

William Stallings Data and Computer Communications

Chapter 2 Protocols and Architecture

Characteristics of Protocols

- ⌘ Direct or indirect
- ⌘ Monolithic or structured
- ⌘ Symmetric or asymmetric
- ⌘ Standard or nonstandard

Direct or Indirect

⌘ Direct

- ☑ Systems share a point to point link or
- ☑ Systems share a multi-point link
- ☑ Data can pass without intervening active agent

⌘ Indirect

- ☑ Switched networks or
- ☑ Internetworks or internets
- ☑ Data transfer depend on other entities

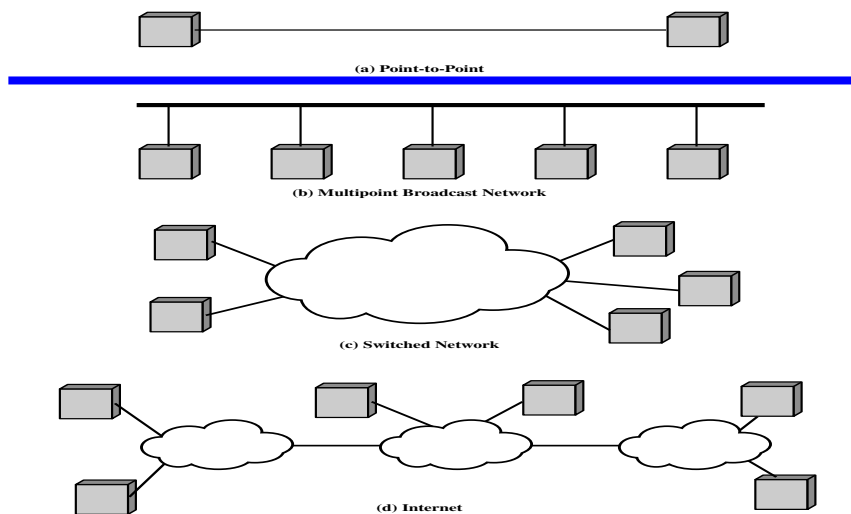


Figure 2.1 Means of Connection of Communicating Systems

Monolithic or Structured

- ⌘ Communication is a complex task
- ⌘ Too complex for single unit
- ⌘ Structured design breaks down problem into smaller units
- ⌘ A set of protocols with hierarchical or layered structure
- ⌘ Higher-level entities rely on lower-level entities to exchange data

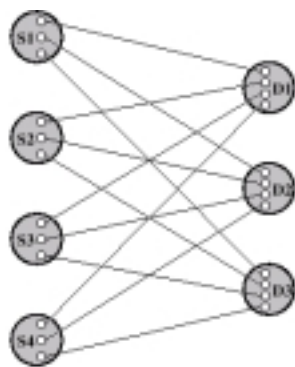
Symmetric or Asymmetric

- ⌘ Symmetric
 - ☒ Communication between peer entities
- ⌘ Asymmetric
 - ☒ May be dictated by the logic of an exchange e.g. client/server
 - ☒ Desire to keep one of the entities or systems simple
 - ☒ A computer that polls and selects a number of terminals

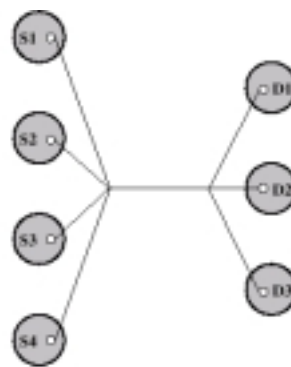
Standard or Nonstandard

- ⌘ Nonstandard protocols built for specific computers and tasks
- ⌘ K sources and L receivers leads to $K \cdot L$ protocols and $2 \cdot K \cdot L$ implementations
- ⌘ If common protocol used, $K + L$ implementations needed

Use of Standard Protocols



(a) Without standards: 12 different protocols;
24 protocol implementations



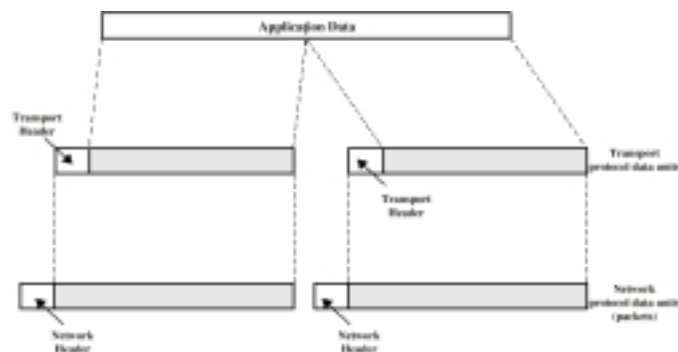
(a) With standards: 1 protocol;
7 implementations

Basic Functions of Protocols

- ⌘ Encapsulation
- ⌘ Segmentation and reassembly
- ⌘ Connection control
- ⌘ Ordered delivery
- ⌘ Flow control
- ⌘ Error control
- ⌘ Addressing
- ⌘ Multiplexing
- ⌘ Transmission services

Encapsulation

- ⌘ Addition of control information to data
 - ☑ Address information
 - ☑ Error-detecting code
 - ☑ Protocol control



Segmentation (Fragmentation)

- ⌘ Data blocks are of bounded size
- ⌘ Application layer messages may be large
- ⌘ Network packets may be smaller
- ⌘ Splitting larger blocks into smaller ones is segmentation (or fragmentation in TCP/IP)
 - ☒ ATM blocks (cells) are 53 octets (bytes) long
 - ☒ Ethernet blocks (frames) are up to 1526 octets long

