



Relational Database Design:

Recap of ICS 334
The first DB course in KFUPM



Announcement

- Read the following paper before next class:
- Code, *Relational Model of Data Large Shared Data Banks*, communication of the ACM, Volume 12, number 6, June 1970.
 - Posted for you in the WebCT.



Lecture objectives

- Briefly mention topics covered in ICS 334.



- Lecture outline

- The main steps in DB Design
- Requirement collection and analysis
- Conceptual data model
- Logical data model
- Physical data model
- Relational Data model



- The main Steps in DB Design

- Requirement collection and analysis
 - How the data is collected and analyzed.
- Data modeling
 - How data should be stored in the database.
 - Includes:
 - Conceptual data model
 - Logical data model
 - Relational data model
 - physical data model
- Functional modeling
 - How the data is processed.



- Requirement collection and analysis

- The goals are:
 - To determine the data requirements of the DB in terms of primitive objects
 - To classify and describe the information about these objects
 - To identify and classify the relationships among the objects
 - To determine the types of transactions that will be executed on the DB and the interactions between the data and the transactions
 - To identify rules governing the integrity of the data
- Requirement analysis can be gathered by **data modeler** from:
 - Existing documents
 - Users
 - Existing systems



-- Example

- Required: A database to record students' enrollment in ICS courses.
- Assume that after requirement analysis we identified the following 3 entities.
 - Course
 - Enroll
 - Student

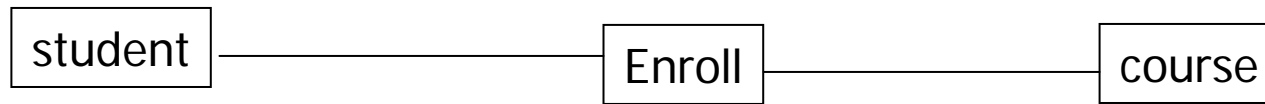


- Conceptual data model

- Objectives:
 - To identify entities:
 - To identify the highest-level relationships among entities.
- No attribute is specified



-- Example:



3 main entities: Student, Enroll, Course



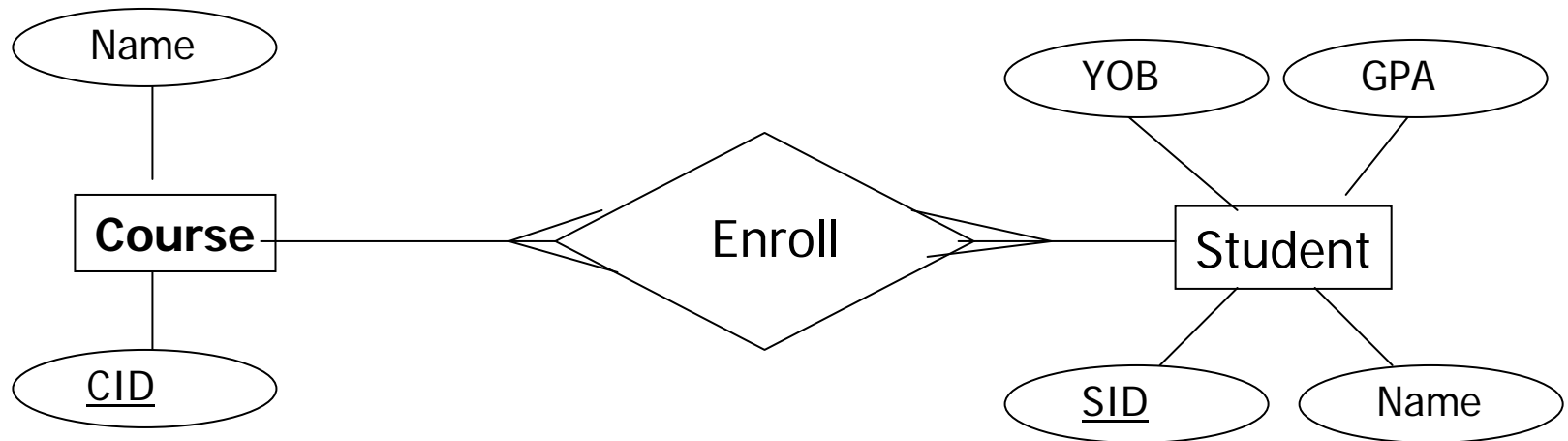
- Logical data model

- Objective:
 - Describe the data in as much detail as possible, without regard to how they will be physically implemented in the database.

- Features include:
 - Entities and relationships among them.
 - All attributes for each entity are specified.
 - The primary key for each entity specified.
 - Foreign keys are specified.
 - Many-to-many relationships are resolved
 - Normalization occurs at this level.



-- Example:



Note: Because relational data model is a logical data model, it should be discussed together with the logical data model. But I prefer discussing it after the physical data model.



