King Fahd University of Petroleum & Minerals Computer Engineering Dept

COE 200 – Fundamentals of Computer Engineering

Term 043

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Binary Logic

- Deals with *binary* variables that take one of two discrete values
- Values of variables are called by a variety of very different names
  - *high* or *low* based on voltage representations in electronic circuits
  - *true* or *false* based on their usage to represent logic states
  - *one* (1) or *zero* (0) based on their values in Boolean algebra
  - *open* or *closed* based on its operation in *gate* logic
  - *on* or *off* based on its operation in *switching* logic
  - *asserted* or *de-asserted* based on its effect in digital systems

Basic Operations - AND

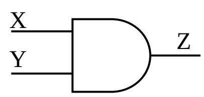
- Another Symbol is ".", e.g.

  Z = X AND Y or

  Z = X.Y or even

  Z = XY
- X and Y are inputs, Z is an output
- Z is equal to 1 if and only if X = 1 and Y = 1; Z = 0 otherwise (similar to the multiplication operation)
- Truth Table:
- Graphical symbol:

X	Y	Z=XY
0	0	0
0	1	0
1	0	0
1	1	1



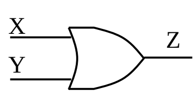
Basic Operations - OR

- Another Symbol is "+", e.g.

  Z = X OR Y or

  Z = X + Y
- X and Y are inputs, Z is an output
- Z is equal to 0 if and only if X = 0 and Y = 0; Z = 1 otherwise (similar to the addition operation)
- Truth Table:
- Graphical symbol:

X	Y	Z=X+Y
0	0	0
0	1	1
1	0	1
1	1	1



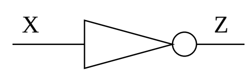
Basic Operations - NOT

- Another Symbol is "‾", e.g.

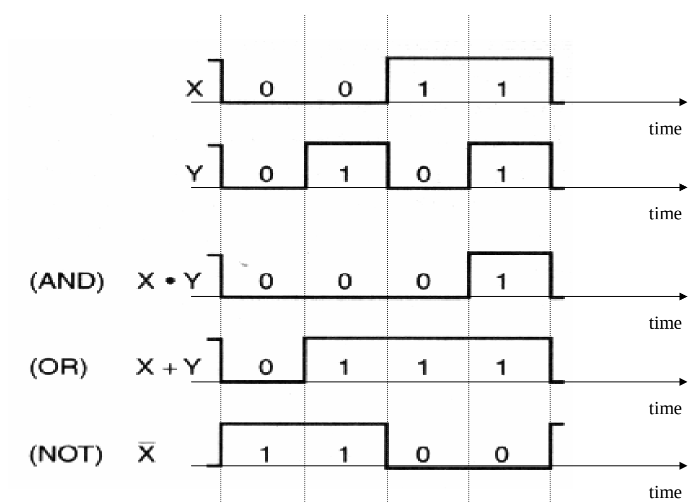
  $Z = \overline{X}$ or $Z = X'$
- X is the input, Z is an output
- Z is equal to 0 if X = 1; Z = 1 otherwise
- Sometimes referred to as the complement or invert operation
- Truth Table:

X	Z=X′
0	1
1	0

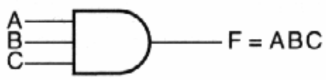
- Graphical symbol:



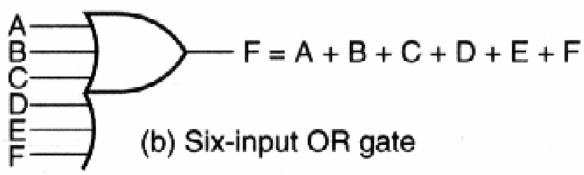
Time Diagrams



Multiple Input Gates



(a) Three-input AND gate



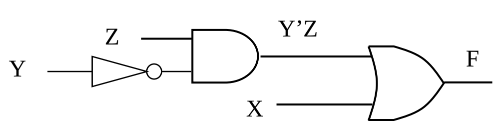
(b) Six-input OR gate

Boolean Algebra

- Consider the following *function*, F

$$F = X + Y'Z$$

- The function F is referred to as a BOOLEAN FUNCTION
- F has two terms: X and Y'Z
- The circuit diagram for F is as shown below



