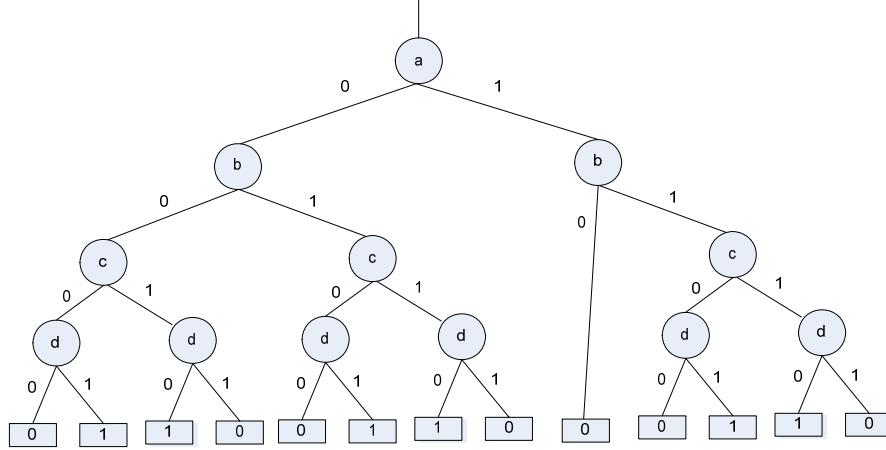


COE 561, Term 081
Digital System Design and Synthesis

HW# 1

Due date: Tuesday, Nov. 11

- Q.1.** Consider the following OBDD with the variable ordering $\{a, b, c, d\}$. Reduce it based on **Reduce** function to obtain the ROBDD. Show the details of your work.



- Q.2.** Consider the functions $f_1 = ab + ac + bc$, $f_2 = a(b \oplus c) + bc$ and $f_3 = a(a \oplus b)' + c(a \oplus b)$:
- Draw the **ROBDD** for the functions **f1**, **f2** and **f3** using the variable order $\{a, b, c\}$.
 - What do you conclude from the results obtained in (i).
- Q.3.** Consider the two functions $f = a \oplus b \oplus c$ and $g = ab + ac + bc$.
- Compute the function $f \oplus g$.
 - Draw the **ITE DAG** for the function $f \oplus g$. Show the details of the ITE algorithm step by step. Use the variable ordering $\{a, b, c\}$
- Q.4.** Consider the following given matrix representing a covering problem:

$$A = \begin{bmatrix} 1 & 0 & 1 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 \end{bmatrix}$$

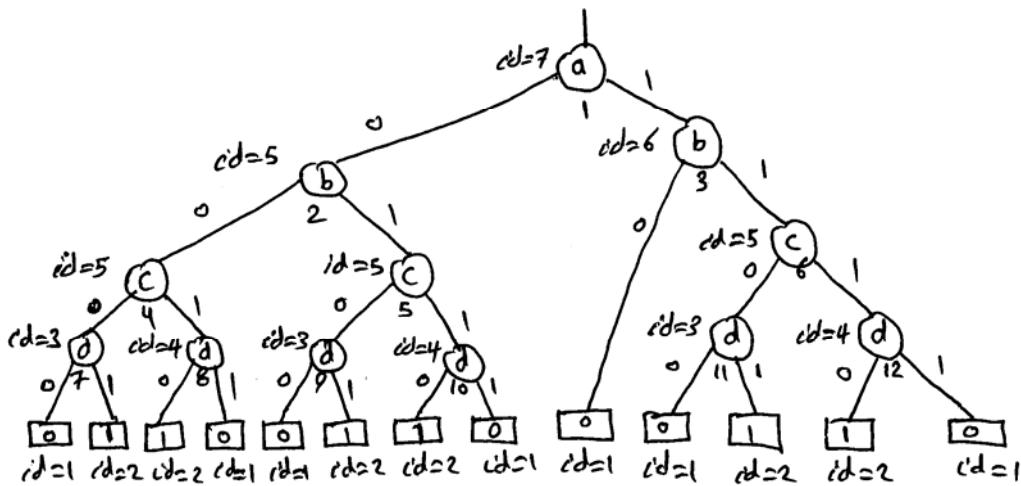
Find a **minimum cover** using **EXACT_COVER** procedure. Show all the details of the algorithm. Assume the following order in branching selection when needed: C₁, C₂, C₃, C₄, C₅, C₆.

