

Multiobjective VLSI Cell Placement

Using Distributed Simulated Evolution
Algorithm

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What this Paper is About

- “ Parallelization of an Evolutionary Heuristic for wire length, power and delay optimized VLSI cell placement is presented
- “ An improved Parallel SimE Algorithm for Cell Placement is proposed and results are compared with a previous approach

Need for Parallelism

- “ For large test cases and multiobjective optimization, SimE has large runtime requirements
- “ SimE, like other stochastic heuristics, is blind and has to be told when to stop
- “ Can consume hours of CPU time depending upon problem size, complexity and stopping criteria

Cost Functions

” Objectives

- . Reducing overall wire length
- . Optimizing power consumption
- . Improving timing performance (delay)

” Constraint

- . Layout width should be within set limit

Wire Length Estimation

- “ Wire length for each net is estimated using an approximate Steiner Tree Algorithm
- “ Total wire length of whole placement is computed by adding individual wire length estimates of each net

$$Cost_{wire} = \sum_{i \in M} l_i$$

- . where l_i is the wire length estimation for net and M denotes total number of nets in circuit.

Power Estimation

” Power consumption p_i of a net i in a circuit can be given as

$$p_i \simeq \frac{1}{2} \cdot C_i \cdot V_{DD}^2 \cdot f \cdot S_i \cdot \alpha$$

- C_i where is total capacitance of net , V_{DD} is the supply voltage, f is the clock frequency,
- S_i is the switching probability of net, and α is a technology dependent constant

Power Estimation (cont.)

- “ Assuming a fix supply voltage and clock frequency, we have:

$$p_i \simeq C_i \cdot S_i$$

- “ The capacitance C_i of cell i is given as:

$$C_i = C_i^r + \sum_{j \in M_i} C_j^g$$

- “ Moreover, C_i^r depends on wire length l_i of net i , so above equation can be written as:

$$p_i \simeq l_i \cdot S_i$$

Power Estimation (cont.)

“ The cost function for estimate of total power consumption in the circuit can be given as:

$$Cost_{power} = \sum_{i \in M} p_i = \sum_{i \in M} (l_i \cdot S_i)$$

Delay Estimation

- “ Delay along the longest path in the circuit
- “ Delay T_π of a path π consisting of nets $\{v_1, v_2, \dots, v_k\}$, is expressed as:

$$T_\pi = \sum_{i=1}^{k-1} (CD_i + ID_i)$$

- . Where CD_i is switching delay of cells driving net v_i , ID_i is interconnect delay of net v_i
- “ Since CD_i is placement independent, delay cost is given by:

$$Cost_{delay} = \max\{T_\pi\}$$

width cost

- “ Given by the maximum of all the row widths in the layout
- “ w_{avg} is minimum possible layout width
 - . obtained by dividing the total width of all the cells in the layout by the number of rows in the layout

width constraint

“ Layout width should not exceed a certain positive ratio α to the average row width

w_{avg}

$$Width - w_{avg} \leq \alpha \times w_{avg}$$

. where Width is the width cost computed

Fuzzy Multi-Objective Function

- “ A cost function that represents the effect of all three objectives in form of a single quantity
- “ Use of fuzzy logic to integrate multiple, possibly conflicting objectives into a scalar cost function

Fuzzy Logic Rule

- “ Fuzzy logic allows us to describe the objectives in terms of linguistic variables
- “ Fuzzy rules are used to find the overall cost of a placement solution
- “ Following rule is used
 - “**IF** a solution has **SMALL** wirelength **AND** **LOW** power consumption **AND** **SHORT** delay **THEN** it is an **GOOD** solution+

