

Data Representation

ICS 233

Computer Architecture and Assembly Language

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Presentation Outline

- ❖ Positional Number Systems
- ❖ Binary and Hexadecimal Numbers
- ❖ Base Conversions
- ❖ Integer Storage Sizes
- ❖ Binary and Hexadecimal Addition
- ❖ Signed Integers and 2's Complement Notation
- ❖ Sign Extension
- ❖ Binary and Hexadecimal subtraction
- ❖ Carry and Overflow
- ❖ Character Storage

Positional Number Systems

Different Representations of Natural Numbers

- XXVII Roman numerals (not positional)
- 27 Radix-10 or **decimal** number (positional)
- 11011_2 Radix-2 or **binary** number (also positional)

Fixed-radix positional representation with k digits

Number N in radix $r = (d_{k-1}d_{k-2} \dots d_1d_0)_r$

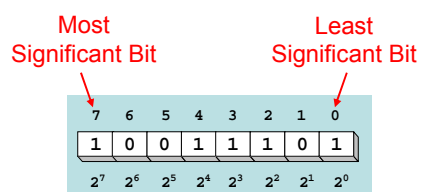
Value = $d_{k-1} \times r^{k-1} + d_{k-2} \times r^{k-2} + \dots + d_1 \times r + d_0$

Examples: $(11011)_2 = 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2 + 1 = 27$

$(2103)_4 = 2 \times 4^3 + 1 \times 4^2 + 0 \times 4 + 3 = 147$

Binary Numbers

- ❖ Each binary digit (called bit) is either 1 or 0
- ❖ Bits have no inherent meaning, can represent
 - ❖ Unsigned and signed integers
 - ❖ Characters
 - ❖ Floating-point numbers
 - ❖ Images, sound, etc.



- ❖ Bit Numbering
 - ❖ Least significant bit (LSB) is rightmost (bit 0)
 - ❖ Most significant bit (MSB) is leftmost (bit 7 in an 8-bit number)

Converting Binary to Decimal

- ❖ Each bit represents a power of 2
- ❖ Every binary number is a sum of powers of 2
- ❖ Decimal Value = $(d_{n-1} \times 2^{n-1}) + \dots + (d_1 \times 2^1) + (d_0 \times 2^0)$
- ❖ Binary $(10011101)_2 = 2^7 + 2^4 + 2^3 + 2^2 + 1 = 157$

7	6	5	4	3	2	1	0
1	0	0	1	1	1	0	1
2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0

Some common powers of 2 →

2^n	Decimal Value	2^n	Decimal Value
2^0	1	2^8	256
2^1	2	2^9	512
2^2	4	2^{10}	1024
2^3	8	2^{11}	2048
2^4	16	2^{12}	4096
2^5	32	2^{13}	8192
2^6	64	2^{14}	16384
2^7	128	2^{15}	32768

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Convert Unsigned Decimal to Binary

- ❖ Repeatedly divide the decimal integer by 2
- ❖ Each remainder is a binary digit in the translated value

Division	Quotient	Remainder
37 / 2	18	1
18 / 2	9	0
9 / 2	4	1
4 / 2	2	0
2 / 2	1	0
1 / 2	0	1

← least significant bit

$37 = (100101)_2$

← most significant bit

← stop when quotient is zero

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Hexadecimal Integers

- ❖ 16 Hexadecimal Digits: 0 – 9, A – F
- ❖ More convenient to use than binary numbers

Binary, Decimal, and Hexadecimal Equivalents

Binary	Decimal	Hexadecimal	Binary	Decimal	Hexadecimal
0000	0	0	1000	8	8
0001	1	1	1001	9	9
0010	2	2	1010	10	A
0011	3	3	1011	11	B
0100	4	4	1100	12	C
0101	5	5	1101	13	D
0110	6	6	1110	14	E
0111	7	7	1111	15	F

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Converting Binary to Hexadecimal

- ❖ Each hexadecimal digit corresponds to 4 binary bits
- ❖ Example:

Convert the 32-bit binary number to hexadecimal

1110 1011 0001 0110 1010 0111 1001 0100

- ❖ Solution:

E	B	1	6	A	7	9	4
1110	1011	0001	0110	1010	0111	1001	0100

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Converting Hexadecimal to Decimal

- ❖ Multiply each digit by its corresponding power of 16

$$\text{Value} = (d_{n-1} \times 16^{n-1}) + (d_{n-2} \times 16^{n-2}) + \dots + (d_1 \times 16) + d_0$$