- Q.5.** Consider the function $F(A, B, C) = AB + \overline{A}C + \overline{B}\overline{C}$.
- (i) Represent the function using **positional cube notation**.
 - (ii) Using positional cube notation, compute the **cofactor** F_A.
 - (iii) Using positional cube notation, compute the **consensus** between the two cubes \overline{AC} and \overline{BC} .
 - (iv) Using positional cube notation, based on the **sharp** operation, compute the complement of the function F.
 - (v) Using positional cube notation, determine if the cube BC is **covered** by the function $F = AB + \overline{A}C + \overline{B}\overline{C}$.
- Q.6.** Consider the function $F(A, B, C, D) = \overline{AC} + A\overline{B} + \overline{AB}C + \overline{AC}\overline{D}$:
- (i) Compute the **complement** of the function using the recursive complementation procedure outlined in section 7.3.4.
 - (ii) Compute all the **prime implicants** of the function using the method outlined in section 7.3.4.

Note that you do not need to use the positional cube notation in your solution of this question.

HW #1

Q1.



First, we set $id(v) = 1$ for all leaf vertices with value 0 and $id(v) = 2$ for leaf vertices with value 1. We initialize ROBDD with two leaf vertices for 0 and 1. Then, we process vertices at level 4, i.e. nodes with index = 4.

$$V = \{7, 8, 9, 10, 11, 12\}.$$

None of the vertices is removed since $id(\text{low}(v)) \neq id(\text{high}(v))$.

We assign keys to all vertices $\in V$.

$$\text{Key}(7) = (1, 2), \text{ Key}(8) = (2, 1), \text{ Key}(9) = (1, 2),$$

$$\text{Key}(10) = (2, 1), \text{ Key}(11) = (1, 2), \text{ Key}(12) = (2, 1).$$

$$\text{oldKey} = (0, 0).$$

We next sort the vertices in V according to their keys. Thus, $V = \{7, 9, 11, 8, 10, 12\}$.

$v = \{7\}$: Since $\text{key}(7) \neq \text{oldKey}$, $\text{nextId} = 3$, $\text{id}(7) = 3$ $\text{oldKey} = (1, 2)$. We add $v = \{7\}$ to the ROBDD.

$v = \{9\}$: Since $\text{Key}(9) = \text{oldKey}$, $\text{id}(9) = 3$.
 $v = \{11\}$: Since $\text{Key}(11) = \text{oldKey}$, $\text{id}(11) = 3$.
 $v = \{8\}$: Since $\text{Key}(8) \neq \text{oldKey}$, $\text{nextId} = 4$. $\text{Key}(8) = 4$
 $\text{oldKey} = (2, 1)$. We add $v = \{8\}$ to the RABDD.

$v = \{10\}$: Since $\text{Key}(10) = \text{oldKey}$, $\text{id}(10) = 4$.
 $v = \{12\}$: Since $\text{Key}(12) = \text{oldKey}$, $\text{id}(12) = 4$.

Next, we process vertices at level 3 with index = c.

$$V = \{4, 5, 6\}$$

None of the vertices is removed since $\text{id}(\text{low}(v)) \neq \text{id}(\text{high}(v))$.
 We assign keys to all vertices $\in V$.

$$\text{Key}(4) = (3, 4), \text{Key}(5) = (3, 4), \text{Key}(6) = (3, 4).$$

$$\text{oldKey} = (0, 0)$$

We sort the vertices according to their keys.

$$V = \{4, 5, 6\}.$$

$v = \{4\}$: Since $\text{Key}(4) \neq \text{oldKey}$, $\text{nextId} = 5$, $\text{id}(4) = 5$,
 $\text{oldKey} = (3, 4)$. We add $v = \{4\}$ to the RABDD.

