## Haykin Chapter 4 Baseband Pulse Transmission

EE571

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### 4.1 Introduction

- In this chapter we study the transmission of digital data (of whatever origin) over a *baseband channel*.
- Baseband transmission of digital data requires the use of a low-pass channel with a bandwidth large enough to accommodate the essential frequency content of the data stream.
- Typically, however, the channel is *dispersive* in that its frequency response deviates from that of an ideal low-pass filter.

### 4.1 Error Sources in Baseband Transmission

### Intersymbol Interference (ISI)

- The result of data transmission over a *dispersive* channel is that each received pulse is affected somewhat by adjacent pulses, thereby giving rise to a common form of interference called *intersymbol interference* (ISI).
- Intersymbol interference is a major source of bit errors in the reconstructed data stream at the receiver output.
- To correct for it, control has to be exercised over the pulse shape in the overall system.
- Thus much of the material covered in this chapter is devoted to *pulse shaping* in one form or another.

## 4.1 Error Sources in Baseband Transmission

- Another source of bit errors in a baseband data transmission system is the *channel noise*.
- Naturally, noise and ISI arise in the system simultaneously.















### Schwarz's inequality

If we have two complex functions  $\mathcal{O}_1(x)$  and  $\mathcal{O}_2(x)$  in the real variable *x*, satisfying the conditions:

$$\int_{-\infty}^{\infty} |\phi_1(x)|^2 dx < \infty \quad \text{AND} \quad \int_{-\infty}^{\infty} |\phi_2(x)|^2 dx < \infty$$

Then

$$\int_{-\infty}^{\infty} \phi_1(x) \phi_2(x) \ dx \bigg|^2 \leq \int_{-\infty}^{\infty} |\phi_1(x)|^2 \ dx \int_{-\infty}^{\infty} |\phi_2(x)|^2 \ dx$$

The equality in (4.9) holds if, and only if, we have

$$\phi_1(x) = k \phi_2^*(x)$$



















### 4.3 Error Rate Due to Noise

- To determine the average probability of error, we consider these two situations separately.
- Suppose that symbol 0 was sent. Then

$$x(t) = -A + w(t), \qquad 0 \le t \le T_b$$

• The matched filter output, sampled at time  $t = T_b$ , is:

$$y = \int_0^{T_0} x(t) dt$$
$$= -A + \frac{1}{T_b} \int_0^{T_b} w(t) dt$$

• which represents the sample value of a random variable Y

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## **Complementary Error Function**

complementary error function is defined as: ٠

$$\operatorname{erfc}(u) = \frac{2}{\sqrt{\pi}} \int_{u}^{\infty} \exp(-z^2) \, dz$$

- which is closely related to the Gaussian distribution. •
- For large positive values of *u*, we have the following *upper bound* on • the complementary error function:

$$\operatorname{erfc}(u) < \frac{\exp(-u^2)}{\sqrt{\pi u}}$$

• Relation to Q-Function:

$$Q(v) = \frac{1}{\sqrt{2\pi}} \int_{v}^{\infty} \exp\left(-\frac{x^2}{2}\right) dx$$

Q-Function:  

$$Q(v) = \frac{1}{2} \operatorname{erfc}\left(\frac{v}{\sqrt{2}}\right)$$

$$\int_{v}^{\infty} \exp\left(-\frac{x^{2}}{2}\right) dx \qquad \operatorname{erfc}(u) = 2Q(\sqrt{2}u)$$

	и	erf(u)	น	erf(u)
	0.00	0.00000	1.10	0.88021
	0.05	0.05637	1.15	0.89612
	0.10	0.11246	1.20	0.91031
	0.15	0.16800	1.25	0.92290
	0.20	0.22270	1.30	0.93401
	0.25	0.27633	1.35	0.94376
ο ( <sup>∞</sup>	0.30	0.32863	1.40	0.95229
$= \frac{z}{z} \left[ \exp(-z^2) dz \right]$	0.35	0.37938	1.45	0.95970
$\sqrt{\pi} J_{\mu}$	0.40	0.42839	1.50	0.96611
	0.45	0.47548	1.55	0.97162
	0.50	0.52050	1.60	0.97635
	0.55	0.56332	1.65	0.98038
	0.60	0.60386	1.70	0.98379
	0.65	0.64203	1.75	0.98667
	0.70	0.67780	1.80	0.98909
	0.75	0.71116	1.85	0.99111
	0.80	0.74210	1.90	0.99279
	0.85	0.77067	1.95	0.99418
	0.90	0.79691	2.00	0.99532
	0.95	0.82089	2.50	0.99959
	1.00	0.84270	3.00	0.99998
	1.05	0.86244	3.30	0.99999









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### 4.3 Error Rate Due to Noise

- This result is intuitively satisfying as it states that, for the transmission of equiprobable binary symbols, we should choose the threshold at the midpoint between the pulse heights —A and +A representing the two symbols 0 and 1.
- Note that for this special case we also have p<sub>01</sub> = p<sub>10</sub>
- A channel for which the conditional probabilities of error  $p_{01}$  and  $p_{10}$  are equal is said to be *binary symmetric*.

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## 4.3 Error Rate: Example (cont)

• Let 
$$u = \sqrt{\frac{E_b}{N_o}}$$
, then for  $P_e = 10^{-6} = \frac{1}{2} erfc(u)$ , we get u=3.3

• Now when the signaling rate is doubled, the new value of  $\mathrm{P}_{\mathrm{e}}$  is:

$$P_e = \frac{1}{2} \operatorname{erfc} \left[ \frac{u}{\sqrt{2}} \right]$$
$$= \frac{1}{2} \operatorname{erfc} (2.33) = 10^{-3}$$