

**How Much Does Transmit Correlation Affect**

**the Sum-Rate of MIMO Downlink Channels?**

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## Outline

- Introduction
- Questions of interest in a broadcast scenario
- System model and multiuser scheduling schemes
- Capacity scaling of DPC with channel correlation
- Capacity scaling of beamforming with channel correlation
- Simulations
- Conclusion

## Introduction to Broadcast Channels

- Multiple antennas add tremendous value to point to point systems
- Research shifted recently to the role of multiple antennas in multiuser systems
- Broadcast scenarios (point to multi-point) are especially important because downlink scheduling is the major bottleneck for broadband wireless networks

## Three Main Questions in a Broadcast Scenario (1)

**Q1)** Quantify the maximum sum rate possible to all users

**A1)** Sum-rate is achieved using dirty paper coding (DPC) (Caire and

Shamai '02, Viswanath and Tse '02, Vishwanath et al. '02, Yu and Cioffi '02)

(-) DPC is computationally complex at both Tx and Rx

(-) Requires a great deal of Feedback (CSI for all users at Tx)

## Three Main Questions in a Broadcast Scenario (2)

The second question is motivated by the drawbacks of DPC

**Q2)** Devise computationally efficient algorithms for capturing capacity

**A2)** Utilize multi-user diversity to achieve performance close to capacity

(+) Opportunist multiple random beamforming coincides asymptotically with DPC (Sharif and Hassibi '06)

$$R = M \log \log n + M \log \frac{P}{M} + o(1)$$

(+) Requires simply SINR feedback to Tx

## Three Main Questions in a Broadcast Scenario (3)

**Q3)** With this promising performance, how does opportunist beam-forming perform under various non-idealities

**A3)** (i) Time correlation (Kountouris and Gesbert '05)

(ii) Frequency correlation (Fakhereddin, Sharif, and Hassibi'06)

(iii) Channel estimation error (Vikali, Sharif, and Hassibi '06)

(iv) Spatial correlation (D. Park and S Y. Park '05)

**Main problem to be addressed:**

- For a Gaussian broadcast channel, we would like to quantify the hit that transmit correlation causes to scaling laws of the sum-rate capacity. We consider DPC and various beamforming schemes.

## System Model

- Base station with  $M$  antennas broadcasting to  $n$  single-antenna users
- Received signal at each antenna

$$Y_i = \sqrt{P}H_i S + W_i, \quad i = 1, \dots, n$$

with  $E[S^* S] = 1$  and Gaussian noise  $W_i \sim CN(0, I)$

- Channel  $H_i$  of  $i$ -th user is  $1 \times M$  vector
  - Distributed as  $CN(0, R)$ ;  $R$  is nonsingular with  $\text{tr}(R) = M$
  - Known perfectly at receiver
  - Follows a block fading model (with coherence interval  $T$ )
  - $H_i$  is independent from one user to another

## Digression: Extreme Value Theory

- Let  $x_1, x_2, \dots, x_n$  be i.i.d random variables with pdf  $f(x)$  and CDF  $F(x)$ . How does  $\max_i x_i$  behave?
- Let  $z$  denote the limit

$$z = \lim_{x \rightarrow \infty} \frac{1 - F(x)}{f(x)}$$

then, for large  $n$  we have with high probability

$$\max_i x_i = z \log(n)$$



## Scaling of DPC under Correlation

- Sum-rate capacity of DPC

$$R_{DPC} = E \left\{ \max_{\{P_1, \dots, P_n, \sum P_i = P\}} \log \det \left( I + \sum_{i=1}^n H_i^* P_i H_i \right) \right\}.$$

- Define  $H_i = H_{w_i} R^{\frac{1}{2}}$  and employ the inequality  $\det(A) \leq \left(\frac{\text{tr}(A)}{M}\right)^M$  to obtain

$$R_{DPC} \leq M \log \left( \frac{1}{M} \text{tr}(R^{-1}) + \max_i \|H_{w_i}\|^2 \frac{P}{M} \right)$$

- For large  $n$ ,  $\max_i \|H_{w_i}\|^2$  behaves as  $\log n$  with high probability.

