Relational Database Design:

Recap of ICS 334
The first DB course in KFUPM
Announcement

- Read the following paper before next class:

- Code, *Relational Model of Data Large Shared Data Banks*,
  communication of the ACM, Volume 12, number 6, June 1970.
  
  - Posted for you in the WebCT.
Lecture objectives

- Briefly mention topics covered in ICS 334.
- Lecture outline

- The main steps in DB Design
- Requirement collection and analysis
- Conceptual data model
- Logical data model
- Physical data model
- Relational Data model
- The main Steps in DB Design

- Requirement collection and analysis
  - How the data is collected and analyzed.

- Data modeling
  - How data should be stored in the database.
  - Includes:
    - Conceptual data model
    - Logical data model
    - Relational data model
    - Physical data model

- Functional modeling
  - How the data is processed.
Requirement collection and analysis

The goals are:

- To determine the data requirements of the DB in terms of primitive objects
- To classify and describe the information about these objects
- To identify and classify the relationships among the objects
- To determine the types of transactions that will be executed on the DB and the interactions between the data and the transactions
- To identify rules governing the integrity of the data

Requirement analysis can be gathered by data modeler from:

- Existing documents
- Users
- Existing systems
-- Example

- **Required**: A database to record students’ enrollment in ICS courses.

- Assume that after requirement analysis we identified the following 3 entities.
  - Course
  - Enroll
  - Student
- Conceptual data model

- Objectives:
  - To identify entities:
  - To identify the highest-level relationships among entities.
- No attribute is specified
-- Example:

3 main entities: Student, Enroll, Course
- Logical data model

Objective:
- Describe the data in as much detail as possible, without regard to how they will be physically implemented in the database.

Features include:
- Entities and relationships among them.
- All attributes for each entity are specified.
- The primary key for each entity specified.
- Foreign keys are specified.
- Many-to-many relationships are resolved.
- Normalization occurs at this level.
-- Example:

Note: Because relational data model is a logical data model, it should be discussed together with the logical data model. But I prefer discussing it after the physical data model.
- Physical Data Model

**Features:**

- Specification all tables and columns
- Foreign keys are used to identify relationships between tables
- Denormalization may occur based on user requirements
- Physical considerations may cause the physical data model to be quite different from the logical data model

**Steps:**

- Convert entities into tables
- Convert relationships into foreign keys
- Convert attributes into columns
- Modify the physical data model based on physical constraints / requirements
The Relational Data Model

- The logical model behind the relational DB (RDB). Data is always represented as relations (2-dim. tables).

Basic Concepts:
- Relation
- Attribute
- Schema
- Tuple
- Keys
- Constraints
- Functional dependency
- Normalization
-- Relation

- The way to represent data is through relations.
- A relation is a two-dimensional table.
- The order of rows and columns can be exchanged, and it is still the same relation.

Example:

<table>
<thead>
<tr>
<th>Course</th>
<th>Title</th>
</tr>
</thead>
<tbody>
<tr>
<td>ICS 102</td>
<td>Java</td>
</tr>
<tr>
<td>ICS 202</td>
<td>Data structures</td>
</tr>
<tr>
<td>ICS 334</td>
<td>Databases</td>
</tr>
</tbody>
</table>
-- Attribute

- An attribute is the name of a column in a relation. It usually describes the meaning of the content in the column.

- Example:

<table>
<thead>
<tr>
<th>CID</th>
<th>Title</th>
</tr>
</thead>
<tbody>
<tr>
<td>ICS 102</td>
<td>Java</td>
</tr>
<tr>
<td>ICS 202</td>
<td>Data structures</td>
</tr>
<tr>
<td>ICS 334</td>
<td>Databases</td>
</tr>
</tbody>
</table>
A schema is a description of a class of relation. It consists of the name of the relation and the set of attributes in the relation. That it is a set of attributes means that the attributes are unordered.

Example:

<table>
<thead>
<tr>
<th>CID</th>
<th>Title</th>
</tr>
</thead>
<tbody>
<tr>
<td>ICS 102</td>
<td>Java</td>
</tr>
<tr>
<td>ICS 202</td>
<td>Data structures</td>
</tr>
<tr>
<td>ICS 334</td>
<td>Databases</td>
</tr>
</tbody>
</table>

Schema for the above relation: Course(CID, Title)
-- Tuple

A tuple is a row in a table. A relation can be seen as a set of tuples.

Example:

<table>
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<th>Title</th>
</tr>
</thead>
<tbody>
<tr>
<td>ICS 102</td>
<td>Java</td>
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</tr>
<tr>
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</tr>
</tbody>
</table>

A relation with 3 tuples
-- Keys

- **Superkey:**
  - A combination of attributes that can be uniquely used to identify a record.
  - A table may have many superkeys

- **Candidate key:**
  - A superkey but without extraneous data.
  - A table can have one or more candidate keys

- **Primary key**
  - One of the candidate keys is chosen to be the primary key.
  - There can only be one primary key in a table

- **Alternate key**
  - All the candidate keys, minus the primary key.
  - The number of alternate keys in a table is one less than the number of candidate keys.
-- constraints

- NOT NULL
  - No null values are allowed in an attribute specified as NOT NULL

- UNIQUE
  - No duplicate values are allowed in a column specified as UNIQUE.

- Primary key
  - Must be UNIQUE and NOT NULL

- Foreign key
  - Must refer to a value of a candidate key in the parent table
  - Can be null
  - Can be duplicate
A functional dependency occurs when one or more attributes in a relation uniquely determine other attribute(s). If A and B are attributes, this can be written $A \rightarrow B$ which would be the same as stating "B is functionally dependent upon A."

Attribute A is the **determinant** and attribute B is the **functionally dependent**.

- If the determinant is part of the primary key (and not the whole key), then the dependency is called **partial dependency**.

- If both the determinant and the functionally dependent attributes are non key attributes, then the dependency is called **transitive dependency**.
-- Normalization

- **UNF (Un Normalized form)**
  - A relation is in UNF if it contains at least one multi-valued attribute.

- **1NF**
  - A relation is in 1NF if it has no multi-valued attribute.

- **2NF**
  - A relation is in 2NF if it is in 1NF and has no partial dependency.

- **3NF**
  - A relation is in 3NF if it is in 2NF and contains no transitive dependency.

- **BCNF**
  - A relation is in BCNF if and only if, every determinant is a candidate key.
END