$v = \{5\}$: Since $\text{Key}(5) = \text{oldKey}$, $\text{id}(5) = 5$.

$v = \{6\}$: Since $\text{Key}(6) = \text{oldKey}$, $\text{id}(6) = 5$.

Next, we process vertices at level 2 with index = b.

$$V = \{2, 3\}.$$

Since $\text{id}(\text{low}(2)) = \text{id}(\text{high}(2)) = 5$, $\text{id}(2) = 5$ and
 vertex 2 is removed from V.

$$\text{Key}(3) = (1, 5), \text{oldKey} = (0, 0)$$

Since $\text{Key}(3) \neq \text{oldKey}$, $\text{nextId} = 6$, $\text{id}(3) = 6$,
 $\text{oldKey} = (1, 5)$. We add $v = \{3\}$ to the RABDD.

Finally, we process vertices at level 1 with index = a.

$$V = \{1\}.$$

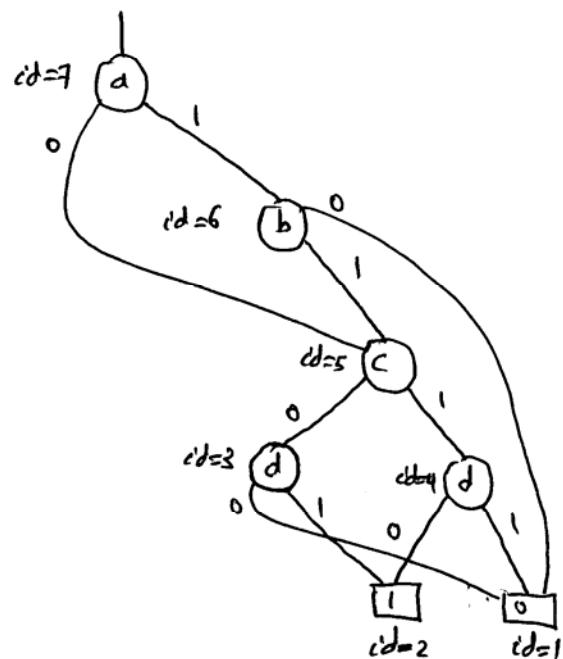
Since $id(\text{low}(1)) \neq id(\text{high}(1))$, the vertex is not removed.

$\text{key}(1) = (5, 6)$, $\text{oldKey} = (0, 0)$.

since $\text{key}(1) \neq \text{oldKey}$, $\text{nextId} = 7$, $\text{id}(1) = 7$.

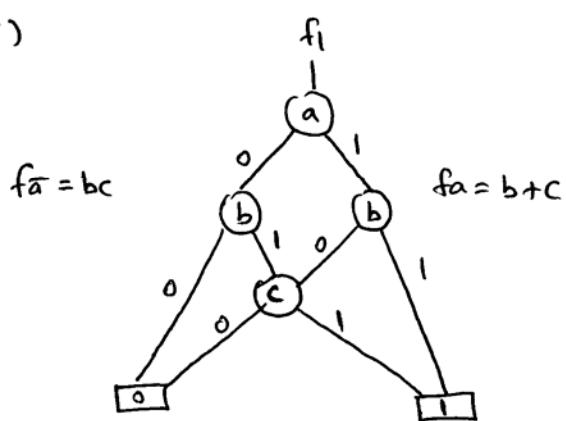
we add $v = \{1\}$ to the ROBDD.