Fuzzy Membership Function

- “ Fuzzy rule is translated to and-like OWA fuzzy operator
- “ Membership $\mu(x)$ of a solution x in fuzzy set GOOD solution is given as:

$$\mu(x) = \begin{cases} \beta \cdot \min_{j=p,d,l} \{\mu_j(x)\} + (1 - \beta) \cdot \frac{1}{3} \sum_{j=p,d,l} \mu_j(x); & \text{if } Width - w_{avg} \leq \alpha \cdot w_{avg}, \\ 0; & \text{otherwise.} \end{cases}$$

- . where $\mu_j(x)$ for $j= p, d, l$, width are membership values in fuzzy sets for power, delay and wire length,
- . β is a constant in the range $[0,1]$

Evolution (SimE)

Algorithm

- “ A general search strategy
- “ Operates on a single solution termed as population
- “ Has a main loop consisting of 3 main steps
 - . Evaluation
 - . Selection
 - . Allocation

Three Operators in SimE Algorithm

- “ **Evaluation** . calculation of goodness of each element of population
- “ **Selection** . process of selecting elements to be reassigned locations in the current solution
- “ **Allocation** . Mutate the population by altering locations of selected cells

Simple Algorithmic Description

Algorithm Slave_Process($CurS, \Phi_s$)

Notation

- (* B is the bias value. *)
- (* $CurS$ is the current solution. *)
- (* Φ_s are the rows assigned to slave s . *)
- (* m_i is module i in Φ_s . *)
- (* g_i is the goodness of m_i . *)

Begin

Receive Placement_And_Indices

Evaluation:

ForEach $m_i \in \Phi_s$ evaluate g_i ;

Selection:

ForEach $m_i \in \Phi_s$ **DO**

Begin

If $Random > Min(g_i + B, 1)$

Then

Begin

$S = S \cup m_i$; Remove m_i from Φ_s

End

End

Sort the elements of S

Allocation:

ForEach $m_i \in S$ **Do**

Begin

Allocate(m_i, Φ_s)

(* Allocate m_i in local partial solution rows Φ_s . *)

End

Send_Partial_Placement_Rows

End. (*Slave_Process*)

Distributed SimE Algorithm

- “ Workload partitioning by dividing rows in a placement (population)
- “ Each PE computes the 3 SimE operators on assigned rows (a sub-population)
- “ Individual Sub-populations are merged after each iteration and new sub-populations created and distributed among PEs

Distributed SimE Algorithm

Master PE does:

- “ Receive partial placement from all PEs, combine them and evaluate fitness,
- “ Re-partition to obtain new allocations,
- “ Distribute new partial placements among PEs

```
Algorithm Parallel_Simulated_Evolution
  Read_User_Input_Parameters
  Read_Input_Files
Begin
  Construct_Initial_Placement
  Repeat
    Generate_Random_Row-Indices
    ParFor
      Slave_Process( $CurS, \Phi_s$ )
    (* Broadcast Cur Placement And Row-Indices. *)
    EndParFor
    ParFor
      Receive_Partial_Placement_Rows
    EndParFor
    Construct_Complete_Solution
    Calculate_Cost
  Until (Stopping_Criteria_is_Satisfied)
  Return_Best_Solution.
End. (*Parallel_Simulated_Evolution*)
```