Thus,

$$\begin{aligned} R_{DPC} &\leq M \log \left( \frac{\text{tr}(R^{-1})}{M} + \frac{P}{M} \log n \right) + \log \det R \\ &= M \log \log n + M \log \left( \frac{P}{M} \right) + M \log \sqrt[M]{\det R} \text{ for large } n \end{aligned}$$

- This is also a lower bound as it the scaling of deterministic beam forming. So, for large  $n$

$$R_{DPC} = M \log \log n + M \log \left( \frac{P}{M} \right) + M \log \sqrt[M]{\det R}$$

- Compare with rate for spatially uncorrelated channel

$$R_{DPC} = M \log \log n + M \log \left( \frac{P}{M} \right)$$

Keep in mind that  $\sqrt[M]{R} \leq \frac{\text{Tr}(R)}{M} = 1$

## What is Random Beam Forming?

- Choose  $M$  random orthonormal vectors  $\phi_m$ ,  $m = 1, \dots, M$  (according to an isotropic distribution)
- Construct the signal

$$S(t) = \sum_{m=1}^M \phi_m s_m(t), \quad t = 1, \dots, T$$

where  $T$  is less than the coherence interval of the channel.

- After  $T$  channel uses we independently choose another isotropic set of orthonormal vectors  $\{\phi_m\}$ , and so on. So we are transmitting  $M$  random beams.
- This is a generalization of the scheme “Opportunistic Beamforming” (Viswanath et al. '02) in which only one random beam is transmitted and proportional fairness is guaranteed.

## Exploit Multi-User Diversity

- Each receiver  $i = 1, \dots, n$  computes the following  $M$  SINRs

$$\text{SINR}_{i,m} = \frac{|H_i \phi_m|^2}{1/\rho + \sum_{n \neq m} |H_i \phi_n|^2}, \quad m = 1, \dots, M$$

and feeds back the best SINR

- Rather than randomly assigning the beams, the transmitter assigns signal  $s_m$  to the user with the best SINR for that signal. Therefore

$$C = E \sum_{m=1}^M \log \left( 1 + \max_{i=1, \dots, n} \text{SINR}_{i,m} \right)$$

- Due to the symmetry of all the random variables involved:

$$C = ME \log \left( 1 + \max_{i=1, \dots, n} \text{SINR}_{i,1} \right)$$

## Other Beamforming Schemes

- Random Beam forming (RBF)  $S(t) = \sum_{m=1}^M \phi_m s_m(t)$

- RBF with Channel whitening

$$S(t) = \sum_{m=1}^M \sqrt{\alpha} R^{-1/2} \phi_m s_m(t)$$

- RBF with general precoding

$$S(t) = \sum_{m=1}^M \sqrt{\alpha} A \phi_m s_m(t)$$

- Deterministic beamforming

$$S(t) = \sum_{m=1}^M \phi_m s_m(t), \quad \phi_m \text{'s are fixed}$$

## How to Determine Scaling of BF Schemes

1. Sum rate

$$\begin{aligned} R_{\text{BF}} &= E \sum_{m=1}^M \log \left( 1 + \max_{i=1, \dots, n} \text{SINR}_{i,m} \right) \\ &= ME \left( 1 + \max_{i=1, \dots, n} \text{SINR}_{i,m} \right) \end{aligned}$$

2. To calculate expectation, condition on beams

$$R_{\text{BF}|\Phi} = ME_{H_i|\Phi} \left( 1 + \max_{i=1, \dots, n} \text{SINR}_{i,m} \right)$$

- $\text{SINR}_{i,m}|\Phi$  is iid over  $i$
- Find the distribution of  $\text{SINR}_{i,m}|\Phi$
- Employ extreme value theory to find  $\max_{i=1, \dots, n} \text{SINR}_{i,m}$

3. Average  $R_{\text{BF}|\Phi}$  over  $\Phi$

## Statistics of $\text{SINR}_{i,m}$ (White Channel)

- $\text{SINR}_{i,m}$  is defined by

$$\text{SINR}_{i,m} = \frac{|H_i \phi_m|^2}{1/\rho + \sum_{n \neq m} |H_i \phi_n|^2}, \quad m = 1, \dots, M$$