Advantages & Disadvantages of Segmentation

⌘ Advantages

- ☒ More efficient error control
- ☒ More equitable access to network facilities with shorter delays
- ☒ Smaller buffers needed
- ☒ More efficient for checkpoints and restart/recovery

⌘ Disadvantages

- ☒ Overheads
- ☒ Increased interrupts at receiver
- ☒ More processing time

Effect of Packet Size on Transmission Time

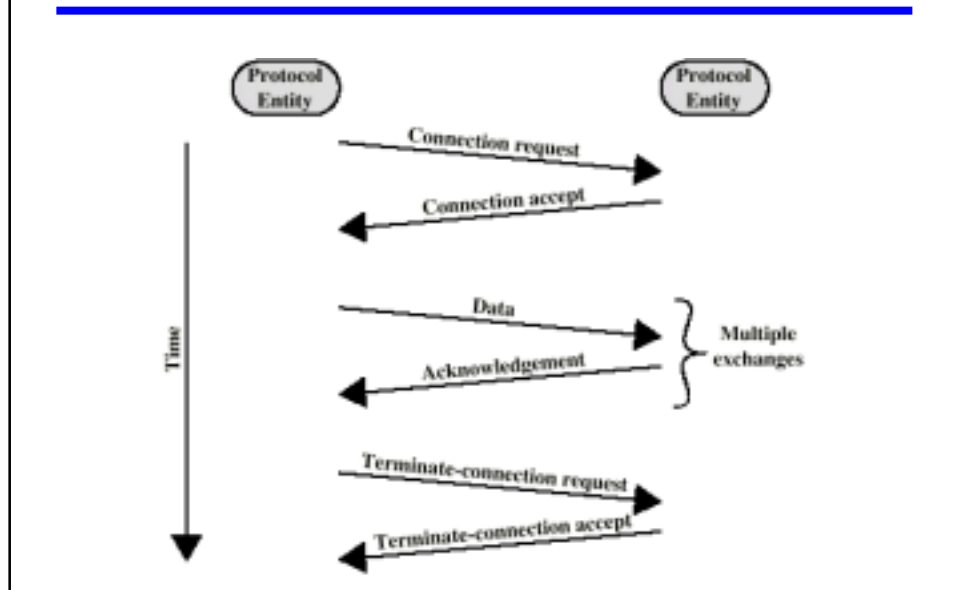
- Message has 40 bytes
- Packet header is 3 bytes
- If packet size is 43 bytes
 - one packet
 - $3 \times 43 = 192$ bytes
- If packet size is 23 bytes
 - two packets
 - $4 \times 23 = 92$ bytes
- If packet size is 11 bytes
 - five packets
 - $7 \times 11 = 77$ bytes
- If packet size is 7 bytes
 - ten packets
 - $12 \times 7 = 84$ bytes



Connection Control

- ⌘ Data transfer is either connectionless or connection-oriented
- ⌘ Connectionless: each PDU is treated independently of other PDUs, e.g. use of Datagram
- ⌘ Connection-oriented (e.g. virtual circuit)
 - ☑ Connection Establishment
 - ☑ Data transfer
 - ☑ Connection termination
- ⌘ Connection interruption and recovery
- ⌘ In connection establishment, protocols negotiate syntax, semantics, and timing
- ⌘ Protocols may allow certain optional features that must be agreed upon by negotiation

Connection Oriented Data Transfer



Datagram

- ⌘ Each packet treated independently
- ⌘ Packets can take any practical route
- ⌘ Packets may arrive out of order
- ⌘ Packets may go missing
- ⌘ Up to receiver to re-order packets and recover from missing packets

Virtual Circuit

- ⌘ Preplanned route established before any packets sent
- ⌘ Call request and call accept packets establish connection (handshake)
- ⌘ Each packet contains a virtual circuit identifier instead of destination address
- ⌘ No routing decisions required for each packet
- ⌘ Clear request to drop circuit
- ⌘ Not a dedicated path

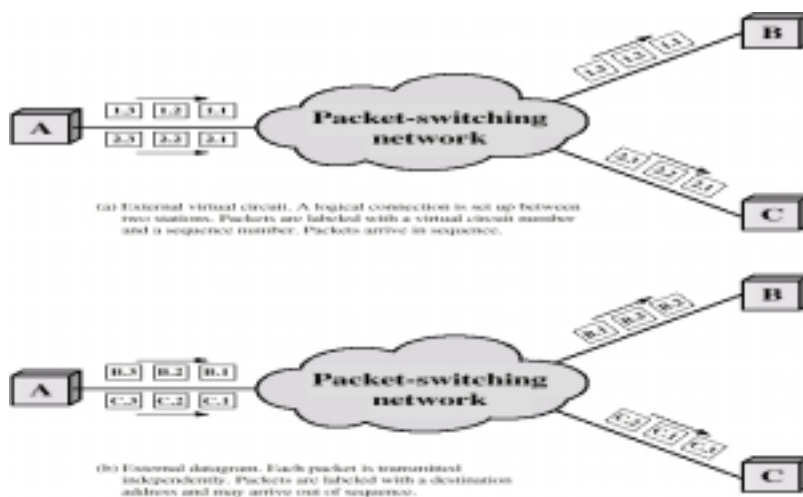
Virtual Circuits vs. Datagram

- ⌘ Virtual circuits
 - ⊠ Network can provide sequencing and error control
 - ⊠ Packets are forwarded more quickly
 - ⊠ No routing decisions to make
 - ⊠ Less reliable
 - ⊠ Loss of a node loses all circuits through that node
- ⌘ Datagram
 - ⊠ No call setup phase
 - ⊠ Better if few packets
 - ⊠ More flexible
 - ⊠ Routing can be used to avoid congested parts of the network