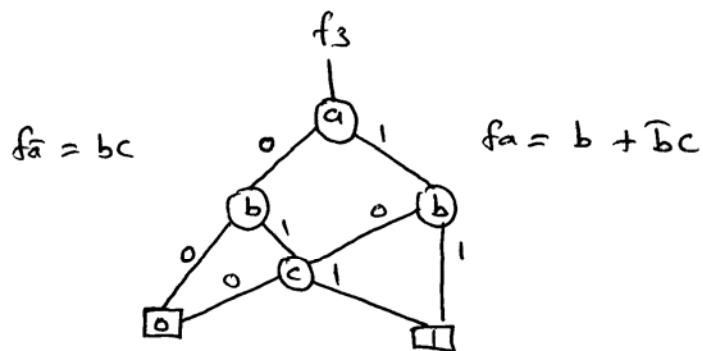
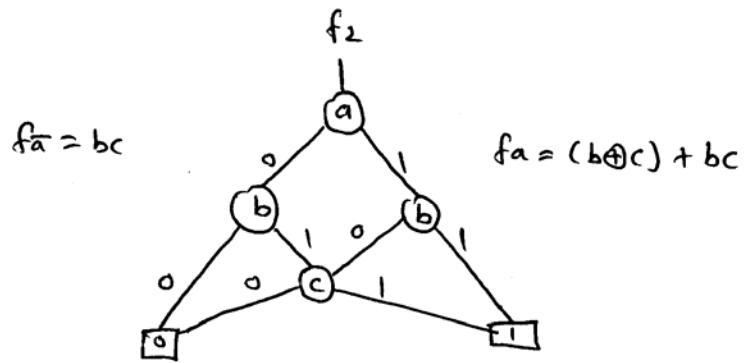
Thus, the ROBDD formed is:



Q₂.

(1)





(ii) Since we have obtained equivalent KOBDS for the three functions, we conclude that $f_1 = f_2 = f_3$.

Q3

(i) We first express the two functions f and g using the orthonormal basis: $\bar{a}\bar{b}$, $\bar{a}b$, $a\bar{b}$, ab

$$f = \bar{a}\bar{b}(c) + \bar{a}b(\bar{c}) + a\bar{b}(\bar{c}) + ab(c)$$

$$g = \bar{a}\bar{b}(0) + \bar{a}b(c) + a\bar{b}(c) + ab(1)$$

$$f \oplus g = \bar{a}\bar{b}(c \oplus 0) + \bar{a}b(\bar{c} \oplus c) + a\bar{b}(\bar{c} \oplus c) + ab(c \oplus 1)$$

$$= \bar{a}\bar{b}(c) + \bar{a}b(1) + a\bar{b}(1) + ab(\bar{c})$$

$$= \bar{a}\bar{b}c + \bar{a}b + a\bar{b} + ab\bar{c}$$

(ii) ITE diagram for the function $f \oplus g$

$$f \oplus g = \text{ITE}(f, \bar{g}, g)$$

$$= \text{ITE}(a \oplus b \oplus c, \overline{ab+ac+bc}, ab+ac+bc)$$

$$- x = a$$

$$t = \text{ITE}(\overline{b \oplus c}, \overline{b+c}, b+c)$$

$$- x = b$$

$$t = \text{ITE}(c, 0, 1) = \bar{c} \quad (\text{trivial case})$$

$$\text{we assign } rd = 3 \Rightarrow t = 3$$

$$e = \text{ITE}(\bar{c}, \bar{c}, c) = 1 \quad (\text{trivial case})$$

$$\Rightarrow e = 2$$

since $t \neq e$, an entry will be added in the table for $(b, 3, 2)$ with $rd = 4$

$$t = 4$$

$$e = \text{ITE}(b \oplus c, \overline{bc}, bc)$$

$$- x = b$$

$$t = \text{ITE}(\bar{c}, \bar{c}, c) = 1 \quad (\text{trivial case})$$

$$\Rightarrow t = 2$$

$$e = \text{ITE}(c, 1, 0) = c \quad (\text{trivial case})$$

$$\text{we assign } rd = 5 \Rightarrow e = 5$$

since $t \neq e$, an entry will be added in the table for $(b, 2, 5)$ with $rd = 6$