- ❖ Examples:

$$(1234)_{16} = (1 \times 16^3) + (2 \times 16^2) + (3 \times 16) + 4 =$$

Decimal Value 4660

$$(3BA4)_{16} = (3 \times 16^3) + (11 \times 16^2) + (10 \times 16) + 4 =$$

Decimal Value 15268

Converting Decimal to Hexadecimal

- ❖ Repeatedly divide the decimal integer by 16
- ❖ Each remainder is a hex digit in the translated value

Division	Quotient	Remainder
422 / 16	26	6
26 / 16	1	A
1 / 16	0	1

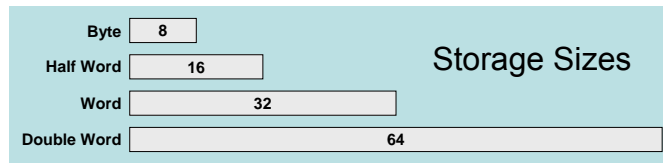
← least significant digit

← most significant digit

stop when
quotient is zero

Decimal 422 = 1A6 hexadecimal

Integer Storage Sizes



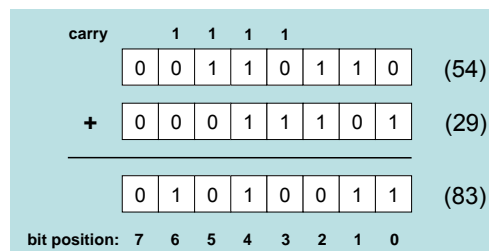
Storage Type	Unsigned Range	Powers of 2
Byte	0 to 255	0 to $(2^8 - 1)$
Half Word	0 to 65,535	0 to $(2^{16} - 1)$
Word	0 to 4,294,967,295	0 to $(2^{32} - 1)$
Double Word	0 to 18,446,744,073,709,551,615	0 to $(2^{64} - 1)$

What is the largest 20-bit unsigned integer?

Answer: $2^{20} - 1 = 1,048,575$

Binary Addition

- ❖ Start with the least significant bit (rightmost bit)
- ❖ Add each pair of bits
- ❖ Include the carry in the addition, if present



Hexadecimal Addition

- ❖ Start with the least significant hexadecimal digits
- ❖ Let Sum = summation of two hex digits
- ❖ If Sum is greater than or equal to 16
 - ❖ Sum = Sum - 16 and Carry = 1
- ❖ Example:

carry:							
			1	1		1	
	1	C	3	7	2	8	6
+	9	3	9	5	E	8	4
<hr/>							
	A	F	C	D	1	0	B

A + B = 10 + 11 = 21
Since 21 ≥ 16
Sum = 21 - 16 = 5
Carry = 1

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Signed Integers

- ❖ Several ways to represent a signed number
 - ❖ Sign-Magnitude
 - ❖ Biased
 - ❖ 1's complement
 - ❖ 2's complement
- ❖ Divide the range of values into 2 equal parts
 - ❖ First part corresponds to the positive numbers (≥ 0)
 - ❖ Second part correspond to the negative numbers (< 0)
- ❖ Focus will be on the 2's complement representation
 - ❖ Has many advantages over other representations
 - ❖ Used widely in processors to represent signed integers

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Two's Complement Representation

❖ Positive numbers

❖ Signed value = Unsigned value

❖ Negative numbers

❖ Signed value = Unsigned value $- 2^n$

❖ n = number of bits

❖ Negative weight for MSB

❖ Another way to obtain the signed value is to assign a negative weight to most-significant bit

1	0	1	1	0	1	0	0
-128	64	32	16	8	4	2	1

$$= -128 + 32 + 16 + 4 = -76$$

8-bit Binary value	Unsigned value	Signed value
00000000	0	0
00000001	1	+1
00000010	2	+2
...
01111110	126	+126
01111111	127	+127
10000000	128	-128
10000001	129	-127
...
11111110	254	-2
11111111	255	-1

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Forming the Two's Complement

starting value	00100100 = +36
step1: reverse the bits (1's complement)	11011011
step 2: add 1 to the value from step 1	+ 1
sum = 2's complement representation	11011100 = -36

Sum of an integer and its 2's complement must be zero:

$$00100100 + 11011100 = 00000000 \text{ (8-bit sum)} \Rightarrow \text{Ignore Carry}$$

Another way to obtain the 2's complement:

Start at the least significant 1

Leave all the 0s to its right unchanged

Complement all the bits to its left

Binary Value

= 00100 100 least significant 1

2's Complement

= 11011 100

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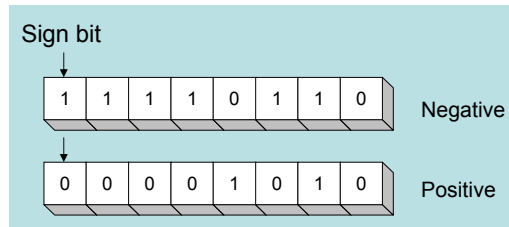
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Sign Bit

- ❖ Highest bit indicates the sign

- ❖ 1 = negative

- ❖ 0 = positive



For Hexadecimal Numbers, check most significant digit

If highest digit is > 7, then value is negative

Examples: 8A and C5 are negative bytes

B1C42A00 is a negative word (32-bit signed integer)

Sign Extension

Step 1: Move the number into the lower-significant bits

Step 2: Fill all the remaining higher bits with the sign bit

- ❖ This will ensure that both magnitude and sign are correct

- ❖ Examples

- ❖ Sign-Extend 10110011 to 16 bits

10110011 = -77 → 11111111 10110011 = -77

- ❖ Sign-Extend 01100010 to 16 bits

01100010 = +98 → 00000000 01100010 = +98

- ❖ Infinite 0s can be added to the left of a positive number

- ❖ Infinite 1s can be added to the left of a negative number

Two's Complement of a Hexadecimal

❖ To form the two's complement of a hexadecimal

- ❖ Subtract each hexadecimal digit from 15
- ❖ Add 1

❖ Examples:

2's complement of **6A3D** = **95C2** + **1** = **95C3**

2's complement of **92F15AC0** = **6D0EA53F** + **1** = **6D0EA540**

2's complement of **FFFFFFF** = **0000000** + **1** = **0000001**

❖ No need to convert hexadecimal to binary

Binary Subtraction

❖ When subtracting $A - B$, convert B to its 2's complement

❖ Add A to $(-B)$

$$\begin{array}{r}
 \text{borrow: } 1 \quad 1 \quad 1 \\
 01001101 \\
 - 00111010 \\
 \hline
 00010011
 \end{array}
 \quad \Rightarrow \quad
 \begin{array}{r}
 \text{carry: } 1 \quad 1 \quad 1 \quad 1 \\
 01001101 \\
 + 11000110 \text{ (2's complement)} \\
 \hline
 00010011 \text{ (same result)}
 \end{array}$$

❖ Final carry is ignored, because

- ❖ Negative number is sign-extended with 1's
- ❖ You can imagine infinite 1's to the left of a negative number
- ❖ Adding the carry to the extended 1's produces extended zeros

Hexadecimal Subtraction

		16 + 5 = 21		
Borrow:	1 1	1	↓	Carry:
	1 1 1 1			
-	B 1 4 F C 6 7 5	→		+
	8 3 9 E A 2 4 7			B 1 4 F C 6 7 5
	-----			+ 7 C 6 1 5 D B 9 (2's complement)
	2 D B 1 2 4 2 E			-----
				2 D B 1 2 4 2 E (same result)

- ❖ When a borrow is required from the digit to the left, then
Add 16 (decimal) to the current digit's value
- ❖ Last Carry is ignored

Ranges of Signed Integers

For n -bit signed integers: Range is -2^{n-1} to $(2^{n-1} - 1)$

Positive range: 0 to $2^{n-1} - 1$

Negative range: -2^{n-1} to -1

Storage Type	Unsigned Range	Powers of 2
Byte	-128 to +127	-2^7 to $(2^7 - 1)$
Half Word	-32,768 to +32,767	-2^{15} to $(2^{15} - 1)$
Word	-2,147,483,648 to +2,147,483,647	-2^{31} to $(2^{31} - 1)$
Double Word	-9,223,372,036,854,775,808 to +9,223,372,036,854,775,807	-2^{63} to $(2^{63} - 1)$

Practice: What is the range of signed values that may be stored in 20 bits?

Carry and Overflow

- ❖ Carry is important when ...
 - ❖ Adding or subtracting **unsigned integers**
 - ❖ Indicates that the **unsigned sum** is out of range
 - ❖ Either < 0 or $>$ maximum unsigned n -bit value
- ❖ Overflow is important when ...
 - ❖ Adding or subtracting **signed integers**
 - ❖ Indicates that the **signed sum** is out of range
- ❖ Overflow occurs when
 - ❖ Adding two positive numbers and the sum is negative
 - ❖ Adding two negative numbers and the sum is positive
 - ❖ Can happen because of the fixed number of sum bits

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Carry and Overflow Examples

- ❖ We can have carry without overflow and vice-versa
- ❖ Four cases are possible (Examples are 8-bit numbers)