- Physical Data Model

- Features:

- Specification all tables and columns
- Foreign keys are used to identify relationships between tables
- Denormalization may occur based on user requirements
- Physical considerations may cause the physical data model to be quite different from the logical data model

- Steps:

- Convert entities into tables
- Convert relationships into foreign keys
- Convert attributes into columns
- Modify the physical data model based on physical constraints / requirements



- The Relational Data Model

- The logical model behind the relational DB (RDB). Data is always represented as relations (2-dim. tables).
- Basic Concepts:
 - Relation
 - Attribute
 - Schema
 - Tuple
 - Keys
 - Constraints
 - Functional dependency
 - Normalization



-- Relation

- The way to represent data is through relations.
- A **relation** is a two-dimensional table.
- The order of rows and columns can be exchanged, and it is still the same relation.

Example:

Relation name →

| CID | Title |
|---------|-----------------|
| ICS 102 | Java |
| ICS 202 | Data structures |
| ICS 334 | Databases |

← Relation



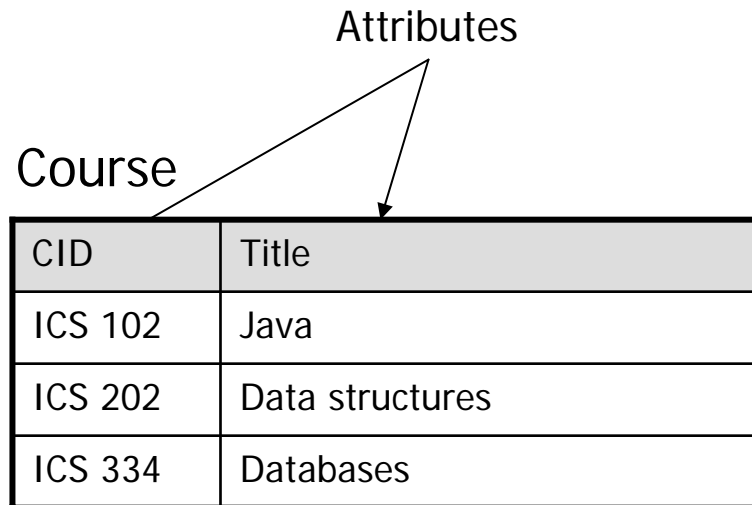
-- Attribute

- An **attribute** is the name of a column in a relation. It usually describes the meaning of the content in the column.
- Example:

Attributes

Course

| CID | Title |
|---------|-----------------|
| ICS 102 | Java |
| ICS 202 | Data structures |
| ICS 334 | Databases |





-- Schema

- A **schema** is a description of a class of relation. It consists of the name of the relation and the set of attributes in the relation.

That it is a set of attributes means that the attributes are unordered.

- Example:

Course

| CID | Title |
|---------|-----------------|
| ICS 102 | Java |
| ICS 202 | Data structures |
| ICS 334 | Databases |

Schema for the above relation: Course(CID, Title)



-- Tuple

- A tuple is a row in a table. A relation can be seen as a set of tuples.
- Example:

Course

| CID | Title |
|---------|-----------------|
| ICS 102 | Java |
| ICS 202 | Data structures |
| ICS 334 | Databases |

← A relation with 3 tuples



-- Keys

- Superkey:
 - A combination of attributes that can be uniquely used to identify a record.
 - A table may have many superkeys
- Candidate key:
 - A superkey but without extraneous data.
 - A table can have one or more candidate keys
- Primary key
 - One of the candidate keys is chosen to be the primary key.
 - There can only be one primary key in a table
- Alternate key
 - All the candidate keys, minus the primary key.
 - The number of alternate keys in a table is one less than the number of candidate keys.



-- constraints

- NOT NULL
 - No null values are allowed in an attribute specified as NOT NULL
- UNIQUE
 - No duplicate values are allowed in a column specified as UNIQUE.
- Primary key
 - Must be UNIQUE and NOT NULL
- Foreign key
 - Must refer to a value of a candidate key in the parent table
 - Can be null
 - Can be duplicate



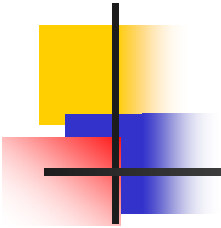
-- Functional dependency

- A functional dependency occurs when one or more attributes in a relation uniquely determine other attribute(s). If A and B are attributes, this can be written $A \twoheadrightarrow B$ which would be the same as stating "B is functionally dependent upon A."
- Attribute A is the **determinant** and attribute B is the **functionally dependent**.
- If the determinant is part of the primary key (and not the whole key), then the dependency is called **partial dependency**.
- If both the determinant and the functionally dependent attributes are non key attributes, then the dependency is called **transitive dependency**.



-- Normalization

- UNF (Un Normalized form)
 - A relation is in UNF if it contains at least one multi-valued attribute
- 1NF
 - A relation is in 1NF if it has no multi-valued attribute.
- 2NF
 - A relations is in 2NF if it is in 1NF and has no partial dependency.
- 3NF
 - A relation is in 3NF if its is in 2NF and contains no transitive dependency.
- BCNF
 - A relation is in BCNF if and only if, every determinant is a candidate key.



END