- Easy to find distribution of  $\text{SINR}_{i,m} | \Phi$  when  $H_i$  is white

$$f(x) = \frac{e^{-\frac{x}{\rho}}}{(1+x)^M} \left( \frac{1}{\rho}(1+x) + M - 1 \right)$$

$$F(x) = 1 - \frac{e^{-\frac{x}{\rho}}}{(1+x)^M}$$

- Finding these statistics in the correlated case is challenging

## Statistics of $\text{SINR}_{i,m}$ Given $\Phi$ (Correlated Case)

- We can show that the CDF of SINR in the correlated case

$$F(x) = 1 - \frac{1}{2\pi^M \det(R)} \lambda_M \prod_{i=1}^{M-1} \frac{\lambda_i \lambda_M}{x(\lambda_i - \lambda_M)} e^{-\frac{1}{\rho} \frac{x}{\lambda_M}}$$

where  $\lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_M$  are the eigenvalues of the matrix

$$A = (1 + x)\Lambda^{1/2} \phi_m \phi_m^* \Lambda^{1/2} - x\Lambda \quad \rho = \frac{P}{M}$$

Note that eigenvalues are a function of  $x$ .

- pdf is given by

$$f(x) = \frac{1}{2\pi^M \det(R)} e^{-\frac{1}{\rho} \frac{x}{\lambda_M}} \prod_{i=1}^{M-1} \frac{\lambda_i \lambda_M}{x(\lambda_i - \lambda_M)} \times \left\{ \frac{1}{\rho} \frac{\|q_M\|_C^2}{\lambda_M} - \|q_M\|_B^2 - \sum_{i=1}^M \frac{1}{\lambda_i} \frac{\lambda_M^2 \|q_i\|_C^2 - \lambda_i^2 \|q_M\|_C^2}{x(\lambda_i - \lambda_M)} \right\}$$

$$\text{where } B = \Lambda^{1/2}(\phi_m \phi_m^* - I)\Lambda^{1/2} \quad C = \Lambda^{1/2} \phi_m \phi_m^* \Lambda^{1/2}$$



## Scaling of Maximum SINR

- Can now show

$$\lim_{x \rightarrow \infty} \frac{1 - F(x)}{f(x)} = \frac{P}{M} \frac{1}{\|\phi_m\|_{\Lambda-1}^2}$$

- Using extreme value theory, we can show that for large  $n$

$$\max_{i=1, \dots, n} \text{SINR}_{i,m} = \frac{P}{M} \frac{1}{\|\phi_m\|_{\Lambda-1}^2} \log n$$

- Conditional sum-rate capacity scales as

$$R_{\text{BF}|\Phi} = M \log \log n + M \log \frac{P}{M} + M \log \left( \frac{1}{\|\phi_m\|_{\Lambda-1}^2} \right)$$

- Sum-rate capacity of random beam-forming

$$R_{\text{RBF}} = M \log \log n + M \log \frac{P}{M} + M E_{\Phi} \log \left( \frac{1}{\|\phi_m\|_{\Lambda-1}^2} \right)$$

## Averaging Over the Random Beams

- Need to obtain CDF of  $\frac{1}{\|\phi_m\|_{\Lambda^{-1}}^2}$  which is challenging.
- The CDF of  $y = \frac{1}{\|\phi\|_{\Lambda^{-1}}^2}$  is given by

$$G(x) = Pr\left(\frac{1}{\|\phi\|_{\Lambda^{-1}}^2} < x\right) = 1 - \sum_i \eta_i \left(\frac{1}{x} - \frac{1}{\lambda_i(\Lambda)}\right)^{M-1} u\left(1 - \frac{x}{\lambda_i(\Lambda)}\right)$$

where  $\eta_i = \frac{1}{\prod_{j \neq i} \left(\frac{1}{\lambda_j(\Lambda)} - \frac{1}{\lambda_i(\Lambda)}\right)}$