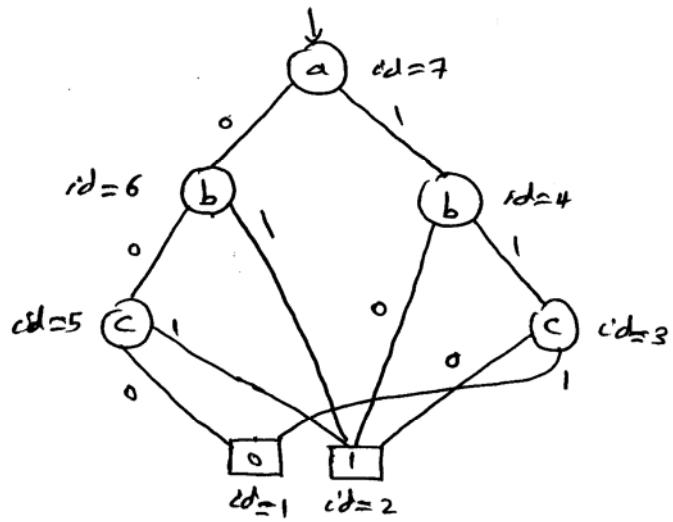
$$e = 6$$

since $t \neq e$, an entry will be added in the table for $(a, 4, 6)$, with $rd = 7$.

Thus, the unique table produced is :

id	var	right child	left child
3	c	1	2
4	b	3	2
5	c	2	1
6	b	2	5
7	a	4	6

The corresponding ITE DAG is :



Q4. The matrix to be covered :

	c_1	c_2	c_3	c_4	c_5	c_6
r_1	1	0	1	1	1	0
r_2	1	0	0	0	0	1
r_3	0	1	0	1	1	0
r_4	0	1	0	1	0	0
r_5	1	1	0	1	0	1
r_6	0	1	1	0	1	0

Initially, there are no essential columns.

c_1 dominates c_6 and hence c_6 is removed

c_5 dominates c_3 and hence c_3 is removed

r_3 dominates r_4 and hence r_3 is removed

r_5 dominates r_2 and hence r_5 is removed

The reduced matrix is :

	c ₁	c ₂	c ₄	c ₅
r ₁	1	0	1	1
r ₂	1	0	0	0
r ₄	0	1	1	0
r ₆	0	1	0	1

We can see now that c₁ is essential and it has to be selected and r₁ and r₂ are removed since they are covered.

The obtained reduced matrix will be:

	c ₂	c ₄	c ₅
r ₄	1	1	0
r ₅	1	0	1

Now we see that c₂ dominates c₄ and it also dominates c₅. Thus, c₄ and c₅ are removed and c₂ becomes essential and it is selected.

c₂ becomes essential and it is selected.
Thus, the selected columns are c₁ and c₂ and the returned solution is $\alpha = (1, 1, 0, 0, 0, 0)$.

Note that in this problem, there was no need for branching as applying reduction techniques led to the minimum cover.

$$Q5. \quad F = AB + \overline{A}C + \overline{B}\overline{C}$$

(i) Positional cube notation:

	A	B	C
AB	01	01	11
\overline{AC}	10	11	01
\overline{BC}	11	10	10

(ii) Cofactor F_A

	A	B	C
B	11	01	11
\overline{BC}	11	10	10

$$F_A = B + \overline{BC}$$

(iii) Consensus between \overline{AC} and \overline{BC}

\overline{AC}	A	B	C
\overline{BC}	10	11	01
	11	10	00
	10	11	00
\overline{AB}	10	10	11

(iv) Complement of F using sharp operation

First, we compute $U \# AB$

	A	B	C
U	11	11	11
AB	01	01	11
	10	11	11
\overline{B}	11	10	11
	11	11	00

Then, we compute $\{\overline{A} + \overline{B}\} \# \overline{AC}$

\overline{A}	A	B	C		\overline{B}	A	B	C
\overline{AC}	10	11	11		\overline{AC}	11	10	11
\overline{AC}	10	11	01		\overline{AC}	10	11	01
	00	11	11	void	\overline{AB}	01	10	11
	10	00	11	void		11	00	11
\overline{AC}	10	11	10		\overline{BC}	11	10	10