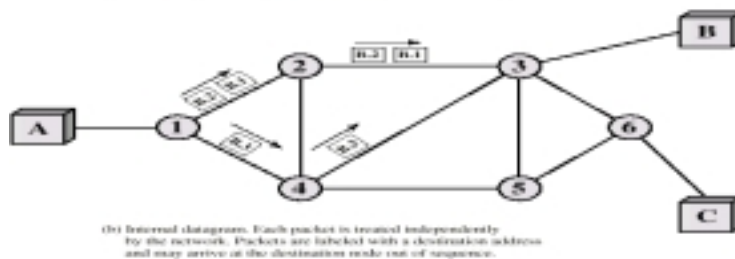
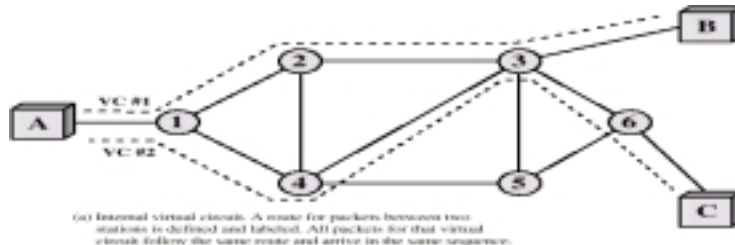
Packet Switching External and Internal Operation

- ⌘ Packet switching - datagrams or virtual circuits
- ⌘ Interface between station and network node
 - ☑ Connection oriented
 - ☑ Station requests logical connection (virtual circuit)
 - ☑ All packets identified as belonging to that connection & sequentially numbered
 - ☑ Network delivers packets in sequence
 - ☑ External virtual circuit service
 - ☑ e.g. X.25
 - ☑ Different from internal virtual circuit operation
 - ☑ Connectionless
 - ☑ Packets handled independently
 - ☑ External datagram service
 - ☑ Different from internal datagram operation

External Virtual Circuit and Datagram Operation



Internal Virtual Circuit and Datagram Operation



Connection Control

- ⌘ Data transfer uses sequencing
- ⌘ Both entities number PDUs and keep track of both incoming and outgoing numbers
- ⌘ Sequence numbers used for
 - ☑ Ordered delivery
 - ☑ Flow control
 - ☑ Error control

Ordered Delivery

- ⌘ PDUs may traverse different paths through network
- ⌘ PDUs may arrive out of order
- ⌘ Sequentially number PDUs to allow for ordering
- ⌘ Problem if sequence numbers repeat after overflow
- ⌘ Have maximum sequence number to be twice maximum number of outstanding PDUs

Flow Control

- ⌘ Done by receiving entity
- ⌘ Limit amount or rate of data
- ⌘ Stop and wait: each PDU must be acknowledged before the next can be sent
- ⌘ Credit systems: allow a number of PDUs to be sent without acknowledgment
 - ☑ Sliding window
- ⌘ Flow control is implemented in several protocols

Error Control

- ⌘ Guard against data loss or damage
- ⌘ Error detection
 - ☑ Sender inserts error detecting code
 - ☑ Receiver checks this code
 - ☑ If OK, acknowledge
 - ☑ If error, discard packet
- ⌘ Retransmission
 - ☑ If no acknowledge in given time, re-transmit
- ⌘ Performed at various levels

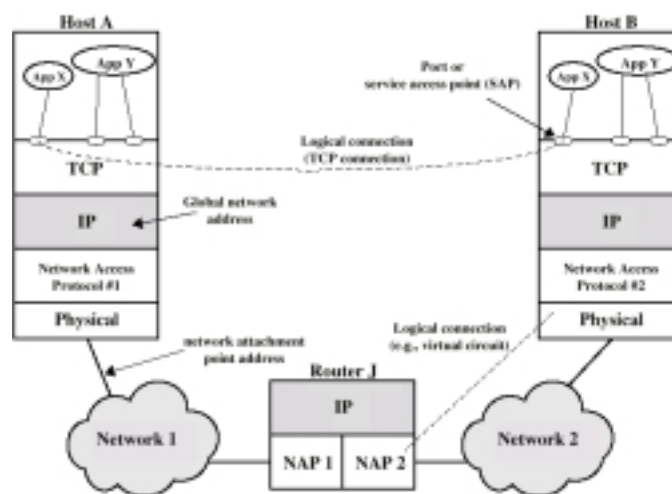
Addressing Issues

- ⌘ Addressing level
- ⌘ Addressing scope
- ⌘ Connection identifiers
- ⌘ Addressing mode

Addressing level

- ⌘ Level in architecture at which entity is named
- ⌘ Unique address for each end system (computer) and intermediate system (router)
 - ☑ IP or internet address (TCP/IP)
 - ☑ Network service access point or NSAP (OSI)
- ⌘ Network level address
 - ☑ IP or internet address (TCP/IP)
 - ☑ Network service access point or NSAP (OSI)
- ⌘ Process within the system
 - ☑ Port number (TCP/IP)
 - ☑ Service access point or SAP (OSI)

