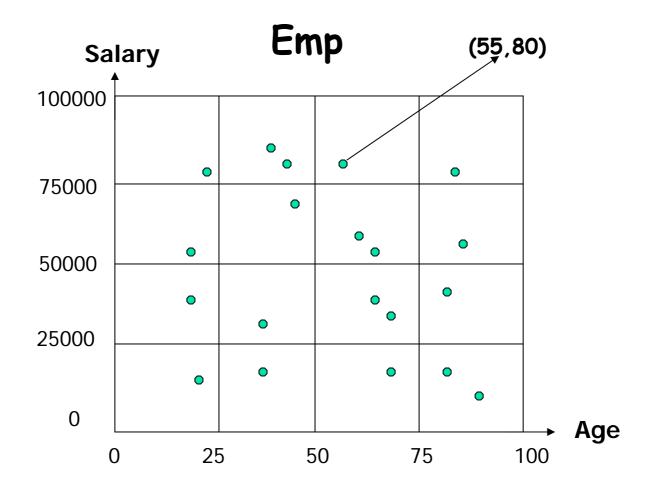


- Multidimensional Data

- Visualizing data as stored in n dimensional space
- Support queries that are not common in SQL
- Applications:
 - Geographical information system
 - Data warehouse
 - Data cubes
 - Multidimensional queries
- Examples:
 - Grid file
 - BANG file
 - BV-Trees



- Grid files





INSERT INTO Emp VALUES(60,70000);





-- Exact Match Query





-- Partial Match Query







-- Range Query





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- Bitmap Indexes

- What is Bitmap Index
- Motivation for Bitmap Indexes
- Compressed Bitmaps
- Managing Bitmap Indexes
- Summary
- Reading List



-- What is Bitmap Index

- A bitmap index for attribute A of relation R is:
 - A collection of bit-vectors
 - The number of bit-vectors = the number of distinct values of A in R.
 - The length of each bit-vector = the cardinality of R.
 - The bit-vector for value v has 1 in position i, if the ith record has v in attribute A, and it has 0 there if not.
- Records are allocated permanent numbers
- There is a mapping between record numbers and record addresses.



-- Example

- Assume relation R with
 - 2 attributes A and B.
 - Attribute A is of type Integer and B is of type String.
 - 6 records, numbered 1 through 6 as follows:

	<u>A,</u>	B
1.	30,	foo
2.	30,	bar
3.	40,	baz
4.	50,	foo
5.	40,	bar
6.	30,	baz

A bit map for attribute B is:

value	Vector
foo	100100
bar	010010
baz	001001



-- Motivation for Bitmap Indexes

- Very efficient when used for partial match queries.
- The offer the advantage of buckets
 - Where we find tuples with specified in several attributes without first retrieving all the record that matched in each of the attributes.
- They can also help answer range queries



-- Compressed Bitmaps

- Assume:
 - The number of records in R are n
 - Attribute A has m distinct values in R
- The size of a bitmap index on attribute A is m*n.
- If m is large, then the number of 1's will be around 1/m.
 - Opportunity to encode
- A common encoding approach is called run-length encoding.



--- Run-length encoding

- Represents runs
 - A run is a sequence of i O's followed by a 1, by some suitable binary encoding of the integer i.
- A run of i O's followed by a 1 is encoded by:
 - First computing how many bits are needed to represent i, Say j
 - Then represent the run by j-1 1's and a single 0 followed by j bits which represent i in binary.
 - The encoding for i = 1 is 01. j = 1
 - The encoding for i = 0 is 00. j = 1
- We concatenate the codes for each run together, and the sequence of bits is the encoding of the entire bit-vector



--- Run-length Encoding: Example

- Let us decode the sequence 11101101001011
- Staring at the beginning (left most bit):
 - First run: The first 0 is at position 4, so
 - j = 4,
 - i = 13
 - Second run: 001011
 - j = 1
 - i = 0
 - Last run: 1011
 - J = 1
 - I = 3
- Our entire run length is thus 13,0,3, hence our bit-vector is:

000000000000110001



-- Managing Bitmap Indexes ...

- Finding bit vectors
 - Create secondary key with the attribute value as a search key
 - Btree
 - Hash
- Finding records
 - Create secondary key with the record number as a search key.



... -- Managing Bitmap Indexes

Handling modification

- Assume record numbers are not changed
 - Deletion
 - Tombstone replaces deleted record
 - Corresponding bit is set to 0
 - Insertion
 - Record assigned the next record number.
 - A bit of value 0 or 1 is appended to each bit vector
 - If new record contains a new value of the attribute, add one bitvector.

Modification

- Change the bit corresponding to the old value of the modified record to 0
- Change the bit corresponding to the new value of the modified record to 1
- If the new value is a new value of A, then insert a new bit-vector.



- Summary

- Multidimensional Data
- Grid Files
 - Insertion
 - Different kinds of queries
- Bitmap Index
 - Compression
 - Management



- Reading List

Section 14.2 and 14.4 of GUW



END