Boolean Algebra – cont'd

- The truth table for F is as follows
- Note:
  - In general, a truth table for an n-variable function, has 2^n rows to cover all possible input combinations
  - The table covers all possible combinations of the inputs
  - To arrive at the F's column one could use an Y'Z column as follows

X	Y	Z	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

The Y'Z column is computed using the Y and Z columns and then using the columns X and Y'Z, the column F is computed
The column Y'Z is **not** an essential part of truth table

X	Y	Z	Y'Z	F
0	0	0	0	0
0	0	1	1	1
0	1	0	0	0
0	1	1	0	0
1	0	0	0	1
1	0	1	1	1
1	1	0	0	1
1	1	1	0	1

Basic Identities

- For the AND operation

  X . 1 = X

  X . 0 = 0

  X . X = X

  X . X' = 0

- For the OR operation

  X + 0 = X

  X + 1 = 1

  X + X = X

  X + X' = 1

- For the NOT operation

  X'' = X

Notice the **duality**: start with one identity
-Replace the AND by OR
-Replace the 1 by 0 and vice versa
You end up with an identity from the other group

Basic Identities (2)

- For the AND operation — OR Operation

	AND	OR
Commutative:	X.Y=Y.X	X+Y=Y+X
Associative:	X(YZ)=(XY)Z	X+(Y+Z)=(X+Y)+Z
Distributive:	X+YZ=(X+Y)(X+Z)	X(Y+Z)=(XY)+(XZ)
DeMorgan's:	(X.Y)′=X′+Y′	(X+Y)′=X′.Y′

- All above properties can be generalized to n > 2 variables: e.g:
  - $(X_1+X_2+\ldots+X_n)' = X_1'.X_2'. \ldots .X_n'$, or
  - $(X_1.X_2. \ldots .X_n)' = X_1'+X_2'+ \ldots +X_n'$

Verifying Basic Identities

- Any identity (not only the basic ones) can be verified using the truth table
- Example: verify that (X+Y)′ = X′Y′

X	Y	X+Y	(X+Y)′	X′	Y′	X′.Y′
0	0	0	1	1	1	1
0	1	1	0	1	0	0
1	0	1	0	0	1	0
1	1	1	0	0	0	0

ALL possible combinations of inputs

Some columns to aid in calculations

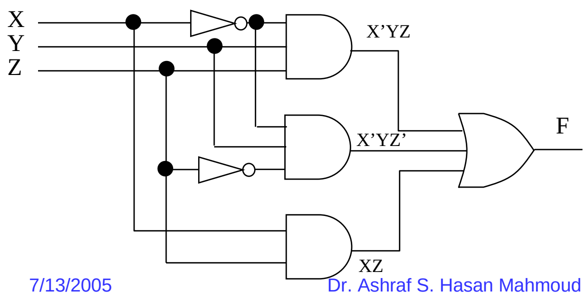
The two quantities in question

Algebraic Manipulation - Example

- Consider the following function, F

$$F = X'YZ + X'YZ' + XZ$$

- The function can be implemented using above expressions as in



We need:
-2 inverters
-3 AND gates
-1 OR gate

Algebraic Manipulation – Example – cont'd

- The function

$$F = X'YZ + X'YZ' + XZ$$

can be simplified "ALGEBRAICALLY" as follows:

F = X'YZ + X'YZ' + XZ
= X'Y(Z + Z') + XZ → by the distributive property
= X'Y(1) + XZ → by the properties of the OR operation
= X'Y + XZ → by the properties of the AND operation

- Therefore F can be written as

$$F = X'Y + XZ$$

- Using this simpler form, one can implement the function as

Algebraic Manipulation – Example – cont'd

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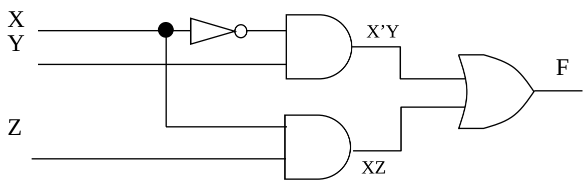
$$F = X'Y + XZ$$

- Using this simpler form, one can implement the function as



We need:
-1 inverters
-2 AND gates
-1 OR gate

Reduced hardware cost

- One can use the truth table method to show that F = X'YZ + X'YZ' + XZ is indeed equal to X'Y + XZ