- Use CDF to show that

$$R_{RBF} = M \log \log n + M \log \frac{P}{M} + \log \lambda_1(\Lambda) + \sum_{i=1}^M \eta_i \log \left(\frac{\lambda_i}{\lambda_1}\right) \sum_{k=1}^{M-1} \frac{1}{k+2} \left(\frac{-1}{\lambda_i}\right)^{M-1-k} \frac{1}{y^{k+2}} \Big|_{\lambda_1}^{\lambda_i}$$

## Sum rate of Deterministic Beam Forming

- Sum-rate of deterministic beam forming

$$R_{BF-D} = M \log \log n + M \log \frac{P}{M} + \sum_{i=1}^M \log \left( \frac{1}{\phi_i^* U^* \Lambda^{-1} U \phi_i} \right)$$

$U^* \Lambda^{-1} U$  is the eigenvalue decomposition of  $R^{-1}$

- Special case:  $U \phi_i$ 's are the columns of identity matrix

$$R_{BF-D} = M \log \log n + M \log \frac{P}{M} + M \log \sqrt[M]{\det R}$$

Since  $\text{tr}(R) = M$ , the geometric mean satisfies  $\det(R) \leq 1$

- Scaling coincides with (D. Park and S Y. Park '05) which focused on the  $M = 2$  case

## Sum rate of RBF with Channel Whitening

- For random beam forming with channel whitening,

$$S(t) = \sum_{m=1}^M \sqrt{\alpha} R^{-1/2} \phi_m s_m(t)$$

- Set  $\alpha = \frac{\text{tr}(R^{-1})}{M}$  to guarantee  $E[S^* S] \leq 1$
- Scaling becomes the same as for white channel case with reduced signal power

$$R_{BF-W} = M \log \log n + M \log \frac{P}{M} + M \log \frac{M}{\text{tr}(R^{-1})}$$

## Simulations

- Consider a base station with  $M = 2$  and  $M = 3$  antennas
- The corresponding correlation matrix is

$$R = \begin{bmatrix} 1 & \alpha \\ \alpha & 1 \end{bmatrix}$$

$$R = \begin{bmatrix} 1 & \alpha & \alpha^2 \\ \alpha & 1 & \alpha \\ \alpha^2 & \alpha & 1 \end{bmatrix}$$

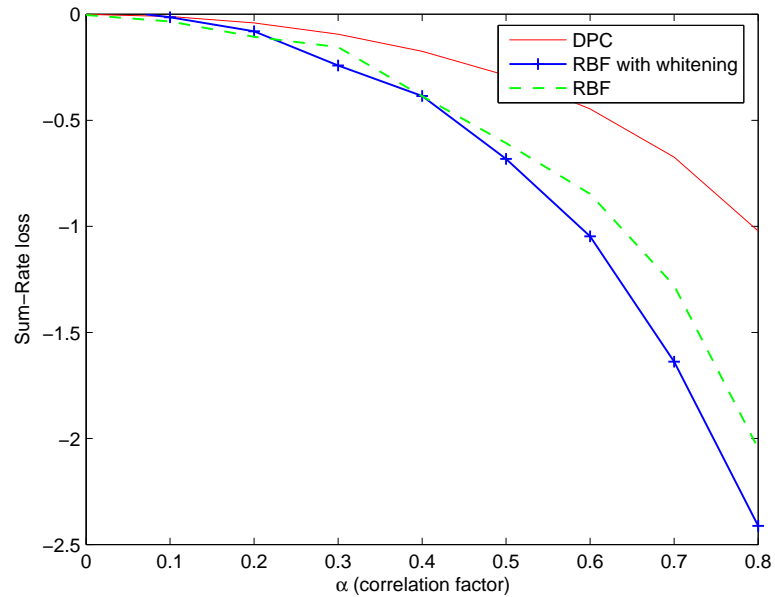


Figure 1: Sum-rate loss versus the correlation factor  $\alpha$  for a system with  $M = 2$  and  $n = 100$ .