Address Concepts



Addressing Scope

⌘ Characteristics of a global address

- ☒ Global nonambiguity
 - ☒ Global address identifies unique system
 - ☒ There is only one system with address X
- ☒ Global applicability
 - ☒ It is possible at any system (any address) to identify any other system (address) by the global address of the other system
 - ☒ Address X identifies that system from anywhere on the network

⌘ Each network must maintain a unique address for each device in the network (network attachment point address)

- ☒ MAC address on IEEE 802 networks
- ☒ X.25 host address

Connection Identifiers

⌘ Connection oriented data transfer (virtual circuits)

⌘ Allocate a connection identifier during the transfer phase

- ☒ Reduced overhead as connection identifiers are shorter than global addresses
 - ☒ In X.25 protocol, connection identifier is a 12-bit virtual circuit number
- ☒ Routing is defined and identified by connection identifier eliminating routing for each PDU
- ☒ Entities may want multiple connections - multiplexing
- ☒ State information
 - ☒ Enables flow control and error control using sequence numbers

Addressing Mode

- ⌘ Usually an address refers to a single system
 - ☒ Unicast address
 - ☒ Sent to one machine or person
- ⌘ May address all entities within a domain
 - ☒ Broadcast
 - ☒ Sent to all machines or users
- ⌘ May address a subset of the entities in a domain
 - ☒ Multicast
 - ☒ Sent to some machines or a group of users

IP Addresses - Class A

- ⌘ 32 bit global internet address
- ⌘ Network part and host part
- ⌘ Class A
 - ☒ Start with binary 0
 - ☒ Network address 7 bits
 - ☒ Host address 24 bits
 - ☒ $2^7 = 128$ class A addresses
 - ☒ Range 1.x.x.x to 126.x.x.x
 - ☒ All 0 reserved
 - ☒ 01111111 (127) reserved for loopback
 - ☒ All allocated

IP Addresses - Class B

- ⌘ Start 10
- ⌘ Network address 14 bits
- ⌘ Host address 16 bits
- ⌘ $2^{14} = 16,384$ class B addresses
- ⌘ Range 128.x.x.x to 191.x.x.x
- ⌘ All allocated

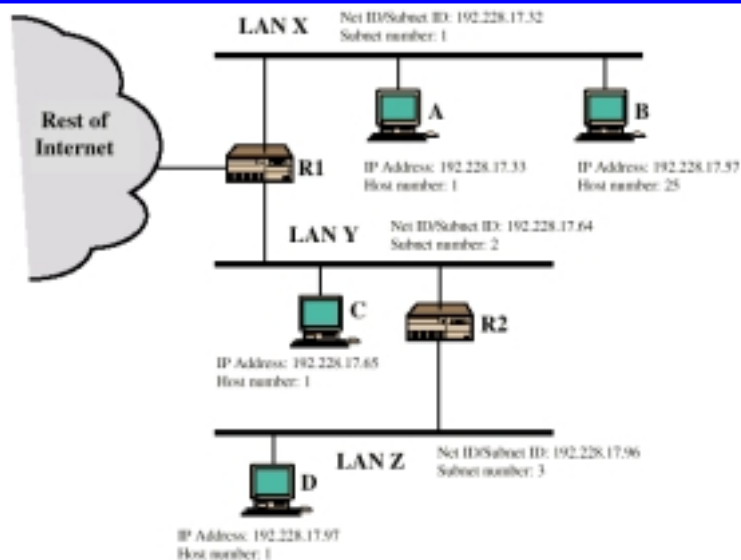
IP Addresses - Class C

- ⌘ Start 110
- ⌘ Network address 21 bits
- ⌘ Host address 8 bits
- ⌘ $2^{21} = 2,097,152$ class C addresses
- ⌘ Range 192.x.x.x to 223.x.x.x
- ⌘ Nearly all allocated
 - ☑ See IPv6

Subnets and Subnet Masks

- ⌘ Allow arbitrary complexity of internetworked LANs within organization
- ⌘ Insulate overall internet from growth of network numbers and routing complexity
- ⌘ To rest of internet site looks like a single network
- ⌘ Each LAN assigned subnet number
- ⌘ Host portion of address partitioned into subnet number and host number
- ⌘ Local routers route within subnetted network
- ⌘ Subnet mask indicates which bits are subnet number and which are host number

Routing Using Subnets



Multiplexing

- ⌘ Combining several signals for transmission on shared medium
- ⌘ Supporting multiple connections on one machine
- ⌘ Mapping of multiple connections at one level to a single connection at another
 - ☑ Carrying a number of connections on one fiber optic cable
 - ☑ Aggregating or bonding ISDN lines to gain bandwidth

Multiplexing

- ⌘ Upward Multiplexing
 - ☑ Multiple higher-level connections are multiplexed on a single lower-level connection
 - ☑ Connecting your PC to ISP for multiple applications, including web, email, ftp, telnet
- ⌘ Downward Multiplexing
 - ☑ Split a single higher-level connection over a number of lower-level connections
 - ☑ Useful for reliability, performance, or efficiency