More Notes on Function

- The function

$$F = X'Y + XZ$$

Can be written as

$$F(X,Y,Z) = X'Y + XZ$$

These are referred to as literals

These are referred to as terms

This is to emphasize the fact that the function has three inputs or variables

More Identities

- Page 37 in the text ➔ VERY IMPORTANT – make sure you can prove/verify all of these identities
- Listing
  1. X + XY = X
  2. XY + XY′ = X
  3. X + X′Y = X + Y
  4. X(X + Y) = X
  5. (X + Y)(X + Y′) = X
  6. X(X′ + Y) = XY
  7. XY + X′Z + YZ = XY + X′Z (the consensus theorem)

The proof/verification of these is in the textbook

More Identities - continued

- Using the duality principle (refer to slide XX) there are other equivalent 7 identities
- Example: The proof of the consensus theorem is as follows

The RHS
= XY + X′Z + YZ
= XY + X′Z + YZ(X + X′)
= XY + X′Z + XYZ + X′YZ
= XY + XYZ + X′Z + X′YZ
= XY(1+Z) + X′Z(1+Y)
= XY + X′Z
= LHS

- The dual of the consensus theorem is given by

(X+Y)(X′+Z)(Y+Z) = (X+Y)(X′+Z)

Complement of a Function

- Using the truth table – complementing F means replacing every 0 with 1 and every 1 with 0 in the F column
- Algebraically, complementing F one can use DeMorgan's rule or the duality principle
- To use the duality principle
  - Replace Each AND with an OR and each OR with an AND
  - Complement each variable and constant

Example

- **Problem:** Find the complement of each of the following two functions $F_1 = X'YZ' + X'Y'Z$, and $F_2 = X(Y'Z' + YZ)$
- **Solution:**

For F_1, applying DeMorgan's rule as many times as necessary

$$F_1' = (X'YZ' + X'Y'Z)'$$
$$= (X'YZ')' \,.\, (X'Y'Z)'$$
$$= (X + Y' + Z) \,.\, (X + Y + Z')$$

Similarly for F_2:

$$F_2' = (X(Y'Z' + YZ))'$$
$$= X' + (Y'Z' + YZ)'$$
$$= X' + (Y'Z')' \,.\, (YZ)'$$
$$= X' + (Y + Z) \,.\, (Y' + Z')$$

Examples

- **Problem 2-2**: Prove the identity of each of the following Boolean equations, using algebraic manipulations.
  a) X'Y' + X'Y + XY = X' + Y
  b) A'B + B'C' + AB + B'C = 1
- **Solution:**

a) LHS = X'Y' + X'Y + XY
= X'Y' + X'Y + X'Y + XY
= X'(Y'+Y) +Y(X + X')
= X' + Y
= RHS

b) LHS = A'B + B'C' + AB + B'C
= (A'+A)B + B'(C'+C)
= B + B'
= 1
= RHS

Examples

- **Problem 2-6**: Simplify the following Boolean expressions to a minimum number of literals:
  a) ABC + ABC' + A'B
  e) (A+B'+AB')(AB+A'C+BC)
- **Solution:**

a) Expression = ABC + ABC' + A'B
= AB(C + C') + A'B
= (A+A')B
= B

e) Expression = (A+B'+AB')(AB+A'C+BC)
= (A+(1+A)B')(AB + A'C)
= (A+B')(AB+A'C)
= A(AB+A'C) + B'(AB+A'C)
= AB + A'B'C

Standard Forms of a Boolean Function

- A Boolean function can be written algebraically in a variety of ways
- Standard form: is an algebraic expression of the function that facilitates simplification procedures and frequently results in more desirable logic circuits (e.g. less number of gates)
- Standard form: contains product terms and sum terms
  - Product term: X'Y'Z
  - Sum term: X + Y + Z'

Standard Forms of a Boolean Function – cont'd

- A minterm: a product term in which all variables (or literals) of the function appear exactly once
- A maxterm: a sum term in which all the variables (or literals) of the function appear exactly once

- Example: for the function F(X,Y,Z),
  - the term X'Y is not a minterm, but XYZ' is a minterm
  - The term X'+Z is not a maxterm, but X+Y'+Z' is maxterm

More on Minterms and Maxterms

- A function of n variables – have 2^n possible minterms and 2^n possible maxterms
-