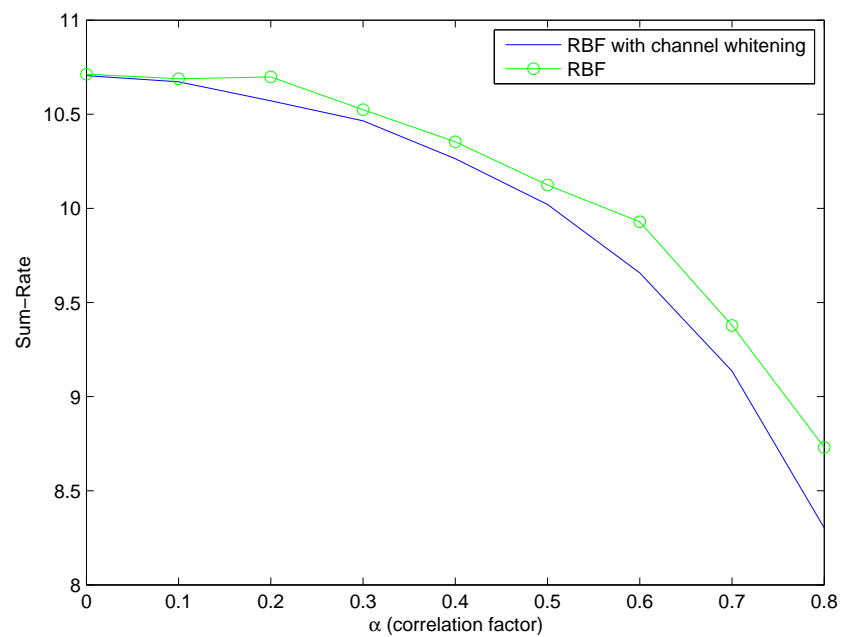


Figure 2: Sum-rate versus the correlation factor  $\alpha$  for a system with  $M = 2$ ,  $P = 10$ , and  $n = 100$ .

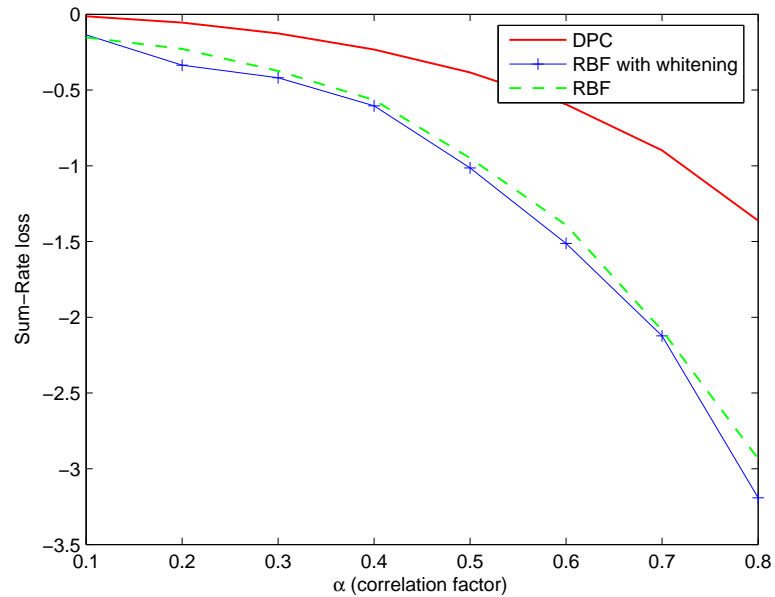


Figure 3: Sum-rate loss versus the correlation factor  $\alpha$  for a system with  $M = 3$  and  $n = 100$ .