Proposed Improvement

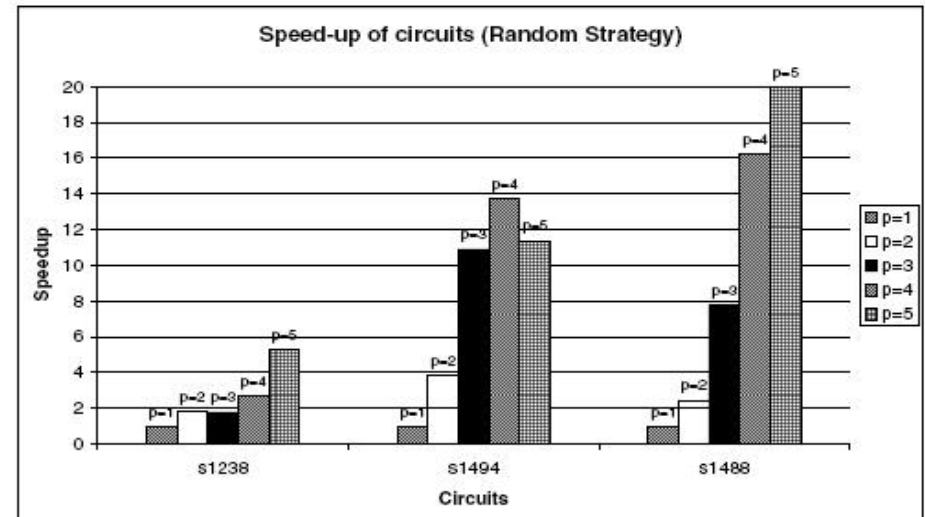
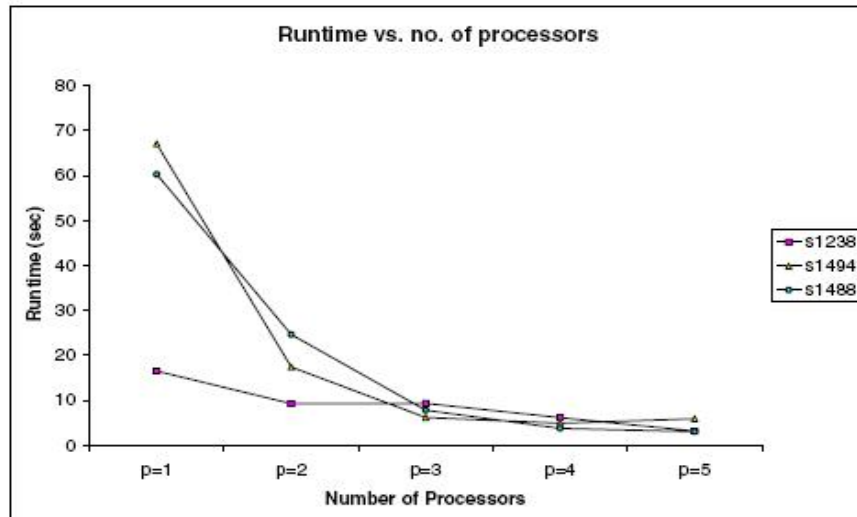
- “ Originally proposed row distribution comprises alternating block and row assignments
- “ Solution qualities inferior due to
 1. Lack of a global placement view to all PEs
 2. Restrictive cell movement due to a fixed allocation pattern
- “ Our solution addresses the second problem

Randomized Rows Assignment

- “ Restrictive cell movement can be alleviated using better row assignment
- “ An assignment that facilitates better inter-mixing among partitions would be intuitively better
- “ Our experimentation with a randomized row assignment gave better results

Experimental Results

Circuit Name	# of Cells	Random Row Distribution					Qual Fixed	Fixed Row Distribution			
		$N_p=1$	$N_p=2$	$N_p=3$	$N_p=4$	$N_p=5$		$N_p=2$	$N_p=3$	$N_p=4$	$N_p=5$
s641	433	UH	4.99	4.97	3.99	3.87	79.7%	9.14	1.08	0.76	0.55
s1238	540	16.5	9.24	9.29	6.12	3.14	95.8%	17.83	8.47	11.30	5.71
s1494	661	67	17.4	6.15	4.88	5.89	82.3%	2.77	1.85	1.76	4.34
s1488	667	60.23	24.6	7.78	3.72	3.02	96.6%	22.0	4.89	5.1	16
s3330	1961	UH	678.02	115	108.5	49.14	33.8%	316	215	4.6	3.4
s5378	2993	UH	1620	338.2	286.6	178.6	46.8%	UH	UH	124.3	95.0



Future work

- “ Evaluation of SimE algorithm parameters for further improvement in parallel version
- “ Use of processor relieve strategy as quality stagnates to enable final solution qualities equivalent to serial version but with improved runtimes

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