Naming Convention for Minterms

- Consider a function F(X, Y)

X	Y	Product Terms	Symbol	m_0	m_1	m_2	m_3
0	0	X'Y'	m_0	1	0	0	0
0	1	X'Y	m_1	0	1	0	0
1	0	XY'	m_2	0	0	1	0
1	1	XY	m_3	0	0	0	1

Variable complemented if 0
Variable not complemented if 1

m_i indicated the i^{th} minterm
For each binary combination of X and Y there is a minterm
The index of the minterm is specified by the binary combination
m_i is equal to 1 for ONLY THAT combination

Naming Convention for Maxterms

- Consider a function F(X, Y)

X	Y	Sum Terms	Symbol	M_0	M_1	M_2	M_3
0	0	X+Y	M_0	0	1	1	1
0	1	X+Y’	M_1	1	0	1	1
1	0	X’+Y	M_2	1	1	0	1
1	1	X’+Y’	M_3	1	1	1	0

Variable complemented if 1
Variable not complemented if 0

M_i indicated the i^{th} maxterm
For each binary combination of X and Y there is a maxterm
The index of the maxterm is specified by the binary combination
M_i is equal to 0 for ONLY THAT combination

More on Minterms and Maxterms

- In general, a function of n variables has
  - 2^n minterms: $m_0, m_1, \ldots, m_{2^n-1}$
  - 2^n maxterms: $M_0, M_1, \ldots, M_{2^n-1}$

- $m_i' = M_i$ or $M_i' = m_i$

Example: for F(X,Y):

$m_2 = XY' \Rightarrow m_2' = X'+Y = M_2$

More on Minterms and Maxterms – cont'd

- A Boolean function can be expressed algebraically from a give truth table by forming the logical sum of ALL the minterms that produce 1 in the function

X	Y	Z	m	F
0	0	0	m_0	1
0	0	1	m_1	0
0	1	0	m_2	1
0	1	1	m_3	0
1	0	0	m_4	0
1	0	1	m_5	1
1	1	0	m_6	0
1	1	1	m_7	1

- **Example**:

Consider the function defined by the truth table

F(X,Y,Z) ➔ 3 variables ➔ 8 minterms

F can be written as

$F = X'Y'Z' + X'YZ' + XY'Z + XYZ$, or

$= m_0 + m_2 + m_5 + m_7$

$= \Sigma m(0,2,5,7)$

More on Minterms and Maxterms – cont'd

- A Boolean function can be expressed algebraically from a give truth table by forming the logical product of ALL the maxterms that produce 0 in the function

X	Y	Z	M	F	F'
0	0	0	M_0	1	0
0	0	1	M_1	0	1
0	1	0	M_2	1	0
0	1	1	M_3	0	1
1	0	0	M4	0	1
1	0	1	M_5	1	0
1	1	0	M_6	0	1
1	1	1	M_7	1	0

- **Example**:

Consider the function defined by the truth table

F(X,Y,Z) ➔ in a manner similar to the previous example, F' can be written as

$F' = m_1 + m_3 + m_4 + m_6$

$= \Sigma m(1,3,4,6)$

Now apply DeMorgan's rule

$F = F'' = [m_1 + m_3 + m_4 + m_6]'$

$= m_1'.m_3'.m_4'.m_6'$

$= M_1.M_3.M_4.M_6$

$= \Pi M(1,3,4,6)$

Note the indices in this list are those that are missing from the previous list in $\Sigma m(0,2,5,7)$

Summary

- A Boolean function can be expressed algebraically as:
  - The logical sum of minterms
  - The logical product of maxterms
- Given the truth table, writing F as
  - Σm_i – for all minterms that produce 1 in the table, or
  - ΠM_i – for all maxterms that produce 0 in the table
- Another way to obtain the Σm_i or ΠM_i is to use ALGEBRA – see next example

Example:

- Write E = Y' + X'Z' in the form of Σm_i and ΠM_i?