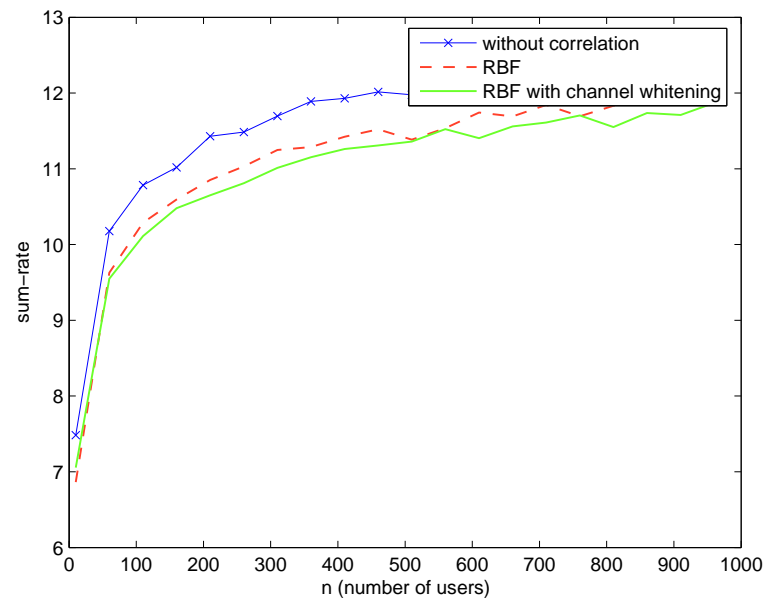


Figure 4: Sum-rate versus the number of users in a system with  $M = 2$  and  $\alpha = 0.5$

## Conclusion

- Studied the effect of spatial correlation on various multiuser schemes for MIMO broadcast channels.
- Considered DPC and random, deterministic, and channel whitening schemes.
- All these techniques exhibit the same scaling for iid channels

$$R_{\text{sum-rate}} = M \log \log n + M \log \frac{P}{M}$$

- In the presence of correlation between transmit antennas, scaling is

$$R_{\text{sum-rate}} = M \log \log n + M \log \frac{P}{M} + M \log c$$

The constant  $0 < c \leq 1$  depends on the scheduling scheme and the eigenvalues of the correlation matrix  $R$ .

## Recent Results:

### Scaling Laws of Group Broadcast Channels

- $K$  groups of users
- Each group of users is interested in the same data
- Worst user of each group is the bottle neck
- Worst user is difficult to define in the multi-antenna case
- For  $K$  groups with  $n$  users each, we show that capacity scales like

$$C = K \frac{P}{nM}$$

- We show that to have a constant rate,  $M$  should grow at least as fast as  $\log n$
- This is a joint work with Amir Dana and Babak Hassibi, Cal Tech.

## Extra Slide: Finding the Distribution of $SINR$

- Consider the  $SINR$  for the first beam

$$SINR_{i,1} = \frac{|H_i \phi_1|^2}{1/\rho + \sum_{n=2}^M |H_i \phi_n|^2},$$

- Define  $S$  by

$$S = -\frac{x}{\rho} + H_i^* ((1+x)\phi_1\phi_1^* - xI)H_i$$

Then

$$\begin{aligned} P(SINR_{i,1} > x) = P(S > 0) &= \int_{-\infty}^{\infty} P(H_i)u(S)dH_i \\ &= \frac{1}{\pi^M \det(R)} \int_{-\infty}^{\infty} e^{-H_i^* R^{-1} H_i} u(S) dH_i \end{aligned}$$

- To evaluate integral, use the integral representation of unit step

$$u(S) = \frac{1}{2\pi} \int_{-\infty}^{\infty} \frac{e^{(j\omega + \beta)S}}{j\omega + \beta} d\omega$$

- Desired probability becomes

$$\begin{aligned}
& P(\text{SINR}_{i,1} > x) \\
&= \frac{1}{2\pi^{M+1} \det(R)} \int_{-\infty}^{\infty} d\omega \frac{1}{j\omega + \beta} \int_{-\infty}^{\infty} dH_i e^{(j\omega + \beta)S - H_i^* R^{-1} H_i} \\
&= \frac{1}{2\pi^{M+1} \det(R)} \int_{-\infty}^{\infty} d\omega \frac{e^{-(j\omega + \beta)\frac{x}{\rho}}}{j\omega + \beta} \int_{-\infty}^{\infty} dH_i e^{-H_i^* \tilde{R}^{-1} H_i} \\
&= \frac{1}{2\pi^{M+1} \det(R)} \int_{-\infty}^{\infty} d\omega \frac{e^{-(j\omega + \beta)\frac{x}{\rho}}}{j\omega + \beta} \frac{1}{\det(\tilde{R})}
\end{aligned}$$