Thus, we obtained the result $\{\overline{AC}, \overline{AB}, \overline{BC}\}$

Then, we perform $\{\overline{AC}, \overline{AB}, \overline{BC}\} \# \overline{BC}$

\overline{AC}	A	B	C		\overline{AB}	A	B	C		\overline{BC}	A	B	C
\overline{AC}	10	11	10		\overline{AB}	01	10	11		\overline{BC}	11	10	10
\overline{BC}	11	10	10		\overline{BC}	11	10	10		\overline{BC}	11	10	10
	00	11	10	void		00	10	11	void		00	10	10
\overline{ABC}	10	01	10			01	00	11	void		11	00	10
	10	11	00	void	\overline{ABC}	01	10	01			11	10	00

Thus, the obtained result is $\overline{ABC} + \overline{AB}C$

(v) To determine if BC is covered by F_j , we need to compute F_{BC} and check if it is tautology.

F_{BC} is :

A	B	C
01	11	11
10	11	11

which is tautology. Thus, BC is covered by F.

$$Q6. \quad F = \bar{A}\bar{C} + A\bar{B} + \bar{A}\bar{B}C + \bar{A}C\bar{D}$$

(i) complement of the function

we will expand on unate variables

$$\begin{aligned}
 F &= B F_B + F_{\bar{B}} \\
 &= B [\bar{A}\bar{C} + \bar{A}C\bar{D}] + [\bar{A}\bar{C} + A + \bar{A}C + \bar{A}C\bar{D}] \\
 &= B [A[0] + [\bar{C} + C\bar{D}]] \\
 &\quad + [D[\bar{A}\bar{C} + A + \bar{A}C] + \bar{D}[\bar{A}\bar{C} + A + \bar{A}C + \bar{A}C]] \\
 &= B [A[0] + [D[\bar{C}] + [\bar{C} + C]]] \\
 &\quad + [D[\bar{A}[\bar{C} + C] + A[1]] \\
 &\quad \quad + \bar{D}[\bar{A}[\bar{C} + C] + A[1]]] \\
 &= B [A[0] + [D[\bar{C}] + 1]] \\
 &\quad + [B[\bar{A}[1] + A[1]] \\
 &\quad \quad + \bar{B}[\bar{A}[1] + A[1]]] \\
 \text{thus, } \bar{F} &= B [A[1] + [D[C] + 0]] \\
 &\quad + [D[\bar{A}[0] + A[0]] \\
 &\quad \quad + \bar{D}[\bar{A}[0] + A[0]]] \\
 &= BA + BDC
 \end{aligned}$$

(ii) Prime implicants

$$F = \bar{A} [\bar{C} + \bar{B}C + C\bar{D}] + A [\bar{B}]$$

$$= \bar{A} [\bar{C}[1] + C[\bar{B} + \bar{D}]] + A[\bar{B}]$$

prime implicants for $\bar{F}\bar{A}\bar{C} = \{1\}$

prime implicants for $\bar{F}\bar{A}C = \{\bar{B}, \bar{D}\}$

prime implicants for $\bar{F}\bar{A} = \overset{\text{SCC}}{\{ \bar{C}, C\bar{B}, C\bar{D}, \bar{B}, \bar{D} \}}$
 $= \{ \bar{C}, \bar{B}, \bar{D} \}$

prime implicants for $F_A = \{ \bar{B} \}$

prime implicants for $F = \text{SCC} \{ A\bar{B}, \bar{A}\bar{C}, \bar{A}\bar{B}, \bar{A}\bar{D}, \bar{B}\bar{C}, \bar{B}, \bar{B}\bar{D} \}$
 $= \{ \bar{A}\bar{C}, \bar{A}\bar{D}, \bar{B} \}$