Transmission Services

⌘ Priority

- ☒ e.g. control messages
- ☒ Assigned on a message basis on connection basis

⌘ Quality of service

- ☒ Minimum acceptable throughput
- ☒ Maximum acceptable delay

⌘ Security

- ☒ Access restrictions

OSI - The Model

⌘ A seven layer model

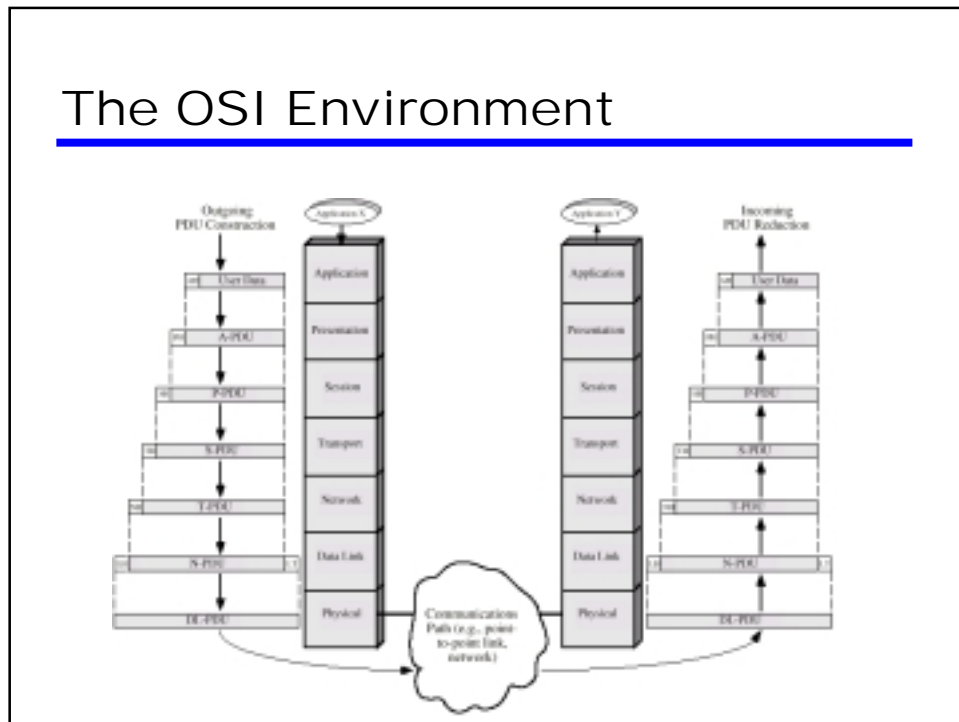
⌘ Each layer performs a subset of the required communication functions

⌘ Each layer relies on the next lower layer to perform more primitive functions

⌘ Each layer provides services to the next higher layer

⌘ Changes in one layer should not require changes in other layers

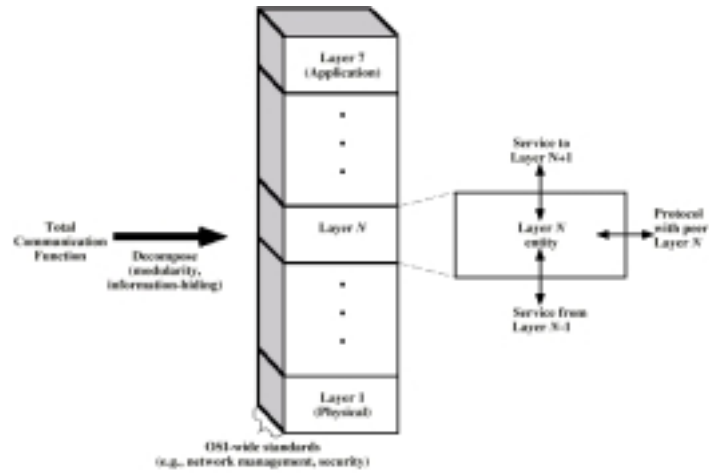
The OSI Environment



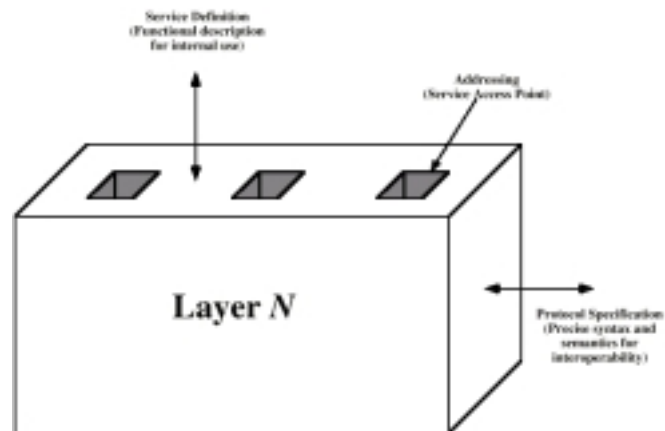
OSI as Framework for Standardization

- ⌘ Defines functions at each layer and facilitates standards-making process.
- ⌘ Standards at each layer can be developed independently and simultaneously.
- ⌘ Changes in standards in one layer need not affect other layers
 - ☑ Well defined boundaries (interface) between layers

OSI as Framework for Standardization



Layer Specific Standards



Elements of Standardization

⌘ Protocol specification

- ☒ Operates between two systems on same layer
- ☒ May involve different operating systems
- ☒ Protocol specification must be precise
 - ☒ Format of data units
 - ☒ Semantics of all fields
 - ☒ allowable sequence of PDUs