- Solution: **Method1**

First construct the Truth Table as shown

Second:

E = Σm(0,1,2,4,5), and

E = ΠM(3,6,7)

X	Y	Z	m	M	E
0	0	0	m_0	M_0	1
0	0	1	m_1	M_1	1
0	1	0	m_2	M_2	1
0	1	1	m_3	M_3	0
1	0	0	m_4	M4	1
1	0	1	m_5	M_5	1
1	1	0	m_6	M_6	0
1	1	1	m_7	M_7	0

Example: cont'd

- Solution: **<u>Method2_a</u>**

E = Y' + X'Z'
= Y'(X+X')(Z+Z') + X'Z'(Y+Y')
= (XY'+X'Y')(Z+Z') + X'YZ'+X'Z'Y'
= XY'Z+X'Y'Z+XY'Z'+X'Y'Z'+ X'YZ'+X'Z'Y'
= $m_5 + m_1 + m_4 + m_0 + m_2 + m_0$
= $m_0 + m_1 + m_2 + m_4 + m_5$
= $\Sigma m(0,1,2,4,5)$

To find the form ΠMi, consider the remaining indices

E = $\Pi M(3,6,7)$

- Solution: **<u>Method2_b</u>**

E = Y' + X'Z'
E' = Y(X+Z)
= YX + YZ
= YX(Z+Z') + YZ(X+X')
= XYZ+XYZ'+X'YZ
E = (X'+Y'+Z')(X'+Y'+Z)(X+Y'+Z')
= $M_7 . M_6 . M_3$
= $\Pi M(3,6,7)$

To find the form Σm_i, consider the remaining indices

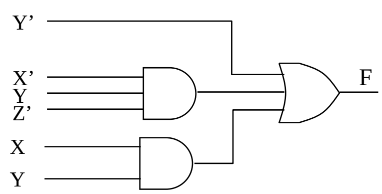
E = $\Sigma m(0,1,2,4,5)$

Exercise

- What is $G(X,Y) = \Sigma m(0,1,2,3)$ equal to?

Implementation – Sum of Products

- Consider F = Y' + X'YZ' + XY
  - Three products: Y' (one literal), X'YZ' (three literals), and XY (two literals)
- The logic diagram



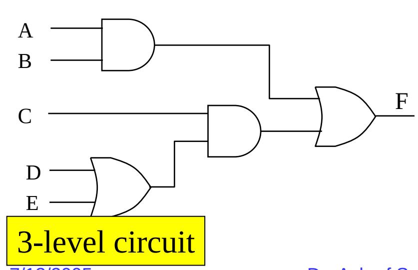
- Two-level implementation:
  - AND-OR
- Each product term requires an AND gate (except one literal terms)
- Logic diagram requires ONE OR gate

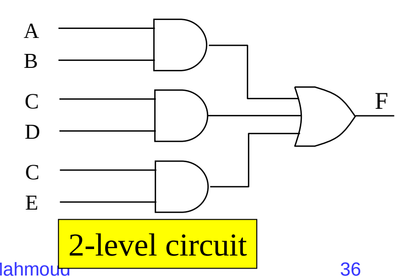
Implementation – Sum of Products – cont'd

- Consider F = AB + C(D+E)
- This expression is NOT in the sum-of-products form
- Use the identities/algebraic manipulation to convert to a standard form (sum of products), as in

F = AB + CD + CE

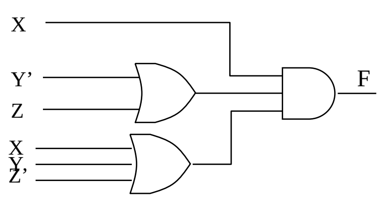
- Logic Diagrams:





Implementation – Product of Sums

- Consider F = X(Y′+Z)(X+Y+Z′)
- This expression is in the product-of-sums form:
  - Thee summation terms: X (one literal), Y′+Z (two literals), and X+Y+Z′ (three literals)
- Logic Diagrams:



- Two-level implementation:
  - OR-AND
- Each sum term requires an OR gate (except one literal terms)
- Logic diagram requires ONE AND gate

Examples:

- Problem 2-10b: Obtain the truth table of the following function and express each function in sum-of-minterms and product-of-maximterms form: (A′+B)(B′+C)

- Solution:

Let F(A,B,C) = (A′+B)(B′+C)

The truth table is as shown in figure

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

F(A,B,C) = A′B′C′+A′B′C+A′BC+ABC
$= \Sigma m(0,1,3,7)$

F(A,B,C) = (A+B′+C)(A′+B+C)(A′+B+C′)(A′+B′+C)
$= \Pi M(2,4,5,6)$