$ \begin{array}{r} 1 \\ 00001111 \\ + 00001000 \\ \hline 00010111 \\ \hline \text{Carry} = 0 \quad \text{Overflow} = 0 \end{array} $	<p style="text-align: right;">15 8 23</p>
---	---

$ \begin{array}{r} 11111 \\ 00001111 \\ + 11111000 \\ \hline 00000111 \\ \hline \text{Carry} = 1 \quad \text{Overflow} = 0 \end{array} $	<p style="text-align: right;">15 248 (-8) 7</p>
--	---

$ \begin{array}{r} 1 \\ 01001111 \\ + 01000000 \\ \hline 10001111 \\ \hline \text{Carry} = 0 \quad \text{Overflow} = 1 \end{array} $	<p style="text-align: right;">79 64 143 (-113)</p>
---	--

$ \begin{array}{r} 1 1 \\ 11011010 \\ + 10011101 \\ \hline 01110111 \\ \hline \text{Carry} = 1 \quad \text{Overflow} = 1 \end{array} $	<p style="text-align: right;">218 (-38) 157 (-99) 119</p>
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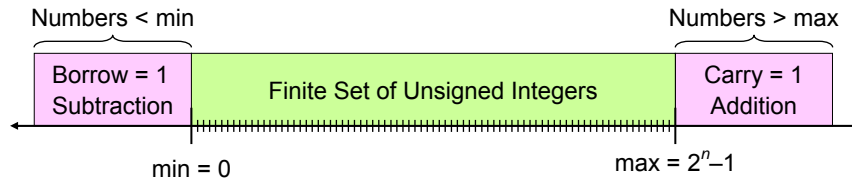
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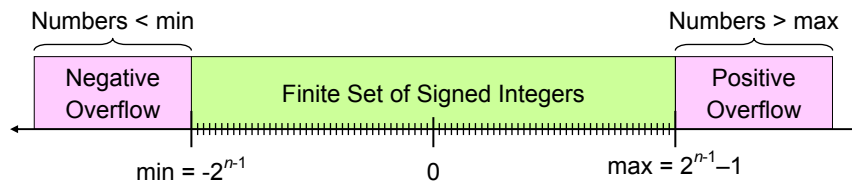
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Range, Carry, Borrow, and Overflow

❖ Unsigned Integers: n -bit representation



❖ Signed Integers: n -bit 2's complement representation



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Character Storage

❖ Character sets

- ❖ Standard ASCII: 7-bit character codes (0 – 127)
- ❖ Extended ASCII: 8-bit character codes (0 – 255)
- ❖ Unicode: 16-bit character codes (0 – 65,535)
- ❖ Unicode standard represents a universal character set
 - Defines codes for characters used in all major languages
 - Used in Windows-XP: each character is encoded as 16 bits
- ❖ UTF-8: variable-length encoding used in HTML
 - Encodes all Unicode characters
 - Uses 1 byte for ASCII, but multiple bytes for other characters

❖ Null-terminated String

- ❖ Array of characters followed by a NULL character

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Printable ASCII Codes

	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
2	space	!	"	#	\$	%	&	'	()	*	+	,	-	.	/
3	0	1	2	3	4	5	6	7	8	9	:	;	<	=	>	?
4	@	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
5	P	Q	R	S	T	U	V	W	X	Y	Z	[\]	^	_
6	`	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o
7	p	q	r	s	t	u	v	w	x	y	z	{		}	~	DEL

❖ Examples:

- ❖ ASCII code for space character = 20 (hex) = 32 (decimal)
- ❖ ASCII code for 'L' = 4C (hex) = 76 (decimal)
- ❖ ASCII code for 'a' = 61 (hex) = 97 (decimal)

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Control Characters

- ❖ The first 32 characters of ASCII table are used for control
- ❖ Control character codes = 00 to 1F (hexadecimal)
 - ❖ Not shown in previous slide
- ❖ Examples of Control Characters
 - ❖ Character 0 is the **NULL** character ⇒ used to terminate a string
 - ❖ Character 9 is the **Horizontal Tab (HT)** character
 - ❖ Character 0A (hex) = 10 (decimal) is the **Line Feed (LF)**
 - ❖ Character 0D (hex) = 13 (decimal) is the **Carriage Return (CR)**
 - ❖ The LF and CR characters are used together
 - They advance the cursor to the beginning of next line
- ❖ One control character appears at end of ASCII table
 - ❖ Character 7F (hex) is the **Delete (DEL)** character

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