⌘ Service definition

- ☒ Functional description of what is provided, not how it is provided

⌘ Addressing

- ☒ Referenced by SAPs
- ☒ Allows multiplexing from higher layer

Service Primitives and Parameters

⌘ Services between adjacent layers expressed in terms of primitives and parameters

⌘ Primitive specifies the function to be performed

⌘ Parameters used to pass data and control information

⌘ Four types of primitives

- ☒ Request
- ☒ Indication
- ☒ Response
- ☒ Confirm

Service Primitives

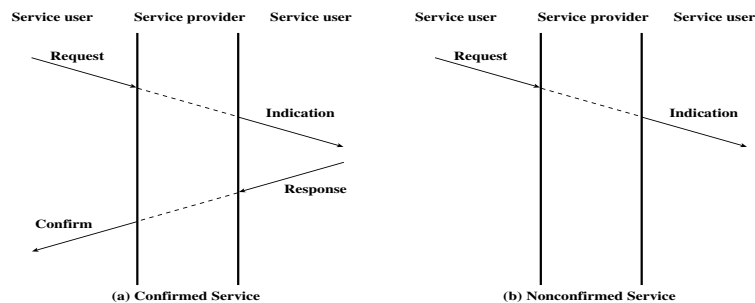


Figure 2.9 Time Sequence Diagrams for Service Primitives

OSI Layers - Physical

⌘ Physical interface between devices

☒ Mechanical

- ☒ Physical properties of interface to transmission medium
- ☒ Specifications of pluggable connector

☒ Electrical

- ☒ Representation of bits in terms of voltage levels
- ☒ Data transmission rates

☒ Functional

- ☒ Functions of individual circuits of physical interface

☒ Procedural

- ☒ Sequence of events by which bit streams are exchanged

⌘ Examples: [EIA-232-F](#), portions of ISDN and LAN standards

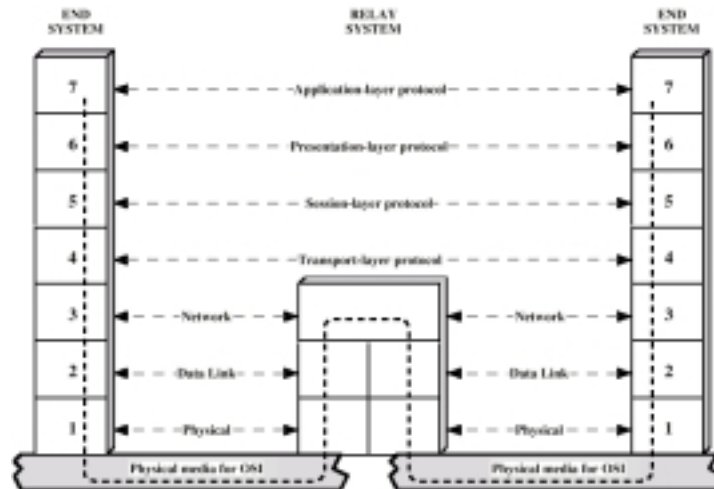
OSI Layers – Data Link

- ⌘ Makes physical link reliable through error detection and control
- ⌘ Activates, maintains and deactivates link
- ⌘ Higher layers may assume error free transmission
- ⌘ Communication through a number of data links require higher layers to perform some error control
- ⌘ Examples: [High-level Data link control \(HDLC\)](#)

OSI Layers - Network

- ⌘ Transfers information across communications network performing switching & routing
- ⌘ Hides data transmission and switching technologies
- ⌘ Not needed on direct links
- ⌘ Highest layer in a network node
- ⌘ System interacts with network
 - ☒ Specification of destination address
 - ☒ Request for network services like priority
- ⌘ Examples: packet level of X.25 standard

Use of a Relay



OSI Layers - Transport

- ⌘ Mechanisms for exchange of data between end systems
- ⌘ Ensures data delivered error free, in sequence, no losses, no duplicates
- ⌘ May optimize the use of network services
- ⌘ Provides quality of service based on acceptable error rates, maximum delay, priority, security
- ⌘ Size and complexity depend on reliability of underlying layers
- ⌘ Examples: connection-oriented TCP (transmission control protocol), connectionless UDP (user datagram protocol)

OSI Layers - Session

- ⌘ Control of dialogues between applications
- ⌘ Dialogue discipline
 - ☑ Full duplex or half duplex
- ⌘ Grouping
 - ☑ Mark data to define groups of data
- ⌘ Recovery
 - ☑ Checkpoint to allow retransmission of all data since last checkpoint due to failure

OSI Layers – Presentation & Application

- ⌘ Presentation
 - ☑ Data formats and coding
 - ☑ Defines syntax used between application entities
 - ☑ Provides for selection and modification of representation used
 - ☑ Data compression and encryption
- ⌘ Application
 - ☑ Interface between application programs and OSI environment

TCP/IP Protocol Suite

- ⌘ Dominant commercial protocol architecture
- ⌘ Specified and extensively used before OSI
- ⌘ Developed by research funded US Department of Defense
- ⌘ Used by the Internet

TCP/IP Approach

- ⌘ Modular and hierarchical like the OSI model
- ⌘ Descriptive in nature compared to prescriptive nature of OSI
 - ☒ OSI dictates that protocols in a layer perform certain functions
 - ☒ In TCP/IP, it is possible to have two protocols in same layer with different functionality
- ⌘ Does not require strict use of all layers
 - ☒ Application level protocols may directly run on top of IP

TCP/IP Protocol Architecture

⌘ Application Layer

- ☒ Communication between processes or applications

⌘ End to end or transport layer (TCP/UDP/...)

- ☒ End to end transfer of data
- ☒ May include reliability mechanism (TCP)
- ☒ Hides detail of underlying network
- ☒ Implemented in end systems only

⌘ Internet Layer (IP)

- ☒ Routing of data
- ☒ Implemented in all end systems and routers

TCP/IP Protocol Architecture

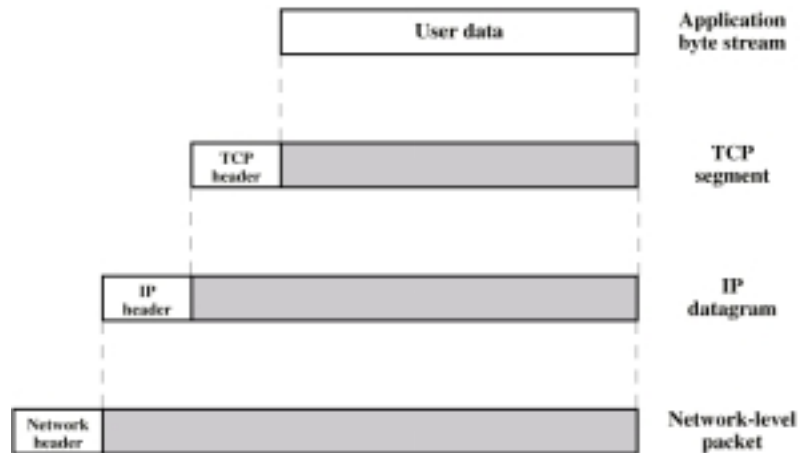
⌘ Network Layer

- ☒ Logical interface between end system and network

⌘ Physical Layer

- ☒ Transmission medium
- ☒ Signal rate and encoding

PDU in TCP/IP



Operation of TCP and IP

⌘ Two levels of addressing

- ☑ Unique host address over global internet, used by IP
- ☑ Unique process (port) address within host, used by TCP

⌘ TCP header

- ☑ Destination port: address to whom data to be delivered
- ☑ Sequence number: used by destination TCP to reorder segments
- ☑ Checksum: code to check error during transmission

Operation of TCP and IP

⌘ IP datagram

- ☑ Created by adding IP header to each segment
- ☑ Header includes destination host address
- ☑ Presented to network access layer for transmission

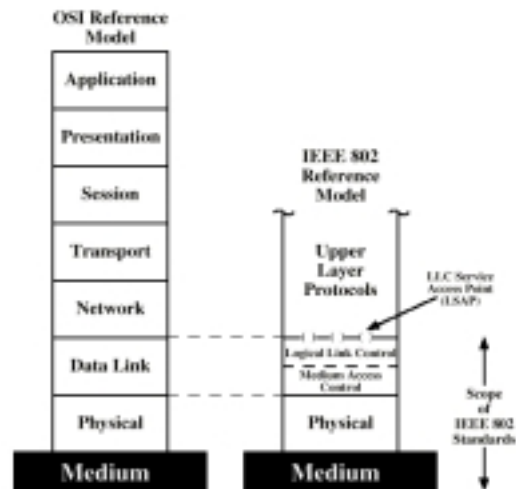
⌘ Packet or frame

- ☑ Created by network access layer by adding its header
- ☑ Header contains information for network to transfer data across it
 - ☑ Destination network address
 - ☑ Facilities request

LAN Protocol Architecture

- ⌘ Lower layers of OSI model
- ⌘ IEEE 802 reference model
- ⌘ Physical
- ⌘ Logical link control (LLC)
- ⌘ Media access control (MAC)

IEEE 802 v OSI



IEEE 802 Layers

⌘ Physical

- ☑ Encoding/decoding
- ☑ Preamble generation/removal
- ☑ Bit transmission/reception
- ☑ Transmission medium and topology

⌘ Logical Link Control

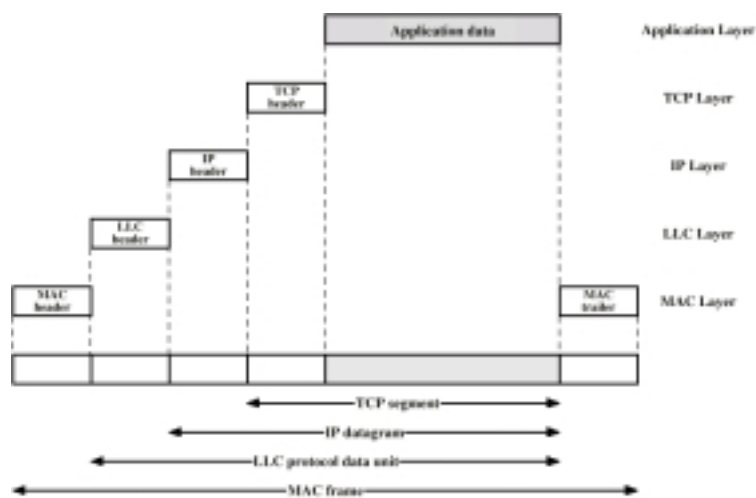
- ☑ Interface to higher levels
- ☑ Flow and error control

IEEE 802 Layers

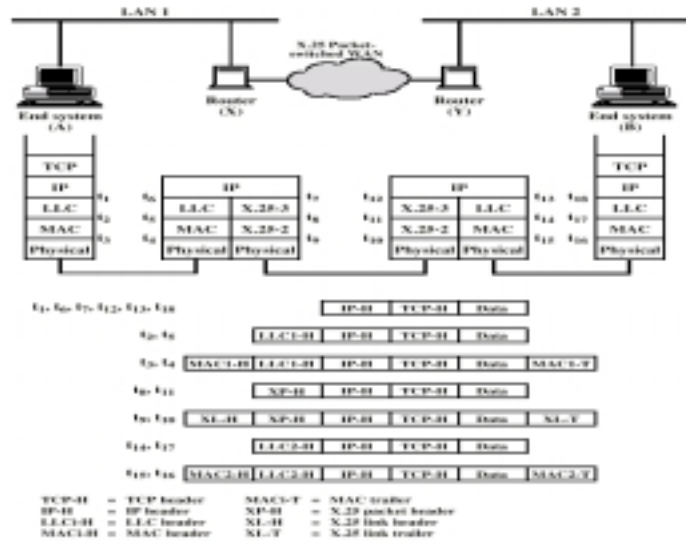
⌘ Media Access Control

- ☒ Assembly of data into frame with address and error detection fields
- ☒ Disassembly of frame
 - ☒ Address recognition
 - ☒ Error detection
- ☒ Govern access to transmission medium
- ☒ Not found in traditional layer 2 data link control
- ☒ For the same LLC, several MAC options may be available

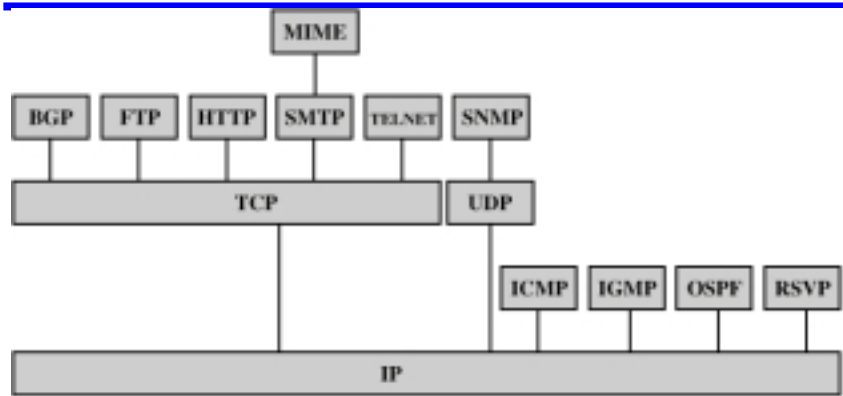
LAN Protocols in Context



IP Operation



Some Protocols in TCP/IP Suite



- BGP = Border Gateway Protocol
- FTP = File Transfer Protocol
- HTTP = Hypertext Transfer Protocol
- ICMP = Internet Control Message Protocol
- IGMP = Internet Group Management Protocol
- IP = Internet Protocol
- MIME = Multi-Purpose Internet Mail Extension
- OSPF = Open Shortest Path First
- RSVP = Resource Reservation Protocol
- SMTP = Simple Mail Transfer Protocol
- SNMP = Simple Network Management Protocol
- TCP = Transmission Control Protocol
- UDP = User Datagram Protocol

Simple Mail Transfer Protocol (SMTP)

- ⌘ Basic email utility
- ⌘ Mechanism to transfer messages across hosts
- ⌘ Features include mailing lists, return receipts, and forwarding
- ⌘ Does not specify message creation; just transfer of message using TCP

File Transfer Protocol (FTP)

- ⌘ Transfer files across systems under user commands
- ⌘ Accommodate both text and binary files
- ⌘ Upon request, sets up connection for exchange of control messages
- ⌘ Upon approval, a second TCP opened for actual data transfer
 - ☑ Avoids overhead of control information
- ⌘ After file transfer complete, control connection signals completion and accepts new commands

Telnet

- ⌘ Remote login capability
- ⌘ Designed to work with simple scroll-mode terminals
- ⌘ Implemented in two modules
 - ☒ User telnet
 - ☒ Interacts with terminal I/O module to communicate with local terminal
 - ☒ Converts between characteristics of real terminals and network standards
 - ☒ Server telnet
 - ☒ Interacts with an application acting as a terminal handler
 - ☒ Makes remote terminal appear as local to application
- ⌘ Traffic between user and server telnet carried on TCP

Internetworking Terms (1)

- ⌘ Communications Network
 - ☒ Facility that provides data transfer service
- ⌘ An internet
 - ☒ Collection of communications networks interconnected by bridges and/or routers
- ⌘ The Internet - note upper case I
 - ☒ The global collection of thousands of individual machines and networks
- ⌘ Intranet
 - ☒ Corporate internet operating within the organization
 - ☒ Uses Internet (TCP/IP and http) technology to deliver documents and resources

Internetworking Terms (2)

⌘ End System (ES)

- ☒ Device attached to one of the networks of an internet
- ☒ Supports end-user applications or services

⌘ Intermediate System (IS)

- ☒ Device used to connect two networks
- ☒ Permits communication between end systems attached to different networks

Internetworking Terms (3)

⌘ Bridge

- ☒ IS used to connect two LANs using similar LAN protocols
- ☒ Address filter passing on packets to the required network only
- ☒ OSI layer 2 (Data Link)

⌘ Router

- ☒ Connects two (possibly dissimilar) networks
- ☒ Uses internet protocol present in each router and end system
- ☒ OSI Layer 3 (Network)